

3850 Group (Spec. H) User's Manual

RENESAS 8-BIT CISC SINGLE-CHIP MICROCOMPUTER 740 FAMILY / 38000 SERIES

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REVISION HISTORY

3850 Group (Spec. H) User's Manual

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1.0	Aug. 30, 2001	_	First edition issued
1.1	Sep. 10, 2001	3-5	Limits and test conditions into Table 3.1.5 are partly added.
1.02	Aug. 29, 2003	1-6 1-7 1-37 1-38 2-103 2-104 2-105 2-106 2-107 3-4 3-20 3-22 3-24	 Fig. 4 is partly revised. Table 2 is partly added. Note of Table 3 is added. Fig. 42 is partly revised. Fig. 43 is partly revised. Clause name of "2.11 Flash memory mode" is revised. Notes of Fig. 2.11.3 are partly revised. Table 2.11.2 is partly revised. Explanations of "[Beginning procedure]" of "2.11.6 CPU rewrite mode" are partly added. Explanations of this page are added. Parameter of Table 3.1.4 is partly revised. Fig. 3.2.16 is partly revised. Fig. 3.2.20 is partly revised.
1.03	Sep. 18, 2003	1-51	Fig. 52 is partly revised.

Preface

This user's manual describes Renesas's CMOS 8-bit microcomputers 3850 Group (Spec. H).

After reading this manual, the user should have a through knowledge of the functions and features of the 3850 Group (Spec. H), and should be able to fully utilize the product. The manual starts with specifications and ends with application examples. For details of software, refer to the "740 Family Software Manual".

The user who is using the 3850 Group (standard) needs to refer to not this manual but "3850/3851 Group User's Manual".

BEFORE USING THIS MANUAL

This user's manual consists of the following three chapters. Refer to the chapter appropriate to your conditions, such as hardware design or software development. Chapter 3 also includes necessary information for systems development. You must refer to that chapter.

1. Organization

• CHAPTER 1 HARDWARE

This chapter describes features of the microcomputer and operation of each peripheral function.

• CHAPTER 2 APPLICATION

This chapter describes usage and application examples of peripheral functions, based mainly on setting examples of relevant registers.

• CHAPTER 3 APPENDIX

This chapter includes necessary information for systems development using the microcomputer, such as the electrical characteristics, the notes, and the list of registers.

*For the mask ROM confirmation form, the ROM programming confirmation form, and the mark specifications, refer to the "Renesas Technology Corp." Homepage (http://www.renesas.com/en/rom).

2. Structure of register

The figure of each register structure describes its functions, contents at reset, and attributes as follows :

		Bits Contents immediately	Bit attribu	ote 2) Ites	
b7 b6 b5 b4 b3 b2 b1 b0	СРС	mode register (CPUM) [Address : :	````	Ľ	
	в	Name	Function	At reset	(R W)
		Processor mode bits	0 : Single-chip mode	0	00
	2	Stack page selection bit	11:J 0:0 page 1:1 page	0	0.0
	3	Nothing arranged for these bits. The		0	0:×
	4	bits. When these bits are read out,	the contents are "0."	0	0 ×
	5	Fix this bit to "0."		1	00
	6	Main clock (XIN-XOUT) stop bit	0 : Operating 1 : Stopped	*	00
i	7	Internal system clock selection bit	0 : XIN-XOUT selected 1 : XCIN-XCOUT selected	*	0 0
0"0" at 1"1" at ? Unde	nedia rese rese fined	itely after reset release t release	ed for control of the corr	espondii	ng function
Note 2: Bit attributes.		The attributes of control register bits only and read and write. In the figur			
	С	Read enabled	VWrite)Write enabled KWrite disabled		

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CHAPTER 1 HARDWARE

DESCRIPTION **FEATURES APPLICATION PIN CONFIGURATION** FUNCTIONAL BLOCK **PIN DESCRIPTION** PART NUMBERING **GROUP EXPANSION** FUNCTIONAL DESCRIPTION NOTES ON PROGRAMMING NOTES ON USAGE DATA REQUIRED FOR MASK ORDERS DATA REQUIRED FOR One Time PROM PROGRAMMING ORDERS ROM PROGRAMMING METHOD FUNCTIONAL DESCRIPTION SUPPLEMENT

HARDWARE DESCRIPTION/FEATURES/APPLICATION/PIN CONFIGURATION

DESCRIPTION

The 3850 group (spec. H) is the 8-bit microcomputer based on the 740 family core technology.

The 3850 group (spec. H) is designed for the household products and office automation equipment and includes serial I/O functions, 8-bit timer, and A-D converter.

FEATURES

Basic machine-language instructions	
\bullet Minimum instruction execution time 0.5 μs	

(at 8 MHz oscillation frequency)

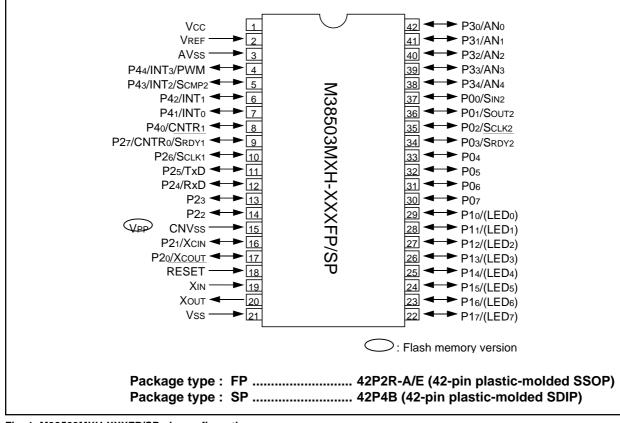
Memory size	
ROM	8K to 32K bytes
RAM	512 to 1024 bytes
• Programmable input/output ports .	
Interrupts	15 sources, 14 vectors
•Timers	8-bit X 4
•Serial I/O1 8-bit X 1(UART or Clock-synchronized)
●Serial I/O2	8-bit X 1(Clock-synchronized)
• PWM	8-bit X 1
●A-D converter	10-bit X 5 channels
Watchdog timer	16-bit X 1
Clock generating circuit	Built-in 2 circuits
(connect to external ceramic resona	tor or quartz-crystal oscillator)

Power source voltage
In high-speed mode 4.0 to 5.5 V
(at 8 MHz oscillation frequency)
In middle-speed mode 2.7 to 5.5 V
(at 8 MHz oscillation frequency)
In low-speed mode 2.7 to 5.5 V $$
(at 32 kHz oscillation frequency)
Power dissipation
In high-speed mode34 mW
(at 8 MHz oscillation frequency, at 5 V power source voltage)
In low-speed mode
Except M38507F8FP/SP 60 μW
M38507F8FP/SP 450 μ W
(at 32 kHz oscillation frequency, at 3 V power source voltage)
●Operating temperature range20 to 85°C

APPLICATION

Office automation equipment, FA equipment, Household products, Consumer electronics, etc.

PIN CONFIGURATION (TOP VIEW)





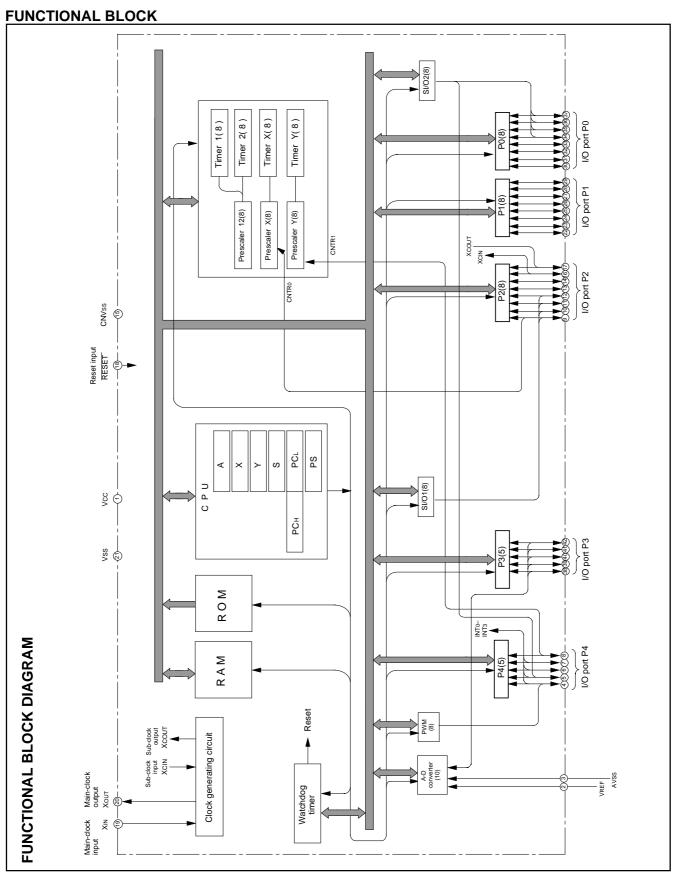


Fig. 2 Functional block diagram

PIN DESCRIPTION

PIN DESCRIPTION

Table 1 Pin description

Pin	Name	Functions	Function except a port function
Vcc, Vss	Power source	•Apply voltage of 2.7 V – 5.5 V to Vcc, and 0 V to Vss.	
CNVss	CNVss input	•This pin controls the operation mode of the chip. •Normally connected to Vss.	
RESET	Reset input	•Reset input pin for active "L".	
Xin	Clock input	Input and output pins for the clock generating circuit. Connect a ceramic resonator or quartz-crystal oscillator the oscillation frequency.	between the XIN and XOUT pins to set
Хоит	Clock output	•When an external clock is used, connect the clock sour pin open.	rce to the XIN pin and leave the XOUT
P00/SIN2 P01/SOUT2 P02/SCLK2 P03/SRDY2	I/O port P0	 •8-bit CMOS I/O port. •I/O direction register allows each pin to be individually programmed as either input or output. •CMOS compatible input level. 	Serial I/O2 function pin
P04–P07		•CMOS 3-state output structure.	-
P10–P17	I/O port P1	•P10 to P17 (8 bits) are enabled to output large current fo	r LED drive.
P20/XCOUT P21/XCIN P22	-	 •8-bit CMOS I/O port. •I/O direction register allows each pin to be individually programmed as either input or output. 	Sub-clock generating circuit I/O pins (connect a resonator)
P23	I/O port P2	•CMOS compatible input level.	
P24/RxD P25/TxD P26/SCLK1	-	 •P20, P21, P24 to P27: CMOS3-state output structure. •P22, P23: N-channel open-drain structure. 	Serial I/O1 function pin
P27/CNTR0/ SRDY1	-		Serial I/O1 function pin/ Timer X function pin
P30/AN0– P34/AN4	I/O port P3	 •8-bit CMOS I/O port with the same function as port P0. •CMOS compatible input level. •CMOS 3-state output structure. 	• A-D converter input pin
P40/CNTR1		•8-bit CMOS I/O port with the same function as port P0.	Timer Y function pin
P41/INT0 P42/INT1	I/O port P4	•CMOS compatible input level. •CMOS 3-state output structure.	Interrupt input pins
P43/INT2/SCMP2			Interrupt input pinSCMP2 output pin
P44/INT3/PWM			Interrupt input pinPWM output pin

HARDWARE PART NUMBERING

PART NUMBERING

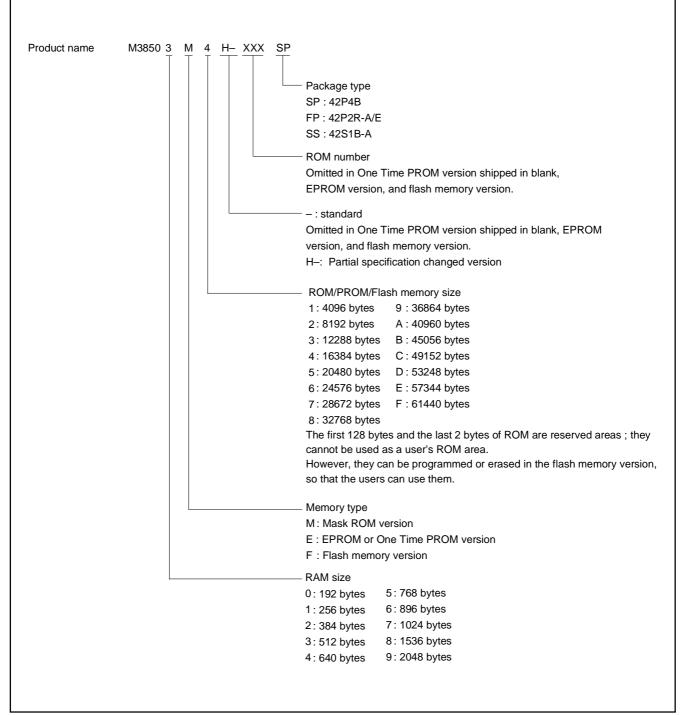


Fig. 3 Part numbering

GROUP EXPANSION

GROUP EXPANSION

Renesas Technology plans to expand the 3850 group (spec. H) as follows.

Memory Type

Support for mask ROM, One Time PROM, and flash memory versions.

Memory Size

Flash memory size	. 32 K bytes
One Time PROM size	. 24 K bytes
Mask ROM size 8 K to	32 K bytes

RAM size	512 to 1 K bytes
Packages	
42P4B	42-pin shrink plastic-molded DIP
42P2R-A/E	42-pin plastic-molded SSOP
42S1B-A	42-pin shrink ceramic DIP (EPROM version)

Memory Expansion Plan

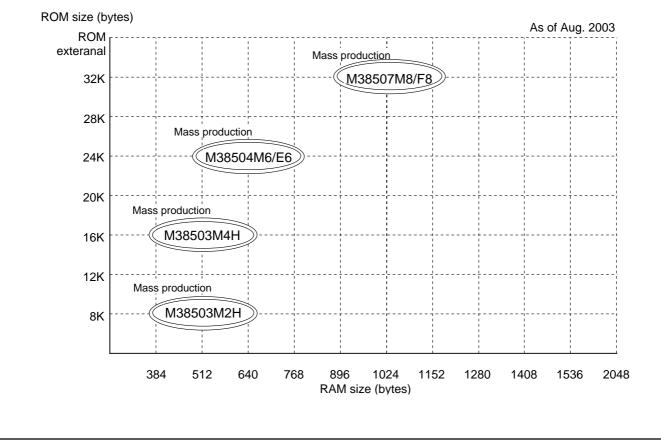


Fig. 4 Memory expansion plan

As of Aug. 2003

Currently support products are listed below.

Table 2 Support products

Product name	ROM size (bytes) ROM size for User in ()	RAM size (bytes)	Package	Remarks	
M38503M2H-XXXSP	8192	512	42P4B	Mask ROM version	
M38503M2H-XXXFP	(8062)		42P2R-A/E	Mask ROM version	
M38503M4H-XXXSP	16384	512	424P4B	Mask ROM version	
M38503M4H-XXXFP	(16254)		42P2R-A/E	Mask ROM version	
M38504M6-XXXSP		640	424P4B	Mask ROM version	
M38504E6-XXXSP				One Time PROM version	
M38504E6SP	24576 (24446)			One Time PROM version (blank)	
M38504E6SS			42S1B-A	EPROM version	
M38504M6-XXXFP	(24440)		42P2R-A/E	Mask ROM version	
M38504E6-XXXFP				One Time PROM version	
M38504E6FP				One Time PROM version (blank)	
M38507M8-XXXSP			42P4B	Mark DOM was in a	
M38507M8-XXXFP	32768		42P2R-A/E	Mask ROM version	
M38507F8SP	(32638)	1024	424P4B		
M38507F8FP			42P2R-A/E	Flash memory version	

Table 33850 group (standard) and 3850 group (spec. H)corresponding products

3850 group (spec. H)
M38503M2H-XXXFP/SP
M38503M4H-XXXFP/SP
M38504M6-XXXFP/SP
M38504E6-XXXFP/SP
M38504E6FP/SP
M38504E6SS
M38507M8-XXXFP/SP
M38507F8FP/SP

Note: The user who is using the 3850 Group (standard) needs to refer to not this manual but "3850/3851 Group User's Manual".

Table 4 Differences between 3850 group (standard) and 3850 group (spec. H)

	3850 group (standard)	3850 group (spec. H)	
Serial I/O	1: Serial I/O (UART or Clock-synchronized)	2: Serial I/O1 (UART or Clock-synchronized)	
		Serial I/O2 (Clock-synchronized)	
A-D converter	Unserviceable in low-speed mode	Serviceable in low-speed mode	
Large current port	5: P13–P17	8: P10–P17	

Notes on differences between 3850 group (standard) and 3850 group (spec. H)

(1) The absolute maximum ratings of 3850 group (spec. H) is smaller than that of 3850 group (standard).

•Power source voltage Vcc = -0.3 to 6.5 V

•CNVss input voltage VI = -0.3 to Vcc +0.3 V

- (2) The oscillation circuit constants of XIN-XOUT, XCIN-XCOUT may be some differences between 3850 group (standard) and 3850 group (spec. H).
- (3) Do not write any data to the reserved area and the reserved bit. (Do not change the contents after reset.)
- (4) Fix bit 3 of the CPU mode register to "1".
- (5) Be sure to perform the termination of unused pins.

FUNCTIONAL DESCRIPTION

FUNCTIONAL DESCRIPTION CENTRAL PROCESSING UNIT (CPU)

The 3850 group (spec. H) uses the standard 740 Family instruction set. Refer to the table of 740 Family addressing modes and machine instructions or the 740 Family Software Manual for details on the instruction set.

- Machine-resident 740 Family instructions are as follows:
 - The FST and SLW instructions cannot be used.
 - The STP, WIT, MUL, and DIV instructions can be used.

[Accumulator (A)]

The accumulator is an 8-bit register. Data operations such as data transfer, etc., are executed mainly through the accumulator.

[Index Register X (X)]

The index register X is an 8-bit register. In the index addressing modes, the value of the OPERAND is added to the contents of register X and specifies the real address.

[Index Register Y (Y)]

The index register Y is an 8-bit register. In partial instruction, the value of the OPERAND is added to the contents of register Y and specifies the real address.

[Stack Pointer (S)]

The stack pointer is an 8-bit register used during subroutine calls and interrupts. This register indicates start address of stored area (stack) for storing registers during subroutine calls and interrupts. The low-order 8 bits of the stack address are determined by the contents of the stack pointer. The high-order 8 bits of the stack address are determined by the stack page selection bit. If the stack page selection bit is "0", the high-order 8 bits becomes "0016". If the stack page selection bit is "1", the high-order 8 bits becomes "0116".

The operations of pushing register contents onto the stack and popping them from the stack are shown in Figure 6.

Store registers other than those described in Figure 6 with program when the user needs them during interrupts or subroutine calls.

[Program Counter (PC)]

The program counter is a 16-bit counter consisting of two 8-bit registers PCH and PCL. It is used to indicate the address of the next instruction to be executed.

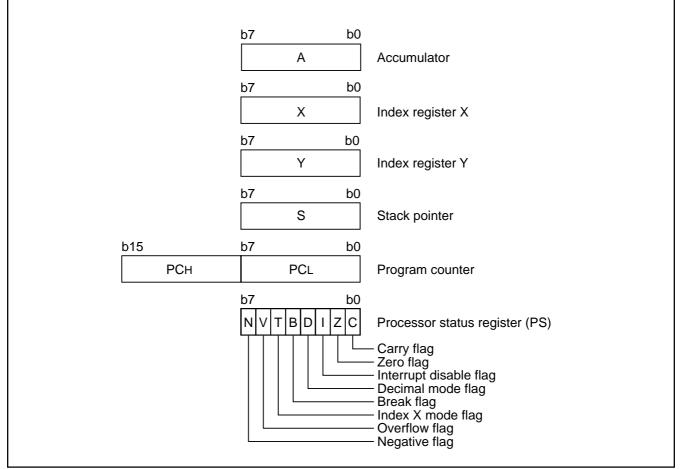


Fig. 5 740 Family CPU register structure

FUNCTIONAL DESCRIPTION

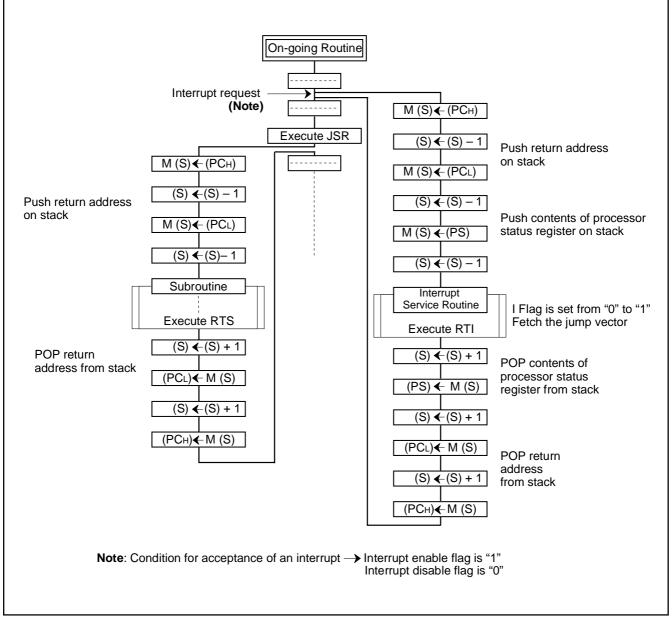


Fig. 6 Register push and pop at interrupt generation and subroutine call

Table 5 Push and pop instructions of accumulator or processor status register

	Push instruction to stack	Pop instruction from stack		
Accumulator	PHA	PLA		
Processor status register	PHP	PLP		

FUNCTIONAL DESCRIPTION

[Processor status register (PS)]

The processor status register is an 8-bit register consisting of 5 flags which indicate the status of the processor after an arithmetic operation and 3 flags which decide MCU operation. Branch operations can be performed by testing the Carry (C) flag, Zero (Z) flag, Overflow (V) flag, or the Negative (N) flag. In decimal mode, the Z, V, N flags are not valid.

•Bit 0: Carry flag (C)

The C flag contains a carry or borrow generated by the arithmetic logic unit (ALU) immediately after an arithmetic operation. It can also be changed by a shift or rotate instruction.

•Bit 1: Zero flag (Z)

The Z flag is set if the result of an immediate arithmetic operation or a data transfer is "0", and cleared if the result is anything other than "0".

•Bit 2: Interrupt disable flag (I)

The I flag disables all interrupts except for the interrupt generated by the BRK instruction.

Interrupts are disabled when the I flag is "1".

•Bit 3: Decimal mode flag (D)

The D flag determines whether additions and subtractions are executed in binary or decimal. Binary arithmetic is executed when this flag is "0"; decimal arithmetic is executed when it is "1". Decimal correction is automatic in decimal mode. Only the ADC

•Bit 4: Break flag (B)

The B flag is used to indicate that the current interrupt was generated by the BRK instruction. The BRK flag in the processor status register is always "0". When the BRK instruction is used to generate an interrupt, the processor status register is pushed onto the stack with the break flag set to "1".

•Bit 5: Index X mode flag (T)

When the T flag is "0", arithmetic operations are performed between accumulator and memory. When the T flag is "1", direct arithmetic operations and direct data transfers are enabled between memory locations.

•Bit 6: Overflow flag (V)

The V flag is used during the addition or subtraction of one byte of signed data. It is set if the result exceeds +127 to -128. When the BIT instruction is executed, bit 6 of the memory location operated on by the BIT instruction is stored in the overflow flag. •Bit 7: Negative flag (N)

The N flag is set if the result of an arithmetic operation or data transfer is negative. When the BIT instruction is executed, bit 7 of the memory location operated on by the BIT instruction is stored in the negative flag.

Table 6 Set and clear instructions of each bit of processor status register

	C flag	Z flag	I flag	D flag	B flag	T flag	V flag	N flag
Set instruction	SEC	-	SEI	SED	_	SET	-	_
Clear instruction	CLC	_	CLI	CLD	_	CLT	CLV	_

[CPU Mode Register (CPUM)] 003B16

The CPU mode register contains the stack page selection bit, etc. The CPU mode register is allocated at address 003B16.

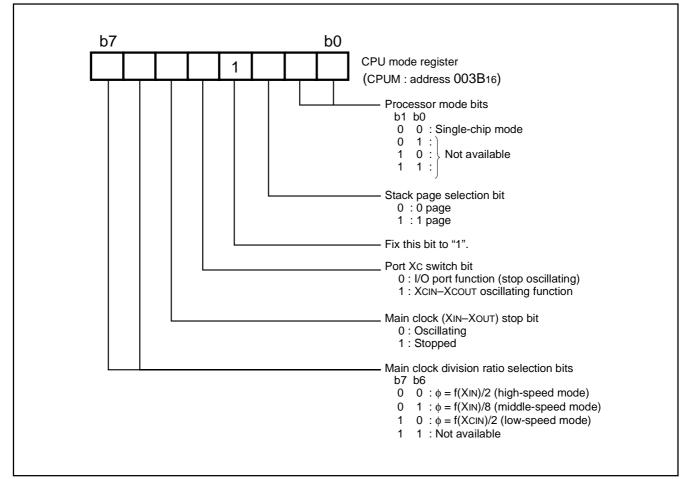


Fig. 7 Structure of CPU mode register

FUNCTIONAL DESCRIPTION

MEMORY

Special Function Register (SFR) Area

The Special Function Register area in the zero page contains control registers such as I/O ports and timers.

RAM

RAM is used for data storage and for stack area of subroutine calls and interrupts.

ROM

The first 128 bytes and the last 2 bytes of ROM are reserved for device testing and the rest is user area for storing programs.

Interrupt Vector Area

The interrupt vector area contains reset and interrupt vectors.

Zero Page

Access to this area with only 2 bytes is possible in the zero page addressing mode.

Special Page

Access to this area with only 2 bytes is possible in the special page addressing mode.

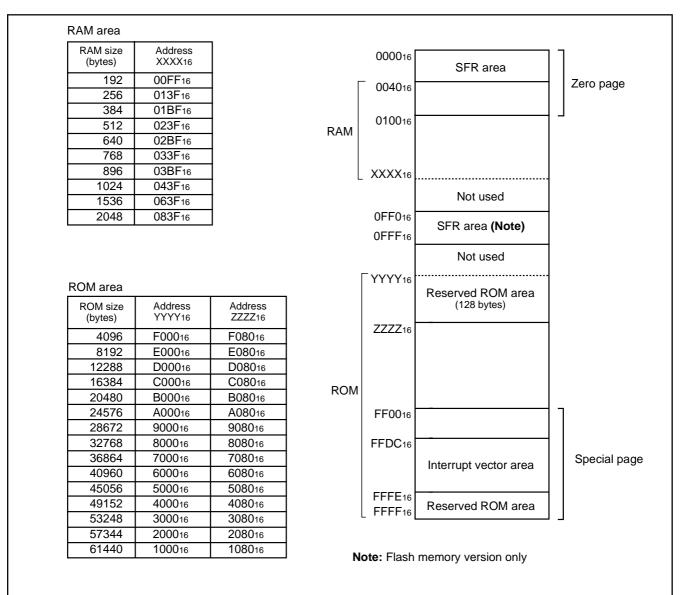


Fig. 8 Memory map diagram

FUNCTIONAL DESCRIPTION

000016	Port P0 (P0)	002016	Prescaler 12 (PRE12)
000116	Port P0 direction register (P0D)	002116	Timer 1 (T1)
000216	Port P1 (P1)	002216	Timer 2 (T2)
000316	Port P1 direction register (P1D)	002316	Timer XY mode register (TM)
000416	Port P2 (P2)	002416	Prescaler X (PREX)
000516	Port P2 direction register (P2D)	002516	Timer X (TX)
000616	Port P3 (P3)	002616	Prescaler Y (PREY)
000716	Port P3 direction register (P3D)	002716	Timer Y (TY)
000816	Port P4 (P4)	002816	Timer count source selection register (TCSS)
000916	Port P4 direction register (P4D)	002916	
000A16		002A16	
000B16		002B16	Reserved *
000C16		002C16	Reserved *
000D16		002D16	Reserved *
000E16		002E16	Reserved *
000F16		002F16	Reserved *
001016		003016	Reserved *
001116		003116	Reserved *
001216	Reserved *	003216	
001316	Reserved *	003316	
001416	Reserved *	003416	A-D control register (ADCON)
001516	Serial I/O2 control register 1 (SIO2CON1)	003516	A-D conversion low-order register (ADL)
001616	Serial I/O2 control register 2 (SIO2CON2)	003616	A-D conversion high-order register (ADH)
001716	Serial I/O2 register (SIO2)	003716	Reserved *
001816	Transmit/Receive buffer register (TB/RB)	003816	MISRG
001916	Serial I/O1 status register (SIOSTS)	003916	Watchdog timer control register (WDTCON)
001A16	Serial I/O1 control register (SIOCON)	003A16	Interrupt edge selection register (INTEDGE)
001B16	UART control register (UARTCON)	003B16	CPU mode register (CPUM)
001C16	Baud rate generator (BRG)	003C16	Interrupt request register 1 (IREQ1)
001D16	PWM control register (PWMCON)	003D16	Interrupt request register 2 (IREQ2)
001E16	PWM prescaler (PREPWM)	003E16	Interrupt control register 1 (ICON1)
001F16	PWM register (PWM)	003F16	Interrupt control register 2 (ICON2)
		0FFE16	Flash memory control register (FMCR)

 $\boldsymbol{\ast}$ Reserved : Do not write any data to this addresses, because these areas are reserved.

Fig. 9 Memory map of special function register (SFR)

FUNCTIONAL DESCRIPTION

I/O PORTS

The I/O ports have direction registers which determine the input/ output direction of each individual pin. Each bit in a direction register corresponds to one pin, and each pin can be set to be input port or output port.

When "0" is written to the bit corresponding to a pin, that pin becomes an input pin. When "1" is written to that bit, that pin becomes an output pin.

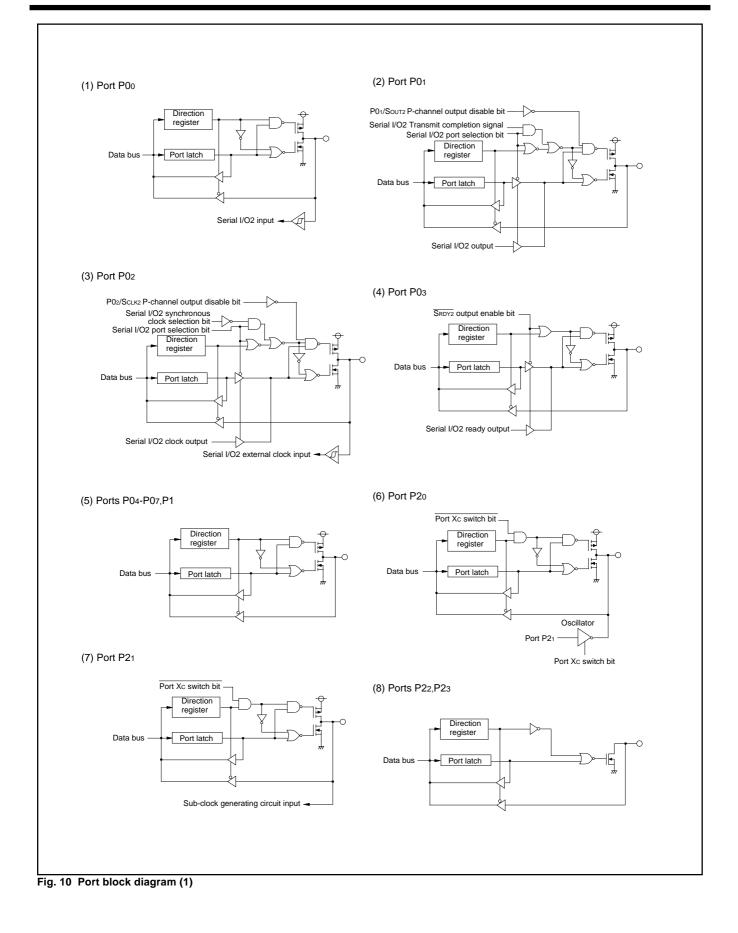
If data is read from a pin which is set to output, the value of the port output latch is read, not the value of the pin itself. Pins set to input are floating. If a pin set to input is written to, only the port output latch is written to and the pin remains floating.

Table 7 I/O port function

Pin	Name	Input/Output	I/O Structure	Non-Port Function	Related SFRs	Ref.No.
P00/SIN2						(1)
P01/SOUT2				Carial 1/02 function 1/0	Carial 1/02 control register	(2)
P02/SCLK2	Port P0			Serial I/O2 function I/O	Serial I/O2 control register	(3)
P03/SRDY2			CMOS compatible input level			(4)
P04–P07		Port P1	CMOS 3-state output			(5)
P10-P17	Port P1					(5)
P20/XCOUT				Sub-clock generating	CPU mode register	(6)
P21/XCIN				circuit		(7)
P22			CMOS compatible			
P23			input level N-channel open-drain			(8)
	Port P2		output			
P24/RxD	101112	Input/output, individual bits		Serial I/O1 function I/O	Serial I/O1 control register	(9)
P25/TxD						(10)
P26/SCLK1						(11)
P27/CNTR0/SRDY1				Serial I/O1 function I/O	Serial I/O1 control register	(12)
				Timer X function I/O	Timer XY mode register	(12)
P30/AN0-	Port P3		A-D conversion input	A-D control register	(13)	
P34/AN4	FULFS		CMOS compatible input level	A-D conversion input		(10)
P40/CNTR1				Timer Y function I/O	Timer XY mode register	(14)
P41/INT0			CMOS 3-state output	External interrupt input	Interrupt edge selection	(15)
P42/INT1					register	(15)
P43/INT2/SCMP2			External interrupt input	Interrupt edge selection		
	Port P4			SCMP2 output	register	(16)
					Serial I/O2 control register	
P44/INT3/PWM				External interrupt input PWM output	Interrupt edge selection register PWM control register	(17)

Note: When bits 5 to 7 of Ports P3 and P4 are read out, the contents are undefined.

HARDWARE FUNCTIONAL DESCRIPTION



FUNCTIONAL DESCRIPTION

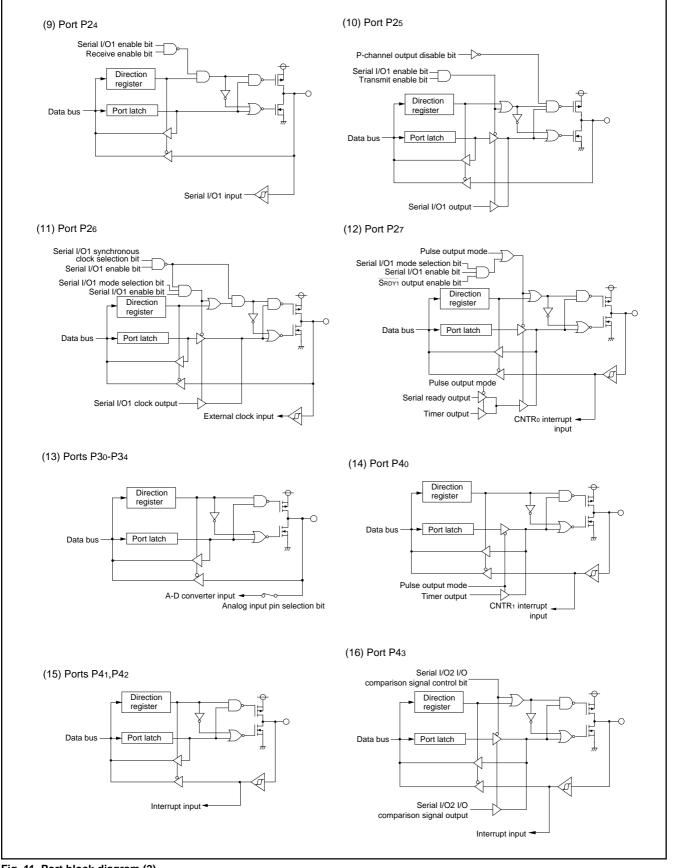


Fig. 11 Port block diagram (2)

FUNCTIONAL DESCRIPTION

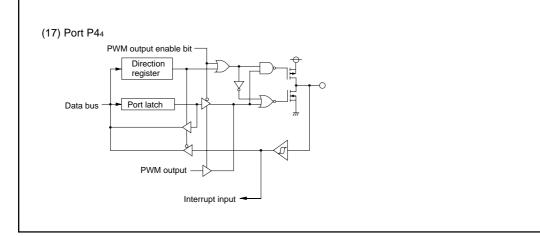


Fig. 12 Port block diagram (3)

FUNCTIONAL DESCRIPTION

INTERRUPTS

Interrupts occur by 15 sources among 15 sources: six external, eight internal, and one software.

Interrupt Control

Each interrupt is controlled by an interrupt request bit, an interrupt enable bit, and the interrupt disable flag except for the software interrupt set by the BRK instruction. An interrupt occurs if the corresponding interrupt request and enable bits are "1" and the interrupt disable flag is "0".

Interrupt enable bits can be set or cleared by software.

Interrupt request bits can be cleared by software, but cannot be set by software.

The BRK instruction cannot be disabled with any flag or bit. The I (interrupt disable) flag disables all interrupts except the BRK instruction interrupt.

When several interrupts occur at the same time, the interrupts are received according to priority.

Interrupt Operation

By acceptance of an interrupt, the following operations are automatically performed:

- 1. The contents of the program counter and the processor status register are automatically pushed onto the stack.
- 2. The interrupt disable flag is set and the corresponding interrupt request bit is cleared.
- 3. The interrupt jump destination address is read from the vector table into the program counter.

■Notes

When setting the followings, the interrupt request bit may be set to "1".

•When setting external interrupt active edge

Related register: Interrupt edge selection register (address 3A16) Timer XY mode register (address 2316)

•When switching interrupt sources of an interrupt vector address where two or more interrupt sources are allocated

Related register: Interrupt edge selection register (address 3A16) When not requiring for the interrupt occurrence synchronized with these setting, take the following sequence.

①Set the corresponding interrupt enable bit to "0" (disabled).

- ②Set the interrupt edge select bit or the interrupt source select bit to "1".
- ③Set the corresponding interrupt request bit to "0" after 1 or more instructions have been executed.
- Set the corresponding interrupt enable bit to "1" (enabled).

Interrupt Source	Priority	Vector Addresses (Note 1)		Interrupt Request	Remarks	
Interrupt Source	FIIOIIty	High	Low	Generating Conditions	Remarks	
Reset (Note 2)	1	FFFD16	FFFC16	At reset	Non-maskable	
ΙΝΤο	2	FFFB16	FFFA16	At detection of either rising or falling edge of INTo input	External interrupt (active edge selectable)	
Reserved	3	FFF916	FFF816	Reserved		
INT1	4	FFF716	FFF616	At detection of either rising or falling edge of INT1 input	External interrupt (active edge selectable)	
INT2	5	FFF516	FFF416	At detection of either rising or falling edge of INT2 input	External interrupt (active edge selectable)	
INT3/ Serial I/O2	6	FFF316	FFF216	At detection of either rising or falling edge of INT3 input/ At completion of serial I/O2 data reception/transmission	External interrupt (active edge selectable) Switch by Serial I/O2/INT3 interrupt source bit	
Reserved	7	FFF116	FFF016	Reserved		
Timer X	8	FFEF16	FFEE16	At timer X underflow		
Timer Y	9	FFED16	FFEC16	At timer Y underflow		
Timer 1	10	FFEB16	FFEA16	At timer 1 underflow	STP release timer underflow	
Timer 2	11	FFE916	FFE816	At timer 2 underflow		
Serial I/O1 reception	12	FFE716	FFE616	At completion of serial I/O1 data reception	Valid when serial I/O1 is selected	
Serial I/O1 transmission	13	FFE516	FFE416	At completion of serial I/O1 transfer shift or when transmis- sion buffer is empty	Valid when serial I/O1 is selected	
CNTR ₀	14	FFE316	FFE216	At detection of either rising or falling edge of CNTR0 input	External interrupt (active edge selectable)	
CNTR1	15	FFE116	FFE016	At detection of either rising or falling edge of CNTR1 input External interrupt (active edge selectable)		
A-D converter	16	FFDF16	FFDE16	At completion of A-D conversion		
BRK instruction	17	FFDD16	FFDC16	At BRK instruction execution	Non-maskable software interrup	

Table 8 Interrupt vector addresses and priority

Notes 1: Vector addresses contain interrupt jump destination addresses.

2: Reset function in the same way as an interrupt with the highest priority.

FUNCTIONAL DESCRIPTION

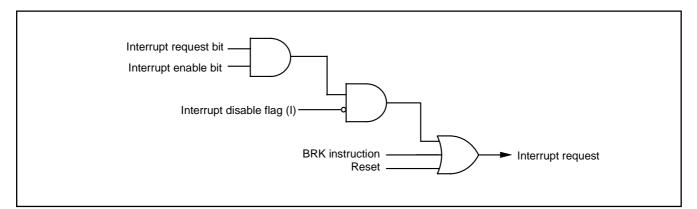
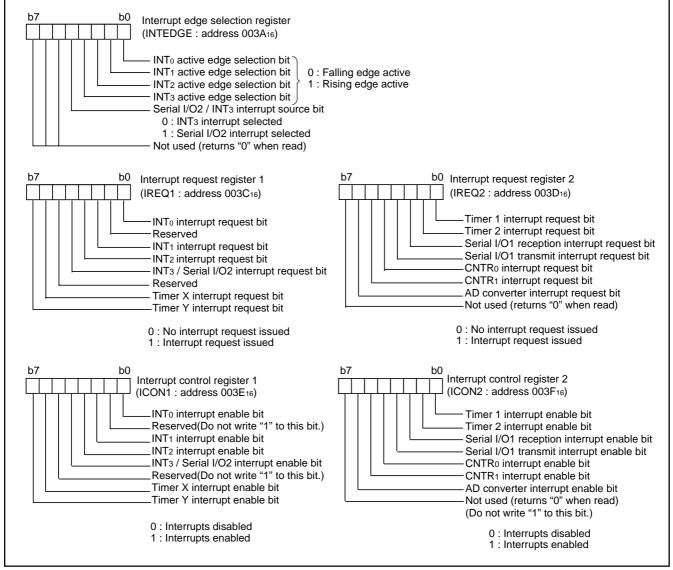


Fig. 13 Interrupt control





TIMERS

The 3850 group (spec. H) has four timers: timer X, timer Y, timer 1, and timer 2.

The division ratio of each timer or prescaler is given by 1/(n + 1), where n is the value in the corresponding timer or prescaler latch. All timers are count down. When the timer reaches "0016", an underflow occurs at the next count pulse and the corresponding timer latch is reloaded into the timer and the count is continued. When a timer underflows, the interrupt request bit corresponding to that timer is set to "1".

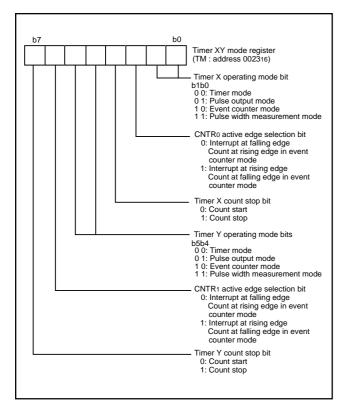


Fig. 15 Structure of timer XY mode register

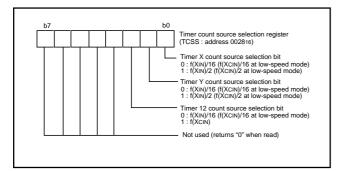


Fig. 16 Structure of timer count source selection register

Timer 1 and Timer 2

The count source of prescaler 12 is the oscillation frequency which is selected by timer 12 count source selection bit. The output of prescaler 12 is counted by timer 1 and timer 2, and a timer underflow sets the interrupt request bit.

Timer X and Timer Y

Timer X and Timer Y can each select in one of four operating modes by setting the timer XY mode register.

(1) Timer Mode

The timer counts the count source selected by Timer count source selection bit.

(2) Pulse Output Mode

The timer counts the count source selected by Timer count source selection bit. Whenever the contents of the timer reach "0016", the signal output from the CNTR0 (or CNTR1) pin is inverted. If the CNTR0 (or CNTR1) active edge selection bit is "0", output begins at " H".

If it is "1", output starts at "L". When using a timer in this mode, set the corresponding port P27 (or port P40) direction register to output mode.

(3) Event Counter Mode

Operation in event counter mode is the same as in timer mode, except that the timer counts signals input through the CNTR0 or CNTR1 pin.

When the CNTR₀ (or CNTR₁) active edge selection bit is "0", the rising edge of the CNTR₀ (or CNTR₁) pin is counted.

When the CNTR0 (or CNTR1) active edge selection bit is "1", the falling edge of the CNTR0 (or CNTR1) pin is counted.

(4) Pulse Width Measurement Mode

If the CNTR0 (or CNTR1) active edge selection bit is "0", the timer counts the selected signals by the count source selection bit while the CNTR0 (or CNTR1) pin is at "H". If the CNTR0 (or CNTR1) active edge selection bit is "1", the timer counts it while the CNTR0 (or CNTR1) pin is at "L".

The count can be stopped by setting "1" to the timer X (or timer Y) count stop bit in any mode. The corresponding interrupt request bit is set each time a timer underflows.

■Note

When switching the count source by the timer 12, X and Y count source bit, the value of timer count is altered in unconsiderable amount owing to generating of a thin pulses in the count input signals.

Therefore, select the timer count source before set the value to the prescaler and the timer.

FUNCTIONAL DESCRIPTION

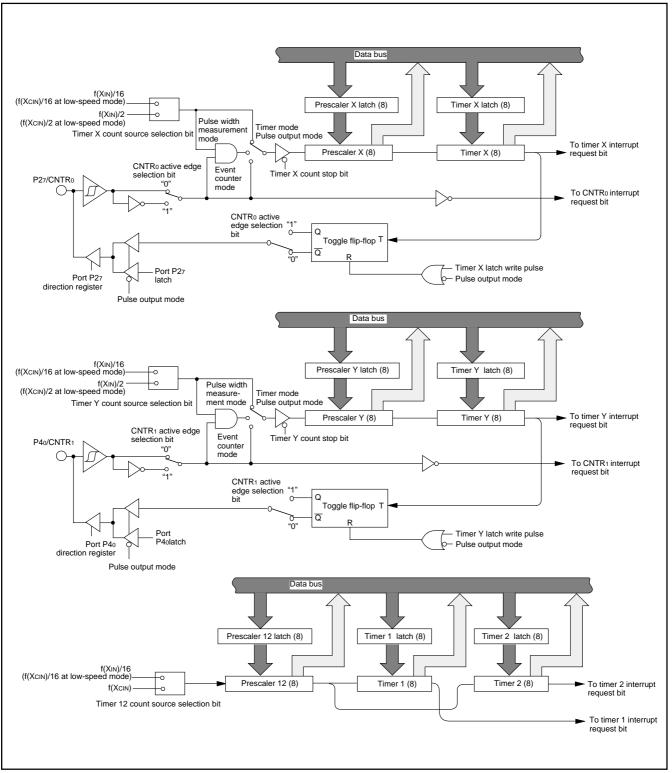


Fig. 17 Block diagram of timer X, timer Y, timer 1, and timer 2

SERIAL I/O •SERIAL I/O1

Serial I/O1 can be used as either clock synchronous or asynchronous (UART) serial I/O. A dedicated timer is also provided for baud rate generation.

(1) Clock Synchronous Serial I/O Mode

Clock synchronous serial I/O mode can be selected by setting the serial I/O1 mode selection bit of the serial I/O1 control register (bit 6 of address 001A16) to "1".

For clock synchronous serial I/O, the transmitter and the receiver must use the same clock. If an internal clock is used, transfer is started by a write signal to the TB/RB.

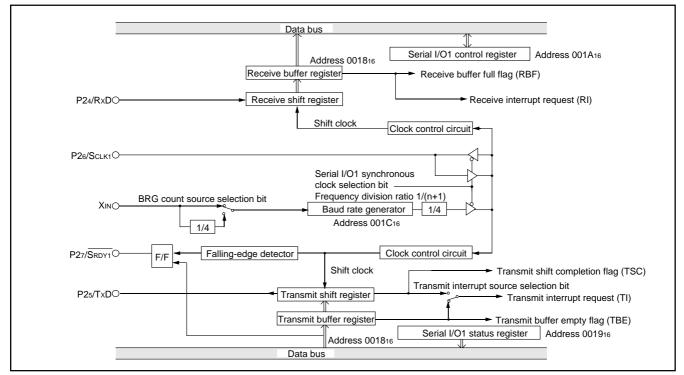


Fig. 18 Block diagram of clock synchronous serial I/O1

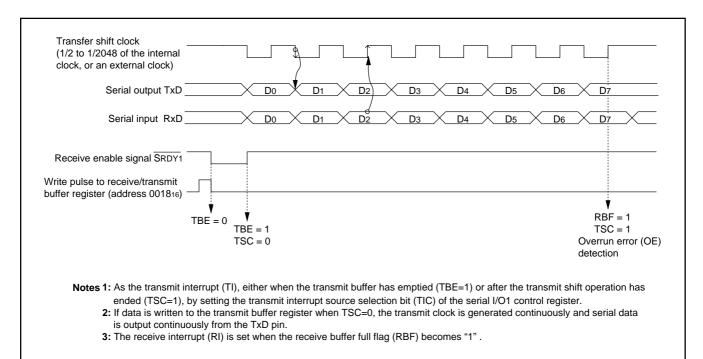


Fig. 19 Operation of clock synchronous serial I/O1 function

(2) Asynchronous Serial I/O (UART) Mode

Clock asynchronous serial I/O mode (UART) can be selected by clearing the serial I/O1 mode selection bit (b6) of the serial I/O1 control register to "0".

Eight serial data transfer formats can be selected, and the transfer formats used by a transmitter and receiver must be identical.

The transmit and receive shift registers each have a buffer, but the

two buffers have the same address in memory. Since the shift register cannot be written to or read from directly, transmit data is written to the transmit buffer register, and receive data is read from the receive buffer register.

The transmit buffer register can also hold the next data to be transmitted, and the receive buffer register can hold a character while the next character is being received.

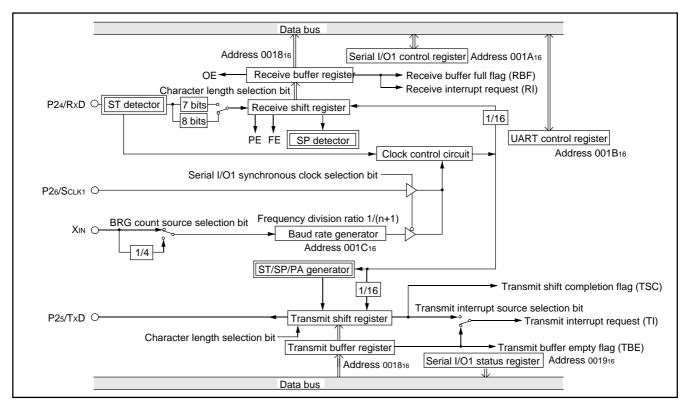


Fig. 20 Block diagram of UART serial I/O1

FUNCTIONAL DESCRIPTION

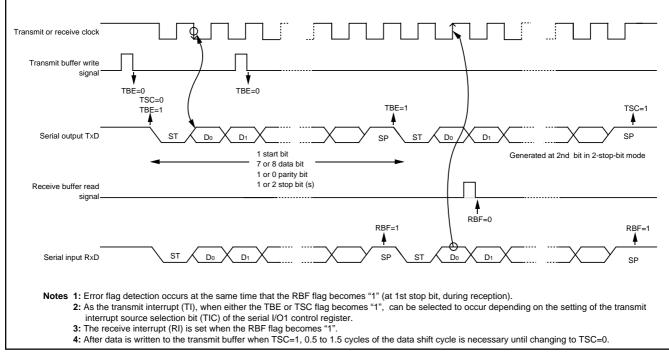


Fig. 21 Operation of UART serial I/O1 function

[Transmit Buffer Register/Receive Buffer Register (TB/RB)] 001816

The transmit buffer register and the receive buffer register are located at the same address. The transmit buffer is write-only and the receive buffer is read-only. If a character bit length is 7 bits, the MSB of data stored in the receive buffer is "0".

[Serial I/O1 Status Register (SIOSTS)] 001916

The read-only serial I/O1 status register consists of seven flags (bits 0 to 6) which indicate the operating status of the serial I/O1 function and various errors.

Three of the flags (bits 4 to 6) are valid only in UART mode.

The receive buffer full flag (bit 1) is cleared to "0" when the receive buffer register is read.

If there is an error, it is detected at the same time that data is transferred from the receive shift register to the receive buffer register, and the receive buffer full flag is set. A write to the serial I/O1 status register clears all the error flags OE, PE, FE, and SE (bit 3 to bit 6, respectively). Writing "0" to the serial I/O1 enable bit SIOE (bit 7 of the serial I/O1 control register) also clears all the status flags, including the error flags.

Bits 0 to 6 of the serial I/O1 status register are initialized to "0" at reset, but if the transmit enable bit (bit 4) of the serial I/O1 control register has been set to "1", the transmit shift completion flag (bit 2) and the transmit buffer empty flag (bit 0) become "1".

[Serial I/O1 Control Register (SIOCON)] 001A16

The serial I/O1 control register consists of eight control bits for the serial I/O1 function.

[UART Control Register (UARTCON)] 001B16

The UART control register consists of four control bits (bits 0 to 3) which are valid when asynchronous serial I/O is selected and set the data format of an data transfer and one bit (bit 4) which is always valid and sets the output structure of the P25/TxD pin.

[Baud Rate Generator (BRG)] 001C16

The baud rate generator determines the baud rate for serial transfer.

The baud rate generator divides the frequency of the count source by 1/(n + 1), where n is the value written to the baud rate generator.

FUNCTIONAL DESCRIPTION

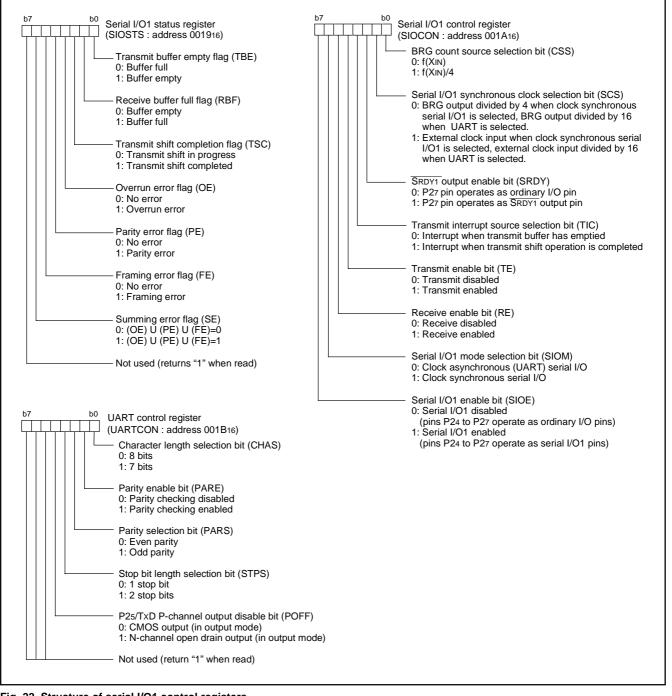


Fig. 22 Structure of serial I/O1 control registers

■Notes on serial I/O

When setting the transmit enable bit of serial I/O1 to "1", the serial I/O1 transmit interrupt request bit is automatically set to "1". When not requiring the interrupt occurrence synchronized with the transmission enalbed, take the following sequence.

①Set the serial I/O1 transmit interrupt enable bit to "0" (disabled).②Set the transmit enable bit to "1".

③Set the serial I/O1 transmit interrupt request bit to "0" after 1 or more instructions have been executed.

④Set the serial I/O1 transmit interrupt enable bit to "1" (enabled).

FUNCTIONAL DESCRIPTION

•SERIAL I/O2

The serial I/O2 can be operated only as the clock synchronous type. As a synchronous clock for serial transfer, either internal clock or external clock can be selected by the serial I/O2 synchronous clock selection bit (b6) of serial I/O2 control register 1.

The internal clock incorporates a dedicated divider and permits selecting 6 types of clock by the internal synchronous clock selection bits (b2, b1, b0) of serial I/O2 control register 1.

Regarding SOUT2 and SCLK2 being output pins, either CMOS output format or N-channel open-drain output format can be selected by the P01/SOUT2, P02/SCLK2 P-channel output disable bit (b7) of serial I/O2 control register 1.

When the internal clock has been selected, a transfer starts by a write signal to the serial I/O2 register (address 001716). After completion of data transfer, the level of the SOUT2 pin goes to high impedance automatically but bit 7 of the serial I/O2 control register 2 is not set to "1" automatically.

When the external clock has been selected, the contents of the serial I/O2 register is continuously sifted while transfer clocks are input. Accordingly, control the clock externally. Note that the SOUT2 pin does not go to high impedance after completion of data transfer.

To cause the SOUT2 pin to go to high impedance in the case where the external clock is selected, set bit 7 of the serial I/O2 control register 2 to "1" when SCLK2 is "H" after completion of data transfer. After the next data transfer is started (the transfer clock falls), bit 7 of the serial I/O2 control register 2 is set to "0" and the SOUT2 pin is put into the active state.

Regardless of the internal clock to external clock, the interrupt request bit is set after the number of bits (1 to 8 bits) selected by the optional transfer bit is transferred. In case of a fractional number of bits less than 8 bits as the last data, the received data to be stored in the serial I/O2 register becomes a fractional number of bits close to MSB if the transfer direction selection bit of serial I/O2 control register 1 is LSB first, or a fractional number of bits close to LSB if the said bit is MSB first. For the remaining bits, the previously received data is shifted.

At transmit operation using the clock synchronous serial I/O, the SCMP2 signal can be output by comparing the state of the transmit pin SOUT2 with the state of the receive pin SIN2 in synchronization with a rise of the transfer clock. If the output level of the SOUT2 pin is equal to the input level to the SIN2 pin, "L" is output from the SCMP2 pin. If not, "H" is output. At this time, an INT2 interrupt request can also be generated. Select a valid edge by bit 2 of the interrupt edge selection register (address 003A16).

[Serial I/O2 Control Registers 1, 2 (SIO2CON1 / SIO2CON2)] 001516, 001616

The serial I/O2 control registers 1 and 2 are containing various selection bits for serial I/O2 control as shown in Figure 23.

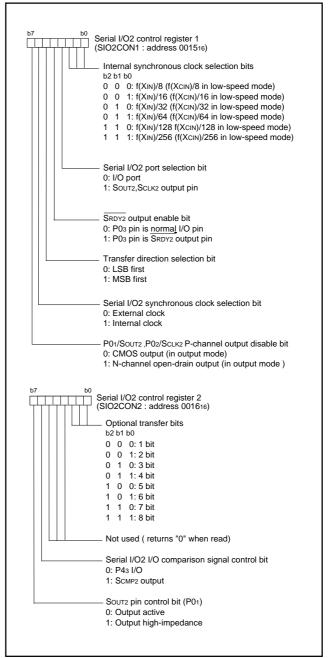


Fig. 23 Structure of Serial I/O2 control registers 1, 2

FUNCTIONAL DESCRIPTION

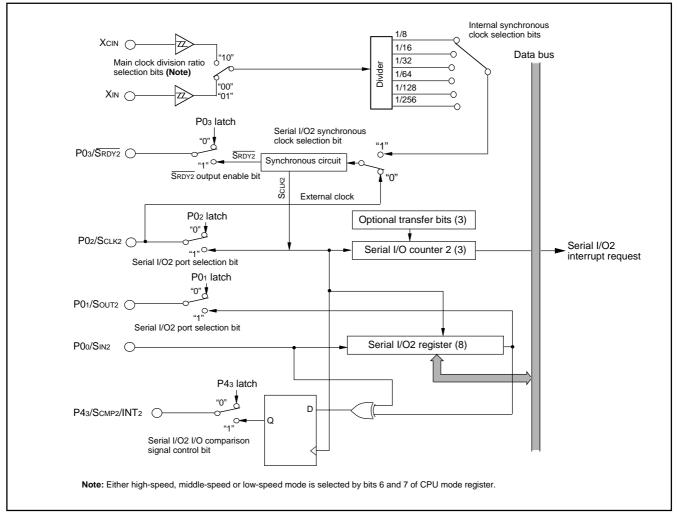


Fig. 24 Block diagram of Serial I/O2

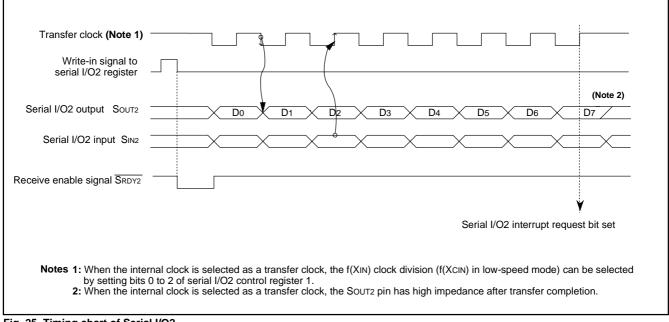
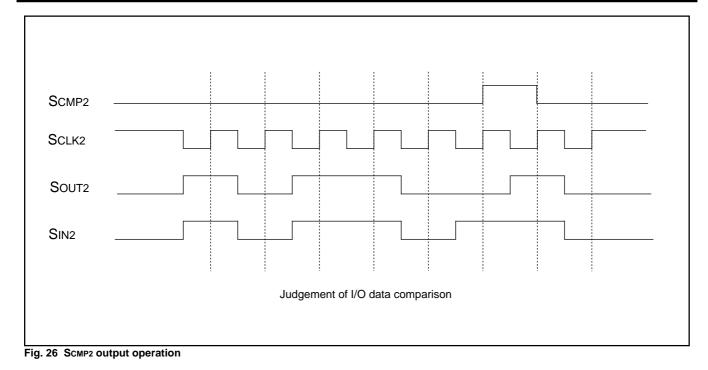


Fig. 25 Timing chart of Serial I/O2

FUNCTIONAL DESCRIPTION



PULSE WIDTH MODULATION (PWM)

The 3850 group (spec. H) has a PWM function with an 8-bit resolution, based on a signal that is the clock input XIN or that clock input divided by 2.

Data Setting

The PWM output pin also functions as port P44. Set the PWM period by the PWM prescaler, and set the "H" term of output pulse by the PWM register.

If the value in the PWM prescaler is n and the value in the PWM register is m (where n = 0 to 255 and m = 0 to 255) :

$$\label{eq:pwm} \begin{split} \text{PWM period} &= 255 \text{ X } (\text{n+1}) \mbox{ / } f(\text{XIN}) \\ &= 31.875 \text{ X } (\text{n+1}) \mbox{ } \mu \text{s} \end{split}$$

 $= 51.075 \times (11+1) \mu s$

(when f(XIN) = 8 MHz,count source selection bit = "0") Output pulse "H" term = PWM period X m / 255

(when f(XIN) = 8 MHz,count source selection bit = "0")

PWM Operation

When bit 0 (PWM enable bit) of the PWM control register is set to "1", operation starts by initializing the PWM output circuit, and pulses are output starting at an "H".

If the PWM register or PWM prescaler is updated during PWM output, the pulses will change in the cycle after the one in which the change was made.

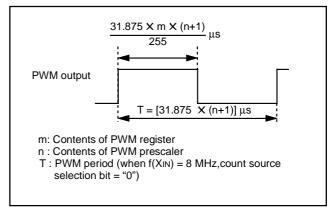


Fig. 27 Timing of PWM period

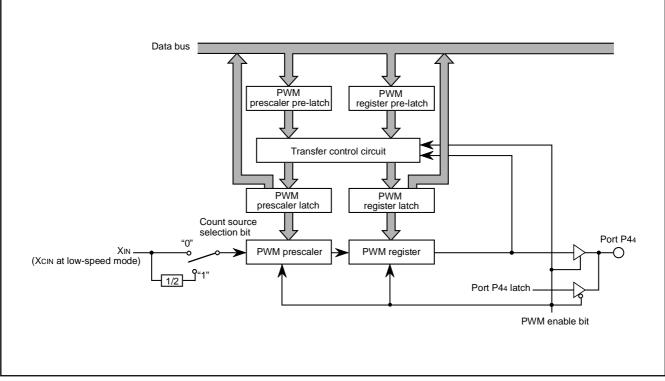


Fig. 28 Block diagram of PWM function

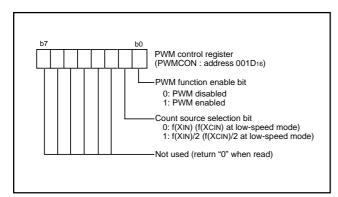
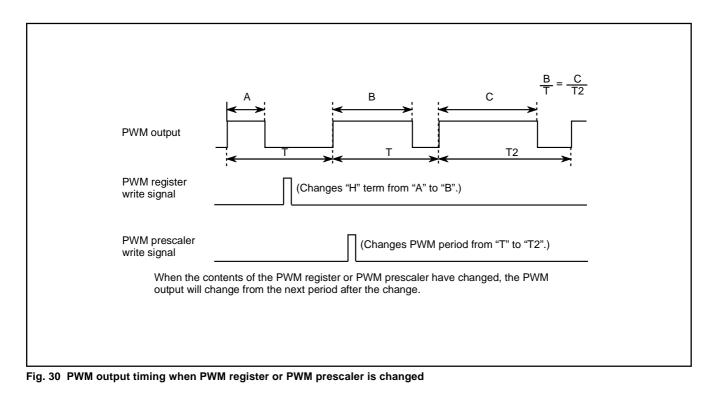


Fig. 29 Structure of PWM control register



■Note

The PWM starts after the PWM function enable bit is set to enable and "L" level is output from the PWM pin. The length of this "L" level output is as follows:

n+1 2 • f(XIN)	sec	(Count source selection bit = 0, where n is the value set in the prescaler)
 f(XIN)	sec	(Count source selection bit = 1, where n is the value set in the prescaler)

A-D CONVERTER [A-D Conversion Registers (ADL, ADH)] 003516, 003616

The A-D conversion registers are read-only registers that store the result of an A-D conversion. Do not read these registers during an A-D conversion.

[AD Control Register (ADCON)] 003416

The AD control register controls the A-D conversion process. Bits 0 to 2 select a specific analog input pin. Bit 4 indicates the completion of an A-D conversion. The value of this bit remains at "0" during an A-D conversion and changes to "1" when an A-D conversion ends. Writing "0" to this bit starts the A-D conversion.

Comparison Voltage Generator

The comparison voltage generator divides the voltage between AVss and VREF into 1024 and outputs the divided voltages.

Channel Selector

The channel selector selects one of ports P30/AN0 to P34/AN4 and inputs the voltage to the comparator.

Comparator and Control Circuit

The comparator and control circuit compare an analog input voltage with the comparison voltage, and the result is stored in the A-D conversion registers. When an A-D conversion is completed, the control circuit sets the A-D conversion completion bit and the A-D interrupt request bit to "1".

Note that because the comparator consists of a capacitor coupling, set $f(X_{\text{IN}})$ to 500 kHz or more during an A-D conversion.

When the A-D converter is operated at low-speed mode, f(XIN) and f(XCIN) do not have the lower limit of frequency, because of the A-D converter has a built-in self-oscillation circuit.

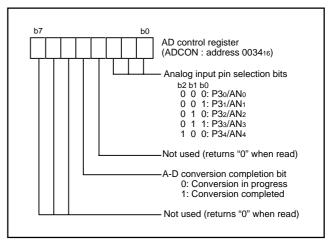
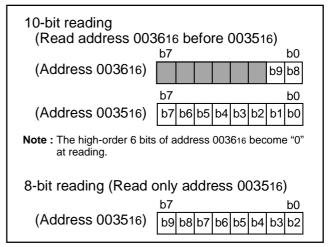


Fig. 31 Structure of AD control register





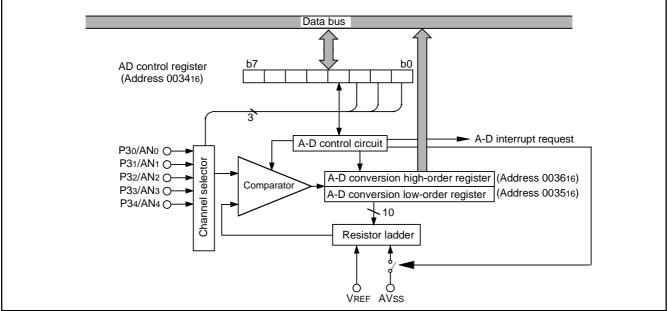


Fig. 33 Block diagram of A-D converter

WATCHDOG TIMER

The watchdog timer gives a mean of returning to the reset status when a program cannot run on a normal loop (for example, because of a software run-away). The watchdog timer consists of an 8-bit watchdog timer L and an 8-bit watchdog timer H.

Standard Operation of Watchdog Timer

When any data is not written into the watchdog timer control register (address 003916) after reset, the watchdog timer is in the stop state. The watchdog timer starts to count down by writing an optional value into the watchdog timer control register (address 003916) and an internal reset occurs at an underflow of the watchdog timer H.

Accordingly, programming is usually performed so that writing to the watchdog timer control register (address 003916) may be started before an underflow. When the watchdog timer control register (address 003916) is read, the values of the high-order 6 bits of the watchdog timer H, STP instruction disable bit, and watchdog timer H count source selection bit are read.

Initial value of watchdog timer

At reset or writing to the watchdog timer control register (address 003916), each watchdog timer H and L is set to "FF16".

•Watchdog timer H count source selection bit operation

Bit 7 of the watchdog timer control register (address 003916) permits selecting a watchdog timer H count source. When this bit is set to "0", the count source becomes the underflow signal of watchdog timer L. The detection time is set to 131.072 ms at f(XIN) = 8 MHz frequency and 32.768 s at f(XCIN) = 32 kHz frequency. When this bit is set to "1", the count source becomes the signal divided by 16 for f(XIN) (or f(XCIN)). The detection time in this case is set to 512 μ s at f(XIN) = 8 MHz frequency and 128 ms at f(XCIN) = 32 kHz frequency. This bit is cleared to "0" after reset.

Operation of STP instruction disable bit

Bit 6 of the watchdog timer control register (address 003916) permits disabling the STP instruction when the watchdog timer is in operation.

When this bit is "0", the STP instruction is enabled.

When this bit is "1", the STP instruction is disabled, once the STP instruction is executed, an internal reset occurs. When this bit is set to "1", it cannot be rewritten to "0" by program. This bit is cleared to "0" after reset.

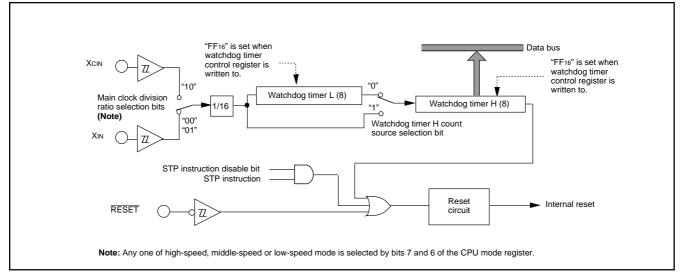
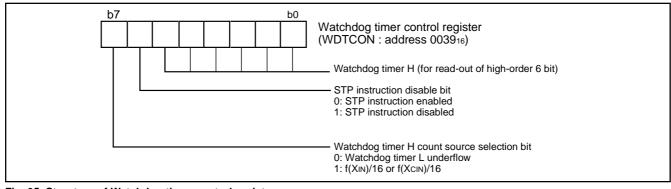
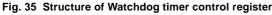


Fig. 34 Block diagram of Watchdog timer





FUNCTIONAL DESCRIPTION

RESET CIRCUIT

To reset the microcomputer, \overrightarrow{RESET} pin must be held at an "L" level for 20 cycles or more of XIN. Then the \overrightarrow{RESET} pin is returned to an "H" level (the power source voltage must be between 2.7 V and 5.5 V, and the oscillation must be stable), reset is released. After the reset is completed, the program starts from the address contained in address FFFD16 (high-order byte) and address FFFC16 (low-order byte). Make sure that the reset input voltage is less than 0.54 V for Vcc of 2.7 V.

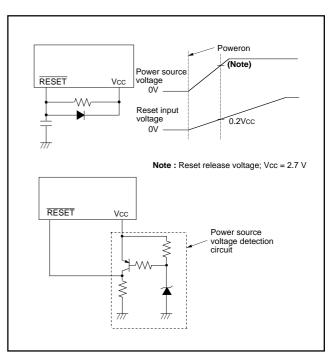


Fig. 36 Reset circuit example

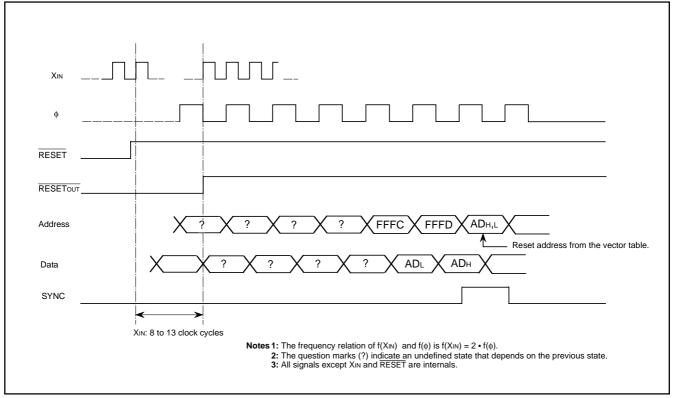


Fig. 37 Reset sequence

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		Address Register contents		Address Register contents
(1) Port P0 (P0)	000016 0016	(34) MISRG	003816 0016
(2) Port P0 d	lirection register (P0D)	000116 0016	(35) Watchdog timer control register (WDTCON)	003916 0 0 1 1 1 1 1 1
(3) Port P1 (P1)	000216 0016	(36) Interrupt edge selection register (INTEDGE)	003A16 0016
(4) Port P1 d	lirection register (P1D)	000316 0016	(37) CPU mode register (CPUM)	003B16 0 1 0 0 1 0 0 0
(5) Port P2 (P2)	000416 0016	(38) Interrupt request register 1 (IREQ1)	003C16 0016
(6) Port P2 d	lirection register (P2D)	000516 0016	(39) Interrupt request register 2 (IREQ2)	003D16 0016
(7) Port P3 (P3)	000616 0016	(40) Interrupt control register 1 (ICON1)	003E16 0016
(8) Port P3 d	lirection register (P3D)	000716 0016	(41) Interrupt control register 2 (ICON2)	003F16 0016
(9) Port P4 (P4)	000816 0016	(42) Processor status register	(PS) X X X X X 1 X X
(10) Port P4 d	lirection register (P4D)	000916 0016	(43) Program counter	(PCH) FFFD16 contents
(11) Serial I/O	2 control register 1 (SIO2CON1)	001516 0016		(PCL) FFFC16 contents
(12) Serial I/O	2 control register 2 (SIO2CON2)	001616 0 0 0 0 1 1 1		
(13) Serial I/O	2 register (SIO2)	001716 X X X X X X X X		
(14) Transmit/	Receive buffer register (TB/RB)	001816 X X X X X X X X		
(15) Serial I/O	1 status register (SIOSTS)	001916 1 0 0 0 0 0 0 0		
(16) Serial I/O	1 control register (SIOCON)	001A16 0016		
(17) UART co	ntrol register (UARTCON)	001B16 1 1 1 0 0 0 0 0		
(18) Baud rate	e generator (BRG)	001C16 X X X X X X X X		
(19) PWM cor	ntrol register (PWMCON)	001D16 0016		
(20) PWM pre	escaler (PREPWM)	001E16 X X X X X X X X		
(21) PWM reg	ister (PWM)	001F16 X X X X X X X X		
(22) Prescaler	r 12 (PRE12)	002016 FF16		
(23) Timer 1 (T1)	002116 0116		
(24) Timer 2 (T2)	002216 0016		
(25) Timer XY	′ mode register (TM)	002316 0016		
(26) Prescaler	r X (PREX)	002416 FF16		
(27) Timer X (TX)	002516 FF16		
(28) Prescaler	r Y (PREY)	002616 FF16]	
(29) Timer Y (TY)	002716 FF16		
(30) Timer cou	int source selection register (TCSS)	002816 0016		
(31) A-D contro	ol register (ADCON)	003416 0 0 0 1 0 0 0 0		
(32) A-D conv	version low-order register (ADL)	003516 X X X X X X X X X		
(33) A-D conv	version high-order register (ADH)	003616 0 0 0 0 0 0 X X		

Note : X : Not fixed

Since the initial values for other than above mentioned registers and RAM contents are indefinite at reset, they must be set.

Fig. 38 Internal status at reset

CLOCK GENERATING CIRCUIT

The 3850 group (spec. H) has two built-in oscillation circuits. An oscillation circuit can be formed by connecting a resonator between XIN and XOUT (XCIN and XCOUT). Use the circuit constants in accordance with the resonator manufacturer's recommended values. No external resistor is needed between XIN and XOUT since a feed-back resistor exists on-chip. However, an external feed-back resistor is needed between XCIN and XCOUT.

Immediately after power on, only the XIN oscillation circuit starts oscillating, and XCIN and XCOUT pins function as I/O ports.

Frequency Control (1) Middle-speed mode

The internal clock ϕ is the frequency of XIN divided by 8. After reset is released, this mode is selected.

(2) High-speed mode

The internal clock ϕ is half the frequency of XIN.

(3) Low-speed mode

The internal clock ϕ is half the frequency of XCIN.

■Note

If you switch the mode between middle/high-speed and low-speed, stabilize both XIN and XCIN oscillations. The sufficient time is required for the sub-clock to stabilize, especially immediately after power on and at returning from the stop mode. When switching the mode between middle/high-speed and low-speed, set the frequency on condition that $f(XIN) > 3 \cdot f(XCIN)$.

(4) Low power dissipation mode

The low power consumption operation can be realized by stopping the main clock XIN in low-speed mode. To stop the main clock, set bit 5 of the CPU mode register to "1". When the main clock XIN is restarted (by setting the main clock stop bit to "0"), set sufficient time for oscillation to stabilize.

The sub-clock XCIN-XCOUT oscillating circuit can not directly input clocks that are generated externally. Accordingly, make sure to cause an external resonator to oscillate.

Oscillation Control (1) Stop mode

If the STP instruction is executed, the internal clock ϕ stops at an "H" level, and XIN and XCIN oscillation stops. When the oscillation stabilizing time set after STP instruction released bit is "0", the prescaler 12 is set to "FF16" and timer 1 is set to "0116". When the oscillation stabilizing time set after STP instruction released bit is "1", set the sufficient time for oscillation of used oscillator to stabilize since nothing is set to the prescaler 12 and timer 1.

Either XIN or XCIN divided by 16 is input to the prescaler 12 as count source. Oscillator restarts when an external interrupt is received, but the internal clock ϕ is not supplied to the CPU (remains at "H") until timer 1 underflows. The internal clock ϕ is supplied for the first time, when timer 1 underflows. This ensures time for the clock oscillation using the ceramic resonators to be stabilized. When the oscillator is restarted by reset, apply "L" level to the RESET pin until the oscillation is stable since a wait time will not be generated.

(2) Wait mode

If the WIT instruction is executed, the internal clock ϕ stops at an "H" level, but the oscillator does not stop. The internal clock ϕ restarts at reset or when an interrupt is received. Since the oscillator does not stop, normal operation can be started immediately after the clock is restarted.

To ensure that the interrupts will be received to release the STP or WIT state, their interrupt enable bits must be set to "1" before executing of the STP or WIT instruction.

When releasing the STP state, the prescaler 12 and timer 1 will start counting the clock XIN divided by 16. Accordingly, set the timer 1 interrupt enable bit to "0" before executing the STP instruction.

■Note

When using the oscillation stabilizing time set after STP instruction released bit set to "1", evaluate time to stabilize oscillation of the used oscillator and set the value to the timer 1 and prescaler 12.

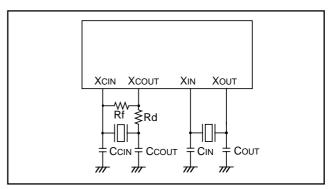


Fig. 39 Ceramic resonator circuit

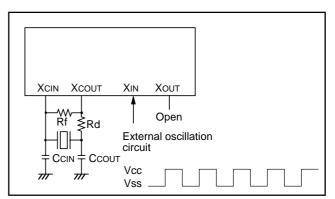


Fig. 40 External clock input circuit

FUNCTIONAL DESCRIPTION

[MISRG (MISRG)] 003816

MISRG consists of three control bits (bits 1 to 3) for middle-speed mode automatic switch and one control bit (bit 0) for oscillation stabilizing time set after STP instruction released.

By setting the middle-speed mode automatic switch start bit to "1" while operating in the low-speed mode and setting the middlespeed mode automatic switch set bit to "1", XIN oscillation automatically starts and the mode is automatically switched to the middle-speed mode.

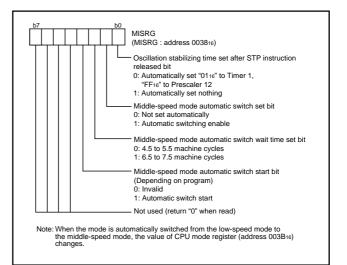
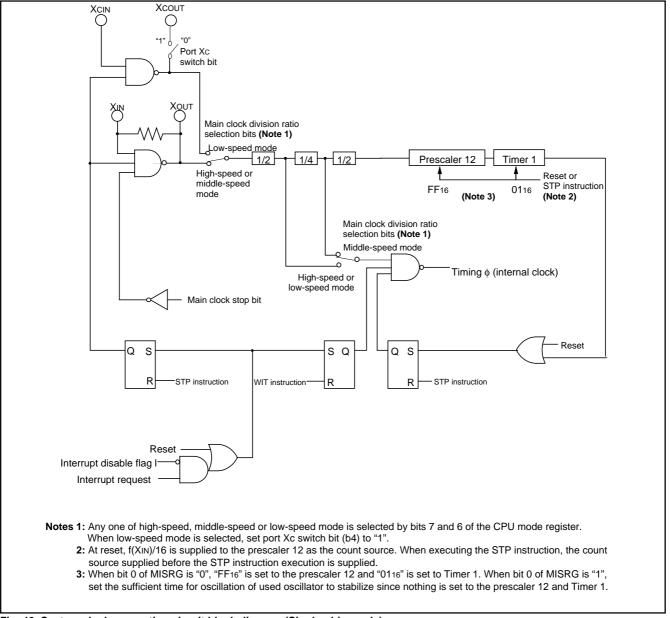
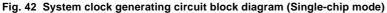
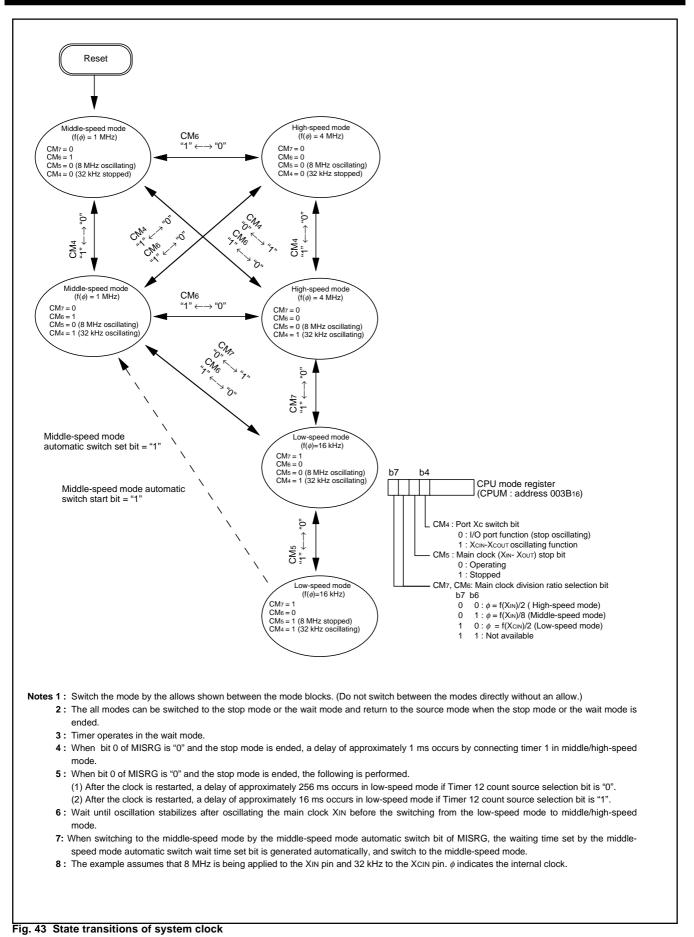


Fig. 41 Structure of MISRG





FUNCTIONAL DESCRIPTION



FLASH MEMORY VERSION Summary

Table 9 shows the summary of the M38507F8 (flash memory version).

Table 9 Summary of M38507F8 (flash memory version)

Item		Specification		
Power source voltage		Vcc = 2.7–5.5 V (Note 1) Vcc = 2.7–3.6 V (Note 2)		
Program/Erase VPP voltage		4.5–5.5 V, f(XIN) = 8 MHz		
Flash memory mode		3 modes (Parallel I/O mode, Standard serial I/O mode, CPU rewrite mode)		
Erase block division	User ROM area	1 block (32 Kbytes)		
	Boot ROM area	1 block (4 Kbytes) (Note 3)		
Program method		Byte program		
Erase method		Batch erasing		
Program/Erase control	method	Program/Erase control by software command		
Number of commands		6 commands		
Number of program/Erase times		100 times		
ROM code protection		Available in parallel I/O mode, and standard serial I/O mode		

Notes 1: The power source voltage must be Vcc = 4.5-5.5 V at program and erase operation.

2: The power source voltage can be Vcc = 3.0-3.6 V also at program and erase operation.

3: The Boot ROM area has had a standard serial I/O mode control program stored in it when shipped from the factory. This Boot ROM area can be rewritten in only parallel I/O mode.

Flash Memory Mode

The M38507F8 (flash memory version) has an internal new DINOR (DIvided bit line NOR) flash memory that can be rewritten with a single power source when Vcc is 5 V, and 2 power sources when Vcc is 3.3-5.0 V.

For this flash memory , three flash memory modes are available in which to read, program, and erase: parallel I/O and standard serial I/O modes in which the flash memory can be manipulated using a programmer and a CPU rewrite mode in which the flash memory can be manipulated by the Central Processing Unit (CPU). Each mode is detailed in the pages to follow.

The flash memory of the M38507F8 is divided into User ROM area and Boot ROM area as shown in Figure 44.

In addition to the ordinary user ROM area to store a microcomputer operation control program, the flash memory has a Boot ROM area that is used to store a program to control rewriting in CPU rewrite and standard serial I/O modes. This Boot ROM area has had a standard serial I/O mode control program stored in it when shipped from the factory. However, the user can write a rewrite control program in this area that suits the user's application system. This Boot ROM area can be rewritten in only parallel I/O mode.

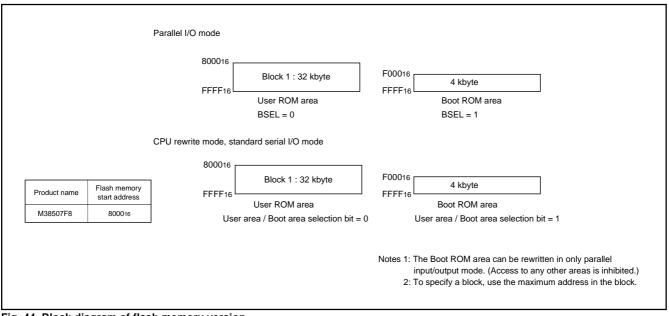


Fig. 44 Block diagram of flash memory version

CPU Rewrite Mode

In CPU rewrite mode, the on-chip flash memory can be operated on (read, program, or erase) under control of the Central Processing Unit (CPU).

In CPU rewrite mode, only the user ROM area shown in Figure 44 can be rewritten; the Boot ROM area cannot be rewritten. Make sure the program and block erase commands are issued for only the user ROM area and each block area.

The control program for CPU rewrite mode can be stored in either user ROM or Boot ROM area. In the CPU rewrite mode, because the flash memory cannot be read from the CPU, the rewrite control program must be transferred to internal RAM area before it can be executed.

Microcomputer Mode and Boot Mode

The control program for CPU rewrite mode must be written into the user ROM or Boot ROM area in parallel I/O mode beforehand. (If the control program is written into the Boot ROM area, the standard serial I/O mode becomes unusable.)

See Figure 44 for details about the Boot ROM area.

Normal microcomputer mode is entered when the microcomputer is reset with pulling CNVss pin low. In this case, the CPU starts operating using the control program in the user ROM area.

When the microcomputer is reset by pulling the P41/INT0 pin high, the CNVss pin high, the CPU starts operating using the control program in the Boot ROM area (program start address is FFFC16, FFFD16 fixation). This mode is called the "boot" mode.

Block Address

Block addresses refer to the maximum address of each block. These addresses are used in the block erase command. In case of the M38507F8, it has only one block.

Outline Performance (CPU Rewrite Mode)

In the CPU rewrite mode, the CPU erases, programs and reads the internal flash memory as instructed by software commands. This rewrite control program must be transferred to internal RAM before it can be executed.

The CPU rewrite mode is accessed by applying $5V \pm 10\%$ to the CNVss pin and writing "1" for the CPU rewrite mode select bit (bit 1 in address 0FFE16). Software commands are accepted once the mode is accessed.

Use software commands to control program and erase operations. Whether a program or erase operation has terminated normally or in error can be verified by reading the status register.

Figure 45 shows the flash memory control register.

Bit 0 is the RY/BY status flag used exclusively to read the operating status of the flash memory. During programming and erase operations, it is "0". Otherwise, it is "1".

Bit 1 is the CPU rewrite mode select bit. When this bit is set to "1" and $5V \pm 10\%$ are applied to the CNVss pin, the M38507F8 accesses the CPU rewrite mode. Software commands are accepted once the mode is accessed. In CPU rewrite mode, the CPU becomes unable to access the internal flash memory directly. Therefore, use the control program in RAM for write to bit 1. To set this bit to "1", it is necessary to write "0" and then write "1" in succession. The bit can be set to "0" by only writing a "0".

Bit 2 is the CPU rewrite mode entry flag. This bit can be read to check whether the CPU rewrite mode has been entered or not.

Bit 3 is the flash memory reset bit used to reset the control circuit of the internal flash memory. This bit is used when exiting CPU rewrite mode and when flash memory access has failed. When the CPU rewrite mode select bit is "1", writing "1" for this bit resets the control circuit. To release the reset, it is necessary to set this bit to "0".

Bit 4 is the User area/Boot area selection bit. When this bit is set to "1", Boot ROM area is accessed, and CPU rewrite mode in Boot ROM area is available. In boot mode, this bit is set "1" automatically. Operation of this bit must be in RAM area.

Figure 46 shows a flowchart for setting/releasing the CPU rewrite mode.

Precautions on CPU Rewrite Mode

Described below are the precautions to be observed when rewriting the flash memory in CPU rewrite mode.

(1) Operation speed

During CPU rewrite mode, set the internal clock frequency 4MHz or less using the main clock division ratio selection bits (bit 6, 7 at 003B16).

(2) Instructions inhibited against use

The instructions which refer to the internal data of the flash memory cannot be used during CPU rewrite mode .

(3) Interrupts inhibited against use

The interrupts cannot be used during CPU rewrite mode because they refer to the internal data of the flash memory.

(4) Watchdog timer

In case of the watchdog timer has been running already, the internal reset generated by watchdog timer underflow does not happen, because of watchdog timer is always clearing during program or erase operation.

(5) Reset

Reset is always valid. In case of CNVss = H when reset is released, boot mode is active. So the program starts from the address contained in address FFFC16 and FFFD16 in boot ROM area.

FUNCTIONAL DESCRIPTION

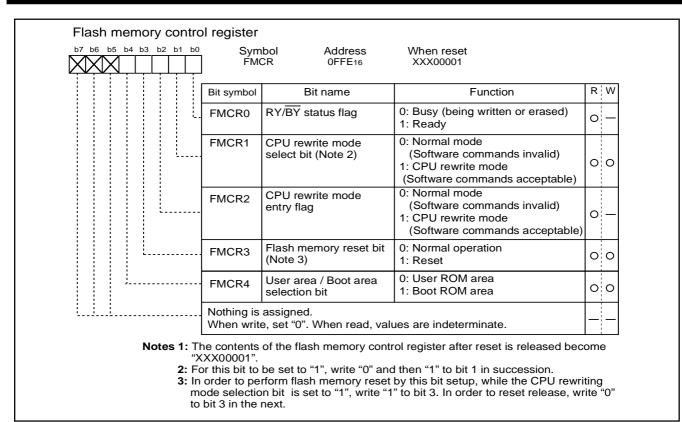


Fig. 45 Flash memory control registers

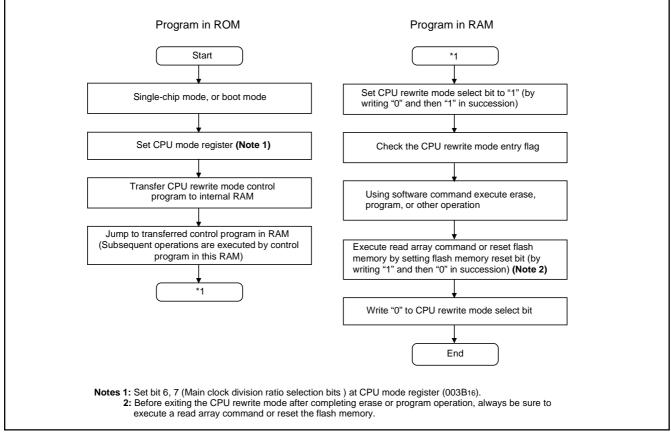


Fig. 46 CPU rewrite mode set/reset flowchart

Software Commands

Table 10 lists the software commands. After setting the CPU rewrite mode select bit to "1", write a software command to specify an erase or program operation.

The content of each software command is explained below.

Read Array Command (FF16)

The read array mode is entered by writing the command code "FF16" in the first bus cycle. When an address to be read is input in one of the bus cycles that follow, the content of the specified address is read out at the data bus (D0-D7).

The read array mode is retained intact until another command is written. And after power on and after recover from deep power down mode, this mode is selected also.

Read Status Register Command (7016)

When the command code "7016" is written in the first bus cycle, the content of the status register is read out at the data bus (D0-D7) by a read in the second bus cycle.

The status register is explained in the next section.

Clear Status Register Command (5016)

This command is used to clear the bits SR1,SR4 and SR5 of the status register after they have been set. These bits indicate that operation has ended in an error. To use this command, write the command code "5016" in the first bus cycle.

Table 10 List of software commands (CPU rewrite mode)

		F	First bus cycle			Second bus cycle		
Command	Cycle number	Mode	Address	Data (Do to D7)	Mode	Address	Data (Do to D7)	
Read array	1	Write	X (Note 4)	FF16				
Read status register	2	Write	Х	7016	Read	Х	SRD (Note 1)	
Clear status register	1	Write	Х	5016				
Program	2	Write	Х	4016	Write	WA (Note 2)	WD (Note 2)	
Erase all block	2	Write	Х	2016	Write	Х	2016	
Block erase	2	Write	Х	2016	Write	BA (Note 3)	D016	

Notes 1: SRD = Status Register Data

2: WA = Write Address, WD = Write Data

3: BA = Block Address (Enter the maximum address of each block.)

4: X denotes a given address in the user ROM area .

Program Command (4016)

Program operation starts when the command code "4016" is written in the first bus cycle. Then, if the address and data to program are written in the 2nd bus cycle, program operation (data programming and verification) will start.

Whether the write operation is completed can be confirmed by reading the status register or the RY/BY status flag. When the program starts, the read status register mode is accessed automatically and the content of the status register is read into the data bus (D0–D7). The status register bit 7 (SR7) is set to "0" at the same time the write operation starts and is returned to "1" upon completion of the write operation. In this case, the read status register mode remains active until the read array command (FF16) is written.

The RY/BY status flag is "0" during write operation and "1" when the write operation is completed as is the status register bit 7.

At program end, program results can be checked by reading the status register.

Erase All Blocks Command (2016/2016)

By writing the command code "2016" in the first bus cycle and the confirmation command code "2016" in the second bus cycle that follows, the system starts erase all blocks(erase and erase verify).

Whether the erase all blocks command is terminated can be confirmed by reading the status register or the RY/BY status flag. When the erase all blocks operation starts, the read status register mode is accessed automatically and the content of the status register can be read out. The status register bit 7 (SR7) is set to "0" at the same time the erase operation starts and is returned to "1" upon completion of the erase operation. In this case, the read status register mode remains active until the read array command (FF16) is written.

The RY/ \overline{BY} status flag is "0" during erase operation and "1" when the erase operation is completed as is the status register bit 7.

At erase all blocks end, erase results can be checked by reading the status register. For details, refer to the section where the status register is detailed.

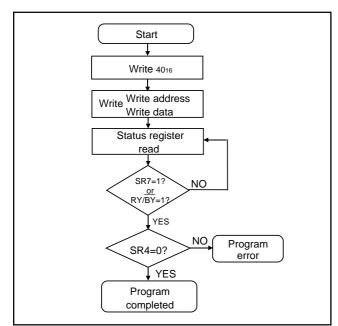


Fig. 47 Program flowchart

Block Erase Command (2016/D016)

By writing the command code "2016" in the first bus cycle and the confirmation command code "D016" in the second bus cycle that follows to the block address of a flash memory block, the system initiates a block erase (erase and erase verify) operation.

Whether the block erase operation is completed can be confirmed by reading the status register or the RY/BY status flag. At the same time the block erase operation starts, the read status register mode is automatically entered, so the content of the status register can be read out. The status register bit 7 (SR7) is set to "0" at the same time the block erase operation starts and is returned to "1" upon completion of the block erase operation. In this case, the read status register mode remains active until the read array command (FF16) is written.

The RY/BY status flag is "0" during block erase operation and "1" when the block erase operation is completed as is the status register bit 7.

After the block erase operation is completed, the status register can be read out to know the result of the block erase operation. For de-

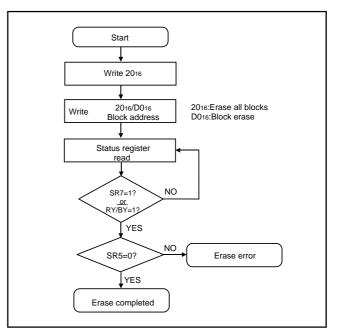


Fig. 48 Erase flowchart

FUNCTIONAL DESCRIPTION

Status Register

The status register shows the operating state of the flash memory and whether erase operations and programs ended successfully or in error. It can be read in the following ways.

- (1) By reading an arbitrary address from the user ROM area after writing the read status register command (7016)
- (2) By reading an arbitrary address from the user ROM area in the period from when the program starts or erase operation starts to when the read array command (FF16) is input

Table 11 shows the status register.

- Also, the status register can be cleared in the following way.
- (1) By writing the clear status register command (5016)
- (2) In the deep power down mode
- (3) In the power supply off state
- After a reset, the status register is set to "8016".
- Each bit in this register is explained below.

Sequencer status (SR7)

After power-on, and after recover from deep power down mode, the sequencer status is set to "1"(ready).

The sequencer status indicates the operating status of the device. This status bit is set to "0" (busy) during write or erase operation and is set to "1" upon completion of these operations.

Erase status (SR5)

The erase status informs the operating status of erase operation to the CPU. When an erase error occurs, it is set to "1". The erase status is reset to "0" when cleared.

Program status (SR4)

The program status informs the operating status of write operation to the CPU. When a write error occurs, it is set to "1". The program status is reset to "0" when cleared.

If "1" is written for any of the SR5 or SR4 bits, the program, erase all blocks, and block erase commands are not accepted. Before executing these commands, execute the clear status register command (5016) and clear the status register.

Also, any commands are not correct, both SR5 and SR4 are set to "1".

Table 11 Definition of each bit in status register

Each bit of	Status name	Definition		
SRD0 bits	Status hame	"1"	"O"	
SR7 (bit7)	Sequencer status	Ready	Busy	
SR6 (bit6)	Reserved	-	-	
SR5 (bit5)	Erase status	Terminated in error	Terminated normally	
SR4 (bit4)	Program status	Terminated in error	Terminated normally	
SR3 (bit3)	Reserved	-	-	
SR2 (bit2)	Reserved	-	-	
SR1 (bit1)	Reserved	-	-	
SR0 (bit0)	Reserved	-	-	

FUNCTIONAL DESCRIPTION

Full Status Check

By performing full status check, it is possible to know the execution results of erase and program operations. Figure 49 shows a full sta-

tus check flowchart and the action to be taken when each error occurs.

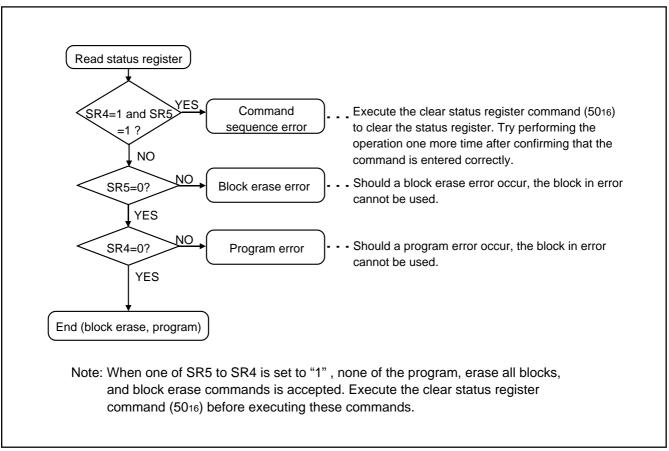


Fig. 49 Full status check flowchart and remedial procedure for errors

Functions To Inhibit Rewriting Flash Memory Version

To prevent the contents of the flash memory version from being read out or rewritten easily, the device incorporates a ROM code protect function for use in parallel I/O mode and an ID code check function for use in standard serial I/O mode.

ROM code protect function

The ROM code protect function is the function inhibit reading out or modifying the contents of the flash memory version by using the ROM code protect control address (FFDB16) during parallel I/O mode. Figure 50 shows the ROM code protect control address (FFDB16). (This address exists in the user ROM area.)

If one of the pair of ROM code protect bits is set to "0", ROM code protect is turned on, so that the contents of the flash memory version are protected against readout and modification. ROM code protect is implemented in two levels. If level 2 is selected, the flash memory is protected even against readout by a shipment inspection LSI tester, etc. When an attempt is made to select both level 1 and level 2, level 2 is selected by default.

If both of the two ROM code protect reset bits are set to "00", ROM code protect is turned off, so that the contents of the flash memory version can be read out or modified. Once ROM code protect is turned on, the contents of the ROM code protect reset bits cannot be modified in parallel I/O mode. Use the serial I/O or some other mode to rewrite the contents of the ROM code protect reset bits.

b7 b6 b5 b4 b3 b2 b1 b0	Symbol ROMCP		Vhen reset FF16
	Bit symbol	Bit name	Function
	Reserved	bit	Always set this bit to "1"
	ROMCP2	ROM code protect level 2 set bit (Note 1, 2)	b3 b2 00: Protect enabled 01: Protect enabled 10: Protect enabled 11: Protect disabled
	ROMCR	ROM code protect reset bit (Note 3)	 b5 b4 00: Protect removed 01: Protect set bit effective 10: Protect set bit effective 11: Protect set bit effective
	ROMCP1	ROM code protect level 1 set bit (Note 1)	b7 b6 00: Protect enabled 01: Protect enabled 10: Protect enabled 11: Protect disabled
readout o 2: When RO inspectior 3: The ROM ROM cod	r modification M code pro- n LSI tester, I code prote e protect le	on in parallel input/output m tect level 2 is turned on, R , etc. also is inhibited. to reset bits can be used to vel 2. However, since these	hip flash memory is protected against hode. OM code readout by a shipment o turn off ROM code protect level 1 and be bits cannot be changed in parallel serial input/output mode or some other

Fig. 50 ROM code protect control address

FUNCTIONAL DESCRIPTION

ID Code Check Function

Use this function in standard serial I/O mode. When the contents of the flash memory are not blank, the ID code sent from the peripheral unit is compared with the ID code written in the flash memory to see if they match. If the ID codes do not match, the commands sent from the peripheral unit are not accepted. The ID code consists of 8-bit data, the areas of which are FFD416 to FFDA16. Write a program which has had the ID code preset at these addresses to the flash memory.

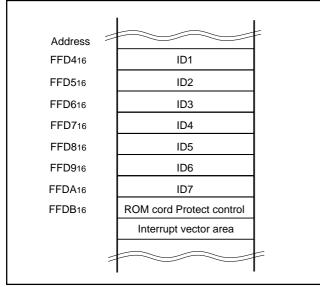


Fig. 51 ID code store addresses

Parallel I/O Mode

The parallel I/O mode is entered by making connections shown in Figure 52 and then turning the Vcc power supply on.

Address

The user ROM is only one block as shown in Figure 44. The block address referred to in this data sheet is the maximum address value of each block.

User ROM and Boot ROM Areas

In parallel I/O mode, the user ROM and boot ROM areas shown in Figure 44 can be rewritten. The BSEL pin is used to choose between these two areas. The user ROM area is selected by pulling the BSEL input low; the boot ROM area is selected by driving the BSEL input high. Both areas of flash memory can be operated on in the same way.

Program and block erase operations can be performed in the user ROM area. The user ROM area and its block is shown in Figure 44.

The user ROM area is 32 Kbytes in size. In parallel I/O mode, it is located at addresses 800016 through FFFF16. The boot ROM area is 4 Kbytes in size. In parallel I/O mode, it is located at addresses F00016 through FFFF16. Make sure program and block erase operations are always performed within this address range. (Access to any location outside this address range is prohibited.)

In the Boot ROM area, an erase block operation is applied to only one 4 Kbyte block. The boot ROM area has had a standard serial I/O mode control program stored in it when shipped from the Renesas factory. Therefore, using the device in standard serial input/output mode, you do not need to write to the boot ROM area.

Functional Outline (Parallel I/O Mode)

In parallel I/O mode, bus operation modes—Read, Output Disable, Standby, Write, and Deep Power Down—are selected by the status of the \overline{CE} , \overline{OE} , \overline{WE} , and \overline{RP} input pins.

The contents of erase, program, and other operations are selected by writing a software command. The data, status register, etc. in memory can only be read out by a read after software command input.

Program and erase operations are controlled using software commands.

The following explains about bus operation modes, software commands, and status register.

Bus Operation Modes Read

The Read mode is entered by pulling the \overline{OE} pin low when the \overline{CE} pin is low and the \overline{WE} and \overline{RP} pins are high. There are two read modes: array, and status register, which are selected by software command input. In read mode, the data corresponding to each software command entered is output from the data I/O pins D0–D7. The read array mode is automatically selected when the device is powered on or after it exits deep power down mode.

Output Disable

The output disable mode is entered by pulling the \overline{CE} pin low and the \overline{WE} , \overline{OE} , and \overline{RP} pins high. Also, the data I/O pins are placed in the high-impedance state.

Standby

The standby mode is entered by driving the \overline{CE} pin high when the \overline{RP} pin is high. Also, the data I/O pins are placed in the high-impedance state. However, if the \overline{CE} pin is set high during erase or program operation, the internal control circuit does not halt immediately and normal power consumption is required until the operation under way is completed.

Write

The write mode is entered by pulling the \overline{WE} pin low when the \overline{CE} pin is low and the \overline{OE} and \overline{RP} pins are high. In this mode, the device accepts the software commands or write data entered from the data I/O pins. A program, erase, or some other operation is initiated depending on the content of the software command entered here. The input data such as address and software command is latched at the rising edge of \overline{WE} or \overline{CE} whichever occurs earlier.

Deep Power Down

The deep power down is entered by pulling the \overline{RP} pin low. Also, the data I/O pins are placed in the high-impedance state. When the device is freed from deep power down mode, the read array mode is selected and the content of the status register is set to "8016". If the \overline{RP} pin is pulled low during erase or program operation, the operation under way is canceled and the data in the relevant block becomes invalid.

Table 12	Relationship between	control signals and bus operation modes
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Mode	Pin name	CE	ŌE	WE	RP	Do to D7
Read	Array	VIL	VIL	Vih	Vін	Data output
Read	Status register	VIL	VIL	Vih	Vін	Status register data output
Output disabled		VIL	Vін	Vін	Vін	High impedance
Stand by	Stand by		Х	Х	Vін	High impedance
	Program	VIL	Vін	VIL	Vін	Command/data input
Write	Erase all blocks	VIL	Vін	VIL	Vін	Command input
	Block erase	VIL	Vін	VIL	Vін	Command input
Deep power down		Х	Х	Х	VIL	High impedance

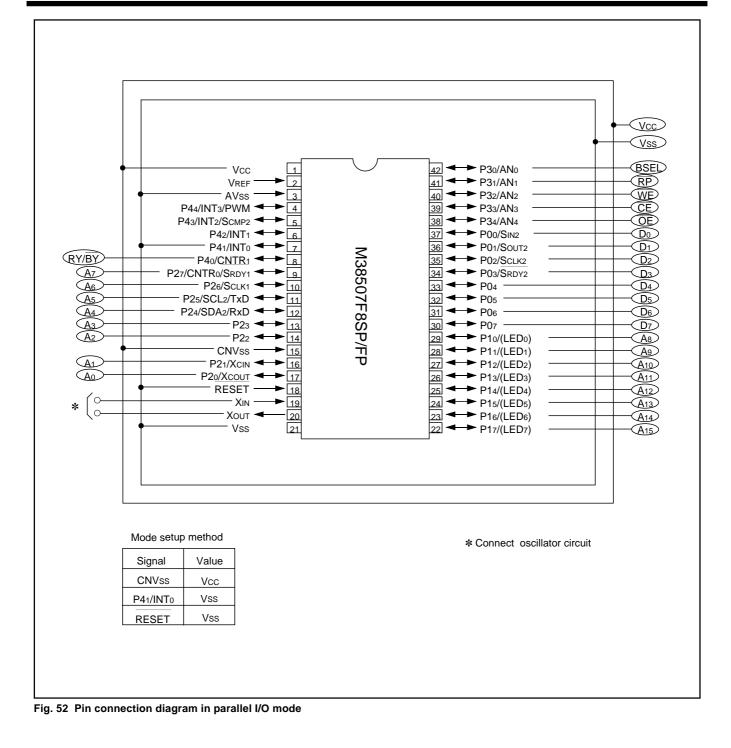
Note: X can be VIL or VIH.

FUNCTIONAL DESCRIPTION

Table 13 Description of Pin Function (Flash Memory Parallel I/O Mode)

Pin name	Signal name	I/O	Function	
Vcc,Vss	Power supply input		Apply 5.0 \pm 0.5 V to the Vcc pin and 0 V to the Vss pin.	
CNVss	CNVss	Ι	Connect this pin to Vcc.	
RESET	Reset input	I	Reset input pin. When reset is held low, more than 20 cycles of clock are required at the XIN pin.	
Xin	Clock input	Ι	Connect a ceramic or crystal resonator between the XIN and XOUT pins. When entering an externally drived clock, enter it from XIN and leave XOUT	
Хоит	Clock output	0	open.	
AVss	Analog power supply input		Connect AVss to Vss.	
Vref	Reference voltage input	Ι	Input AD reference voltage or keep open.	
P00 to P07	Data I/O Do to D7	I/O	These are data D0–D7 input/output pins.	
P10 to P17	Address input A8 to A15	Ι	These are address A8–A15 input pins.	
P20 to P27	Address input Ao to A7	I	These are address A0–A7 input pins.	
P30	BSEL input	I	This is a BSEL input pin.	
P31	RP input	I	This is a \overline{RP} input pin.	
P32	WE input	I	This is a WE input pin.	
P33	CE input	I	This is a \overline{CE} input pin.	
P34	OE input	I	This is a \overline{OE} input pin.	
P40	RY/BY output	0	This is a RY/BY output pin.	
P41	Input P41	I	Enter low signals to this pin.	
P42 to P44	Input P4	 I	Input "H" or "L" or keep open.	

FUNCTIONAL DESCRIPTION



Software Commands

Table 14 lists the software commands. By entering a software command from the data I/O pins (D0–D7) in Write mode, specify the content of the operation, such as erase or program operation, to be performed.

The following explains the content of each software command.

Read Array Command (FF16)

The read array mode is entered by writing the command code "FF16" in the first bus cycle. When an address to be read is input in one of the bus cycles that follow, the content of the specified address is output from the data I/O pins (D0–D7).

The read array mode is retained intact until another command is written.

The read array mode is also selected automatically when the device is powered on and after it exits deep power down mode.

Read Status Register Command (7016)

When the command code "7016" is written in the first bus cycle, the content of the status register is output from the data I/O pins (Do–D7) by a read in the second bus cycle. Since the content of the status register is updated at the falling edge of \overline{OE} or \overline{CE} , the \overline{OE} or \overline{CE} signal must be asserted each time the status is read. The status register is explained in the next section.

Clear Status Register Command (5016)

This command is used to clear the bits SR4,SR5 of the status register after they have been set. These bits indicate that operation has ended in an error. To use this command, write the command code "5016" in the first bus cycle.

Table 14 Software command	list (parallel I/O mode)
---------------------------	--------------------------

Command	Cycle number	First bus cycle			Second bus cycle		
		Mode	Address	Data (Do to D7)	Mode	Address	Data (Do to D7)
Read array	1	Write	X(Note 4)	FF16			
Read status register	2	Write	Х	7016	Read	Х	SRD(Note 1)
Clear status register	1	Write	Х	5016			
Program	2	Write	Х	4016	Write	WA(Note 2)	WD(Note 2)
All block erase	2	Write	Х	2016	Write	Х	2016
Block erase	2	Write	Х	2016	Write	BA(Note 3)	D016

Notes 1: SRD = Status Register Data

2: WA = Write Address, WD = Write Data

3: BA = Block Address (Enter the maximum address of each block)

4: X denotes a given address in the user ROM area or boot ROM area.

Program Command (4016)

The program operation starts when the command code "4016" is written in the first bus cycle. Then, if the address and data to program are written in the 2nd bus cycle, program operation (data programming and verification) will start.

Whether the write operation is completed can be confirmed by reading the status register or the RY/BY signal status. When the program starts, the read status register mode is accessed automatically and the content of the status register can be read out from the data bus (D0–D7). The status register bit 7 (SR7) is set to "0" at the same time the write operation starts and is returned to "1" upon completion of the write operation. In this case, the read status register mode remains active until the read array command (FF16) is written.

The RY/BY pin is "L" during write operation and "H" when the write operation is completed as is the status register bit 7.

At program end, program results can be checked by reading the status register.

Erase All Blocks Command (2016/2016)

By writing the command code "2016" in the first bus cycle and the confirmation command code "2016" in the second bus cycle that follows, the system starts erase all blocks (erase and erase verify). Whether the erase all blocks command is terminated can be confirmed by reading the status register or the RY/BY signal status . When the erase all blocks operation starts, the read status register mode is accessed automatically and the content of the status register can be read out. The status register bit 7 (SR7) is set to "0" at the same time the erase operation starts and is returned to "1" upon completion of the erase operation. In this case, the read status register mode remains active until the read array command (FF16) is written.

The RY/BY pin is "L" during erase operation and "H" when the erase operation is completed as is the status register bit 7.

At erase all blocks end, erase results can be checked by reading the status register. For details, refer to the section where the status register is detailed.

Block Erase Command (2016/D016)

By writing the command code "2016" in the first bus cycle and the confirmation command code "D016" in the second bus cycle that follows to the block address of a flash memory block, the system initiates a block erase (erase and erase verify) operation.

Whether the block erase operation is completed can be confirmed by reading the status register or the RY/BY signal. At the same time the block erase operation starts, the read status register mode is automatically entered, so the content of the status register can be read out. The status register bit 7 (SR7) is set to "0" at the same time the block erase operation starts and is returned to "1" upon completion of the block erase operation. In this case, the read status register mode remains active until the read array command (FF16) is written.

The RY/BY pin is "L" during block erase operation and "H" when the block erase operation is completed as is the status register bit 7. After the block erase operation is completed, the status register can be read out to know the result of the block erase operation. For details, refer to the section where the status register is detailed.

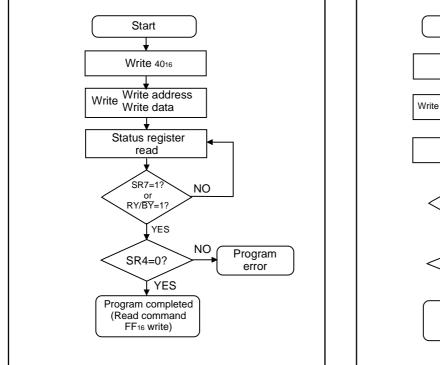


Fig. 53 Page program flowchart

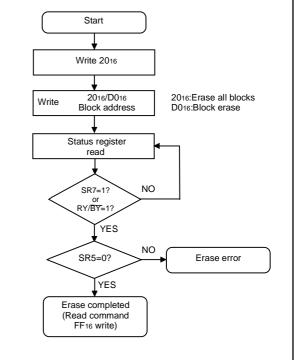


Fig. 54 Block erase flowchart

FUNCTIONAL DESCRIPTION

Status Register

The status register indicates status such as whether an erase operation or a program ended successfully or in error. It can be read under the following conditions.

- In the read array mode when the read status register command (7016) is written and the block address is subsequently read.
- (2) In the period from when the program write or auto erase starts to when the read array command (FF16)

The status register is cleared in the following situations.

- (1) By writing the clear status register command (5016)
- (2) In the deep power down mode
- (3) In the power supply off state

Table 15 gives the definition of each status register bit. When power is turned on or returning from the deep power down mode, the status register outputs "8016".

Sequencer status (SR7)

The sequencer status indicates the operating status of the flash memory. When power is turned on or returning from the deep power down mode, "1" is set for it. This bit is "0" (busy) during the write or erase operations and becomes "1" when these operations ends.

Erase Status (SR5)

The erase status reports the operating status of the erase operation. If an erase error occurs, it is set to "1". When the erase status is cleared, it is set to "0".

Program Status (SR4)

The program status reports the operating status of the write operation. If a write error occurs, it is set to "1". When the program status is cleared, it is set to "0".

If "1" is written for any of the SR5, SR4 bits, the program erase all blocks, block erase, commands are not accepted. Before executing these commands, execute the clear status register command (5016) and clear the status register.

Also, any commands are not correct, both SR5 and SR4 are set to "1".

Full Status Check

Results from executed erase and program operations can be known by running a full status check. Figure 55 shows a flowchart of the full status check and explains how to remedy errors which occur.

Ready/Busy (RY/BY) pin

The RY/ \overline{BY} pin is an output pin (N-chanel open drain output) which, like the sequencer status (SR7), indicates the operating status of the flash memory. It is "L" level during auto program or auto erase operations and becomes to the high impedance state (ready state) when these operations end. The RY/ \overline{BY} pin requires an external pull-up.

Table 15 Status register

Each bit of	Status name	Definition			
SRD0 bits	Status name	"1"	"0"		
SR7 (D7)	Sequencer status	Ready	Busy		
SR6 (D6)	Reserved	-	-		
SR5 (D5)	Erase status	Ended in error	Ended successfully		
SR4 (D4)	Program status	Ended in error	Ended successfully		
SR3 (D3)	Reserved	-	-		
SR2 (D2)	Reserved	-	-		
SR1 (D1)	Reserved	-	-		
SR0 (D0)	Reserved	-	-		

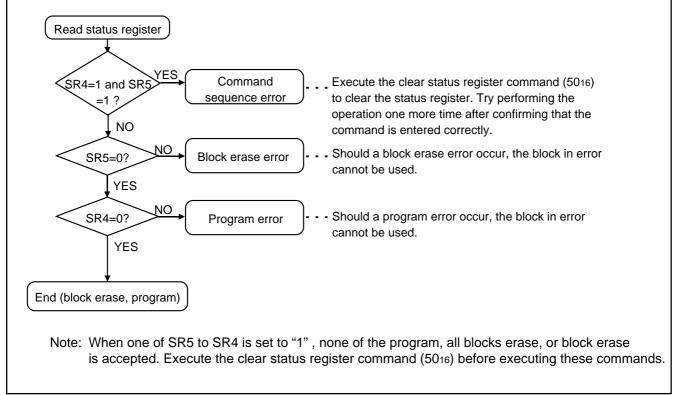


Fig. 55 Full status check flowchart and remedial procedure for errors

FUNCTIONAL DESCRIPTION

Standard serial I/O mode

The standard serial I/O mode inputs and outputs the software commands, addresses and data needed to operate (read, program, erase, etc.) the internal flash memory. This I/O is clock synchronized serial. This modes require a purpose-specific peripheral unit.

The standard serial I/O mode is different from the parallel I/O mode in that the CPU controls flash memory rewrite (uses the CPU's rewrite mode), rewrite data input and so forth. The standard serial I/O mode is started by connecting "H" to the P26 (SCLK1) pin and the P41(INT0) pin and "H" to the CNVss pin (when VCc = 4.5 V to 5.5 V, connect to Vcc; when Vcc = 2.7 V to 4.5 V, supply 4.5 V to 5.5 V to Vpp from an external source), and releasing the reset operation. (In the ordinary command mode, set CNVss pin to "L" level.)

This control program is written in the boot ROM area when the product is shipped from Renesas Technology Corp. Accordingly, make note of the fact that the standard serial I/O mode cannot be used if the boot ROM area is rewritten in the parallel I/O mode. Figure 56 shows the pin connections for the standard serial I/O mode. Serial data I/O uses SI/O1 data serially in 8-bit units.

To use standard serial I/O mode. The operation uses the four SI/O1 pins SCLK1, RxD, TxD and $\overline{\text{SRDY1}}$ (BUSY). The SCLK1 pin is the transfer clock input pin through which an external transfer clock is input. The TxD pin is for CMOS output. The $\overline{\text{SRDY1}}$ (BUSY) pin outputs an "L" level when ready for reception and an "H" level when reception starts.

In the standard serial I/O mode, only the user ROM area indicated in Figure 44 can be rewritten. The boot ROM cannot.

In the standard serial I/O mode, a 7-byte ID code is used. When there is data in the flash memory, commands sent from the peripheral unit (programmer) are not accepted unless the ID code matches.

Overview of standard serial I/O mode

In standard serial I/O mode, software commands, addresses and data are input and output between the MCU and peripheral units (serial programer, etc.) using 4-wire clock-synchronized serial I/O (SI/O1).

In reception, software commands, addresses and program data are synchronized with the rise of the transfer clock that is input to the SCLK1 pin, and are then input to the MCU via the RxD pin. In transmission, the read data and status are synchronized with the fall of the transfer clock, and output from the TxD pin.

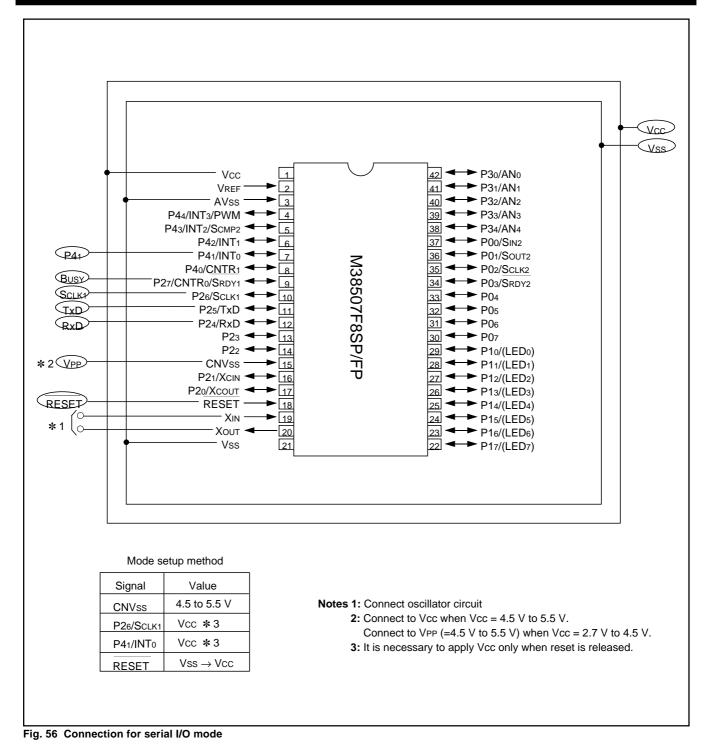
The TxD pin is for CMOS output. Transfer is in 8-bit units with LSB first.

When busy, such as during transmission, reception, erasing or program execution, the $\overline{SRDY1}$ (BUSY) pin is "H" level. Accordingly, always start the next transfer after the $\overline{SRDY1}$ (BUSY) pin is "L" level. Also, data and status registers in memory can be read after inputting software commands. Status, such as the operating state of the flash memory or whether a program or erase operation ended successfully or not, can be checked by reading the status register. Here following are explained software commands, status registers, etc.

FUNCTIONAL DESCRIPTION

Pin	Name	I/O	Description
Vcc,Vss	Power input		Apply program/erase protection voltage to Vcc pin and 0 V to Vss pin.
CNVss	CNVss	I	Connect to Vcc when Vcc = 4.5 V to 5.5 V. Connect to Vpp (= 4.5 V to 5.5 V) when Vcc = 2.7 V to 4.5 V.
RESET	Reset input	I	Reset input pin. While reset is "L" level, a 20 cycle or longer clock must be input to XIN pin.
XIN	Clock input	I	Connect a ceramic resonator or crystal oscillator between XIN and XOUT pins. To input an externally generated clock, input it to XIN pin
Χουτ	Clock output	0	and open Xout pin.
AVss	Analog power supply input		Connect AVss to Vss .
Vref	Reference voltage input	I	Enter the reference voltage for AD from this pin.
P00 to P07	Input port P0	I	Input "H" or "L" level signal or open.
P10 to P17	Input port P1	I	Input "H" or "L" level signal or open.
P20 to P23	Input port P2	I	Input "H" or "L" level signal or open.
P24	RxD input	I	Serial data input pin
P25	TxD output	0	Serial data output pin
P26	SCLK1 input	I	Serial clock input pin
P27	BUSY output	0	BUSY signal output pin
P30 to P34	Input port P3	I	Input "H" or "L" level signal or open.
P40, P42 to P44	Input port P4	I	Input "H" or "L" level signal or open.
P41	Input P41	I	Input "H" level signal, when reset is released.

Table 16 Pin functions (Flash memory standard serial I/O mode)



Software Commands

Table 17 lists software commands. In the standard serial I/O mode, erase operations, programs and reading are controlled by transferring software commands via the RxD pin. Software commands are

explained here below. Basically, the software commands of the standard serial I/O mode is as same as that of the parallel I/O mode, but it is excluded 1 command of block erase, and it is added 3 command of ID check, download function, version data output function.

Table 17	Software commands	(Standard serial	I/O mode 1)

	Control command	1st byte transfer	2nd byte	3rd byte	4th byte	5th byte	6th byte		When ID is not verified
1	Page read	FF ₁₆	Address (middle)	Address (high)	Data output	Data output	Data output	Data output to 259th byte	Not acceptable
2	Page program	41 ₁₆	Address (middle)	Address (high)	Data input	Data input	Data input	Data input to 259th byte	Not acceptable
3	Erase all blocks	A7 ₁₆	D016						Not acceptable
4	Read status register	70 ₁₆	SRD output	SRD1 output					Acceptable
5	Clear status register	5016							Not acceptable
6	ID check	F516	Address (low)	Address (middle)	Address (high)	ID size	ID1	To ID7	Acceptable
7	Download function	FA ₁₆	Size (low)	Size (high)	Check- sum	Data input	To required number of times		Not acceptable
8	Version data output function	FB ₁₆	Version data output	Version data output	Version data output	Version data output	Version data output	Version data output to 9th byte	Acceptable

Notes 1: Shading indicates transfer from flash memory microcomputer to peripheral unit. All other data is transferred from the peripheral unit to the flash memory microcomputer.

2: SRD refers to status register data. SRD1 refers to status register 1 data.

3: All commands can be accepted when the flash memory is totally blank.

4: Address high (A16 to A23) must be "0016".

FUNCTIONAL DESCRIPTION

Page Read Command

This command reads the specified page (256 bytes) in the flash memory sequentially one byte at a time. Execute the page read command as explained here following.

(1) Transfer the "FF16" command code with the 1st byte.

- (2) Transfer addresses A8 to A15 and A16 to A23 ("0016") with the 2nd and 3rd bytes respectively.
- (3) From the 4th byte onward, data (D0–D7) for the page (256 bytes) specified with addresses A8 to A23 will be output sequentially from the smallest address first in sync with the fall of the clock.

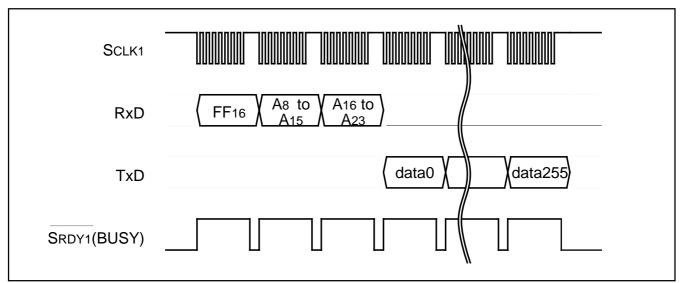


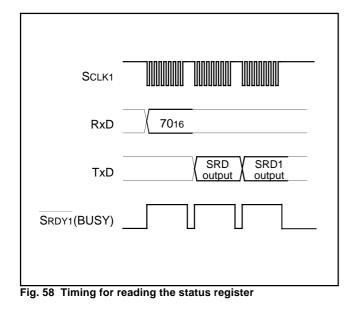
Fig. 57 Timing for page read

Read Status Register Command

This command reads status information. When the "7016" command code is sent with the 1st byte, the contents of the status register (SRD) specified with the 2nd byte and the contents of status register 1 (SRD1) specified with the 3rd byte are read.

Clear Status Register Command

This command clears the bits (SR4–SR5) which are set when the status register operation ends in error. When the "5016" command code is sent with the 1st byte, the aforementioned bits are cleared. When the clear status register operation ends, the $\overline{\text{SRDY1}}$ (BUSY) signal changes from the "H" to the "L" level.



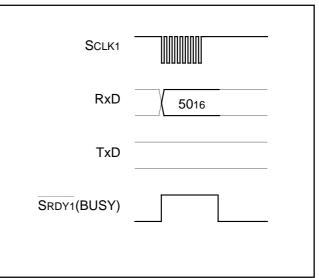


Fig. 59 Timing for clearing the status register

Page Program Command

This command writes the specified page (256 bytes) in the flash memory sequentially one byte at a time. Execute the page program command as explained here following.

- (1) Transfer the "4116" command code with the 1st byte.
- (2) Transfer addresses A8 to A15 and A16 to A23 ("0016") with the 2nd and 3rd bytes respectively.
- (3) From the 4th byte onward, as write data (D0–D7) for the page (256 bytes) specified with addresses A8 to A23 is input sequentially from the smallest address first, that page is automatically written.

When reception setup for the next 256 bytes ends, the $\overline{SRDY1}$ (BUSY) signal changes from the "H" to the "L" level. The result of the page program can be known by reading the status register. For more information, see the section on the status register.

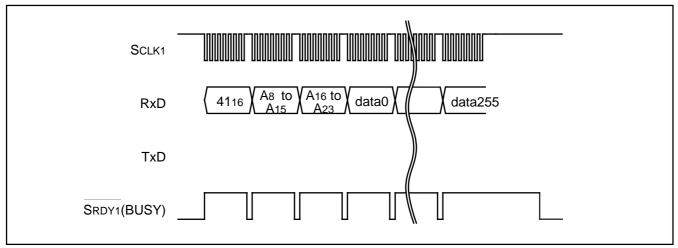


Fig. 60 Timing for the page program

Erase All Blocks Command

This command erases the content of all blocks. Execute the erase all blocks command as explained here following.

(1) Transfer the "A716" command code with the 1st byte.

(2) Transfer the verify command code "D016" with the 2nd byte. With the verify command code, the erase operation will start and continue for all blocks in the flash memory.

When block erasing ends, the $\overline{SRDY1}$ (BUSY) signal changes from the "H" to the "L" level . The result of the erase operation can be known by reading the status register.

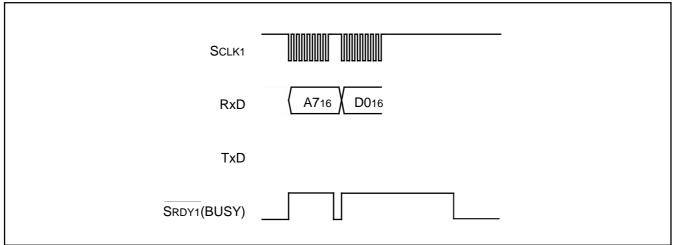


Fig. 61 Timing for erasing all blocks

FUNCTIONAL DESCRIPTION

Download Command

This command downloads a program to the RAM for execution. Execute the download command as explained here following.

- (1) Transfer the "FA16" command code with the 1st byte.
- (2) Transfer the program size with the 2nd and 3rd bytes.
- (3) Transfer the check sum with the 4th byte. The check sum is added to all data sent with the 5th byte onward.

(4) The program to execute is sent with the 5th byte onward.

When all data has been transmitted, if the check sum matches, the downloaded program is executed. The size of the program will vary according to the internal RAM.

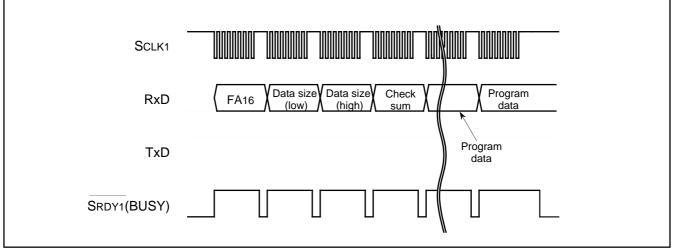


Fig. 62 Timing for download

Version Information Output Command

This command outputs the version information of the control program stored in the boot area. Execute the version information output command as explained here following.

- (1) Transfer the "FB16" command code with the 1st byte.
- (2) The version information will be output from the 2nd byte onward. This data is composed of 8 ASCII code characters.

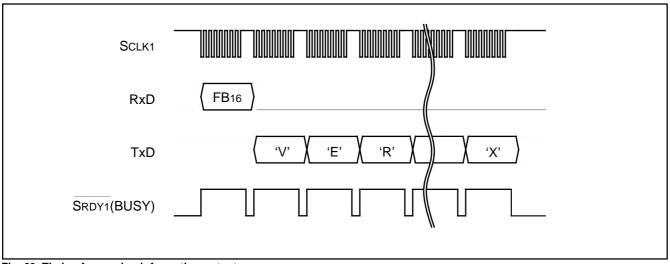


Fig. 63 Timing for version information output

ID Check

This command checks the ID code. Execute the boot ID check command as explained here following.

- (1) Transfer the "F516" command code with the 1st byte.
- (2) Transfer addresses A0 to A7, A8 to A15 and A16 to A23 ("0016") of

the 1st byte of the ID code with the 2nd, 3rd and 4th bytes respectively.

- (3) Transfer the number of data sets of the ID code with the 5th byte.
- (4) The ID code is sent with the 6th byte onward, starting with the 1st byte of the code.

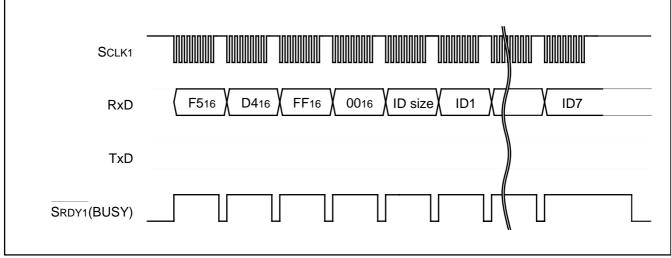


Fig. 64 Timing for the ID check

ID Code

When the flash memory is not blank, the ID code sent from the peripheral units and the ID code written in the flash memory are compared to see if they match. If the codes do not match, the command sent from the peripheral units is not accepted. An ID code contains 8 bits of data. Area is, from the 1st byte, addresses FFD416 to FFDA16. Write a program into the flash memory, which already has the ID code set for these addresses.

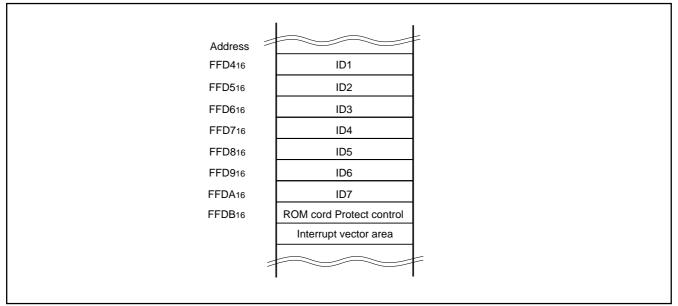


Fig. 65 ID code storage addresses

Status Register (SRD)

The status register indicates operating status of the flash memory and status such as whether an erase operation or a program ended successfully or in error. It can be read by writing the read status register command (7016). Also, the status register is cleared by writing the clear status register command (5016).

Table 18 gives the definition of each status register bit. After clearing the reset, the status register outputs "8016".

Sequencer status (SR7)

Table 18 Status register (SRD)

After power-on and recover from deep power down mode, the sequencer status is set to "1"(ready). The sequencer status indicates the operating status of the device. This status bit is set to "0" (busy) during write or erase operation and is set to "1" upon completion of these operations.

Erase Status (SR5)

The erase status reports the operating status of the auto erase operation. If an erase error occurs, it is set to "1". When the erase status is cleared, it is set to "0".

Program Status (SR4)

The program status reports the operating status of the auto write operation. If a write error occurs, it is set to "1". When the program status is cleared, it is set to "0".

		Defi	nition
SRD0 bits	Status name	"1"	"0"
SR7 (bit7)	Sequencer status	Ready	Busy
SR6 (bit6)	Reserved	-	-
SR5 (bit5)	Erase status	Terminated in error	Terminated normally
SR4 (bit4)	Program status	Terminated in error	Terminated normally
SR3 (bit3)	Reserved	-	-
SR2 (bit2)	Reserved	-	-
SR1 (bit1)	Reserved	-	-
SR0 (bit0)	Reserved	-	-

Status Register 1 (SRD1)

Status register 1 indicates the status of serial communications, results from ID checks and results from check sum comparisons. It can be read after the SRD by writing the read status register command (7016). Also, status register 1 is cleared by writing the clear status register command (5016).

Table 19 gives the definition of each status register bit. "0016" is output when power is turned on and the flag status is maintained even after the reset.

Boot Update Completed Bit (SR15)

This flag indicates whether the control program was downloaded to the RAM or not, using the download function.

Check Sum Consistency Bit (SR12)

This flag indicates whether the check sum matches or not when a program is downloaded for execution using the download function.

ID Check Completed Bits (SR11 and SR10)

These flags indicate the result of ID checks. Some commands cannot be accepted without an ID check.

Data Reception Time Out (SR9)

This flag indicates when a time out error is generated during data reception. If this flag is attached during data reception, the received data is discarded and the microcomputer returns to the command wait state.

	Status nome	Definition		
SKUTDIIS	SRD1 bits Status name -		"O"	
SR15 (bit7)	Boot update completed bit	Update completed	Not Update	
SR14 (bit6)	Reserved	-	-	
SR13 (bit5)	Reserved	-	-	
SR12 (bit4)	Checksum match bit	Match	Mismatch	
SR11 (bit3)	ID check completed bits	00	Not verified	
SR10 (bit2)		01	Verification mismatch	
		10	Reserved	
		11	Verified	
SR9 (bit1)	Data reception time out	Time out	Normal operation	
SR8 (bit0)	Reserved	-	-	

Table 19 Status register 1 (SRD1)

Full Status Check

Results from executed erase and program operations can be known

by running a full status check. Figure 66 shows a flowchart of the full status check and explains how to remedy errors which occur.

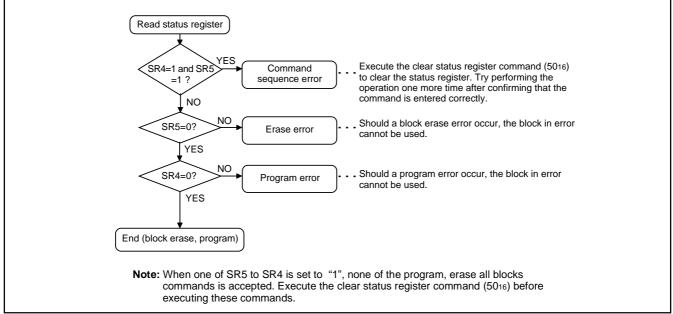


Fig. 66 Full status check flowchart and remedial procedure for errors

Example Circuit Application for The Standard Serial I/O Mode

Figure 67 shows a circuit application for the standard serial I/O mode. Control pins will vary according to programmer, therefore see the peripheral unit manual for more information.

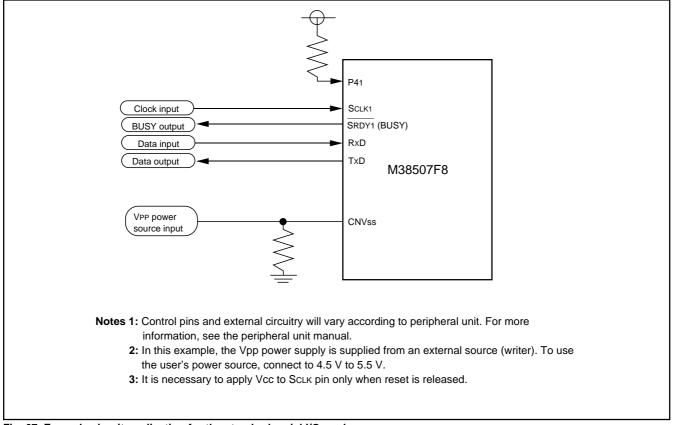


Fig. 67 Example circuit application for the standard serial I/O mode

FUNCTIONAL DESCRIPTION

Flash memory Electrical characteristics

Table 20 Absolute maximum ratings

Symbol	Parameter	Conditions	Ratings	Unit
Vcc	Power source voltage		-0.3 to 6.5	V
VI	Input voltage P00–P07, P10–P17, P20, P21, P24–P27, P30–P34, P40–P44, VREF		-0.3 to Vcc +0.3	V
VI	Input voltage P22, P23		-0.3 to 5.8	V
VI	Input voltage RESET, XIN	All voltages are based on VSS. Output transistors are cut off.	-0.3 to Vcc +0.3	V
VI	Input voltage CNVss	+ - +	-0.3 to Vcc +0.3	V
Vo	Output voltage P00–P07, P10–P17, P20, P21, P24–P27, P30–P34, P40–P44, XOUT		-0.3 to Vcc +0.3	V
Vo	Output voltage P22, P23		-0.3 to 5.8	V
Pd	Power dissipation	Ta = 25 °C	1000 (Note)	mW
Topr	Operating temperature		25±5	°C
Tstg	Storage temperature	1	-40 to 125	°C

Note: The rating becomes 300 mW at the 42P2R-A/E package.

Table 21 Flash memory mode Electrical characteristics (Ta = 25° C, Vcc = 4.5 to 5.5V unless otherwise noted)

Symbol Para			Limits			
	Parameter	Conditions	Min.	Тур.	Max.	Unit
IPP1	VPP power source current (read)	VPP = VCC			100	μA
IPP2	VPP power source current (program)	VPP = VCC			60	mA
IPP3	VPP power source current (erase)	VPP = VCC			30	mA
VIL	"L" input voltage (Note)		0		0.8	V
Vih	"H" input voltage (Note)		2.0		Vcc	V
Vpp	VPP power source voltage		4.5		5.5	V
	No.	Microcomputer mode operation at Vcc = 2.7 to 5.5V	4.5		5.5	V
Vcc V	Vcc power source voltage	Microcomputer mode operation at Vcc = 2.7 to 3.6V	3.0		3.6	V

Note: Input pins for parallel I/O mode.

AC Electrical characteristics (Ta = 25°C, Vcc = 4.5 to 5.5V unless otherwise noted)

Table 22 Read-only mode

Cumb al	Desemator		Limits		Linit
Symbol	Parameter	Min.	Тур.	Max.	Unit
tRC	Read cycle time	200			ns
ta (AD)	Address access time			100	ns
ta (CE)	CE access time			100	ns
ta (OE)	OE access time			80	ns
tCLZ	Output enable time (after CE)	0			ns
tDF(CE)	Output floating time (after CE)			25	ns
tolz	Output enable time (after OE)	0			ns
tDF(OE)	Output floating time (after OE)			25	ns
tPHZ	Output floating time (after PR)			300	ns
tOH	Output valid time (after CE, OE, address)	0			ns
toeh	Write recovery time (before read)	200			ns
tPS	RP recovery time	10			μs

Note : Timing measurement condition is showed in Figure 68.

Table 23 Read / Write mode (WE control)

Symbol	Parameter		Limits			
Symbol	Parameter	Min.	Тур.	Max.	Unit	
tWC	Write cycle time	200			ns	
tAS	Address set up time	100			ns	
tAH	Address hold time	25			ns	
tDS	Data set up time	100			ns	
tDH	Data hold time	25			ns	
tCS	CE set up time	0			ns	
tCH	CE hold time	0			ns	
tWP	WE pulse width	100			ns	
tWPH	"H" write pulse width	50			ns	
tDAP	Program time		25		μs	
tDAE	Erase all blocks time		1.5		s	
tWHRL	RY/BY delay time			200	ns	
tPS	RP recovery time	10			μs	

Note : The read timing parameter in the command write operation mode is same as that of the read-only mode. Typical value is at Vcc = 5.0 V, Ta = 25 °C condition.

Flash memory mode Electrical characteristics

(Ta = 25°C, Vcc = 4.5 to 5.5V unless otherwise noted)

Table 24 Read / Write mode (CE control)

Symbol	Parameter		Limits			
Symbol	Farameter	Min.	Тур.	Max.	Unit	
tWC	Write cycle time	200			ns	
tAS	Address set up time	100			ns	
tah	Address hold time	25			ns	
tDS	Data set up time	100			ns	
tDH	Data hold time	25			ns	
tws	WE set up time	0			ns	
twн	WE hold time	0			ns	
tCEP	CE pulse width	100			ns	
tCEPH	"H" CE pulse width	50			ns	
t DAP	Program time		25		μs	
tDAE	Erase all blocks time		1.5		s	
tehrl	RY/BY delay time			200	ns	
tPS	RP recovery time	10			μs	

Note : The read timing parameter in the command write operation mode is same as that of the read-only mode. Typical value is at Vcc = 5.0 V, Ta = 25 °C condition.

Table 25 Erase and program operation

Parameter	Min.	Тур.	Max.	Unit
Erase all blocks time		1.5		S
Block erase time		1.0		S
Program time (1byte)		25		μs

Table 26 Vcc power up / power down timing

Symbol	Parameter	Min.	Тур.	Max.	Unit
tvcs	RP = VIH set up time (after rised Vcc = Vcc min.)	10			μs

Note : Miserase or miswrite may happen, in case of noise pulse due to the power supply on or off is input to the control pins. Therefore disableing the write mode is need for prevent from memory data break at the power supply on or off. 10µs (min.) waiting time is need to initiate read or write operation after Vcc rises to Vcc min. at power supply on. The memory data is protected owing to keep the RP pin VIL level at power supply off. The RP pin must be kept VIL level for 10µs (min.) after Vcc rises to Vcc min. at the power supply on. The RP pin must be kept VIL level for 10µs (min.) after Vcc rises to Vcc min. at the power supply on. The RP pin must be kept VIL level at power supply off. RP pin doesn't have latch mode, so RP pin must be kept VIH level during read, erase and program operation.

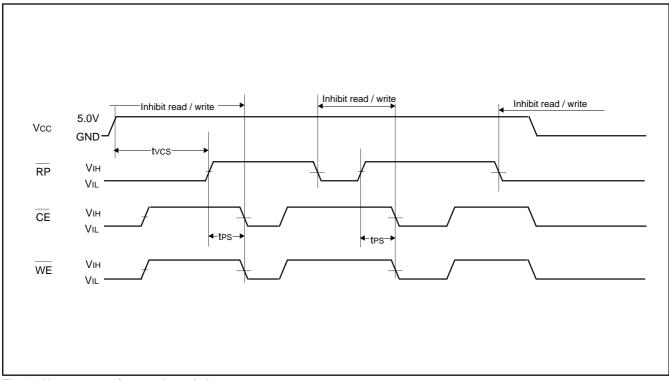


Fig. 68 Vcc power up / power down timing

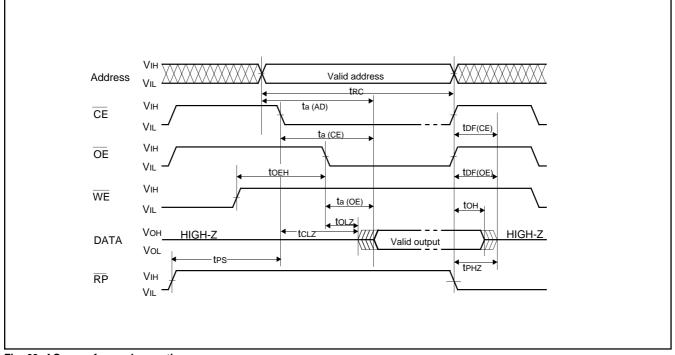


Fig. 69 AC wave for read operation

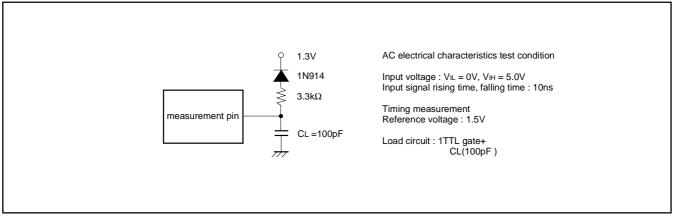
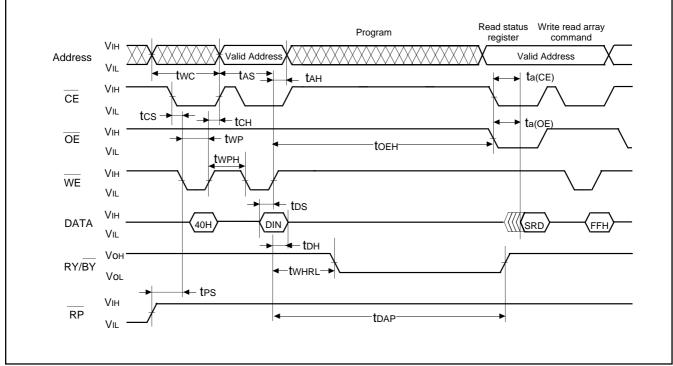
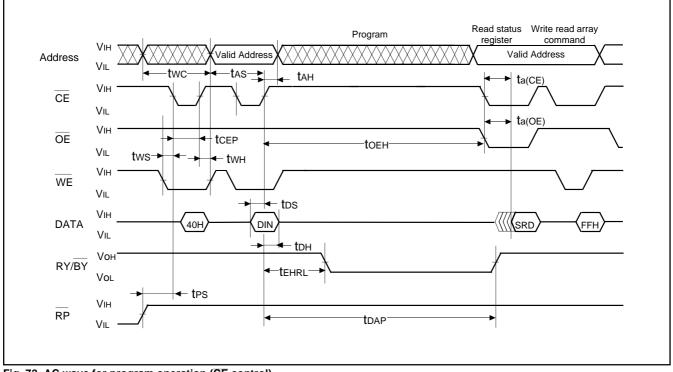
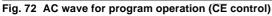


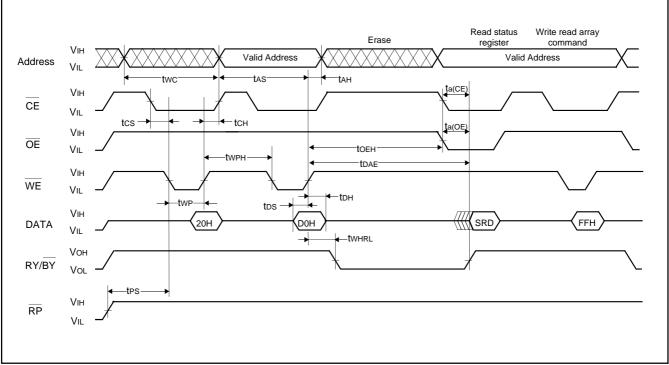
Fig. 70 AC electrical characteristics test condition for read operation

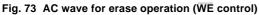












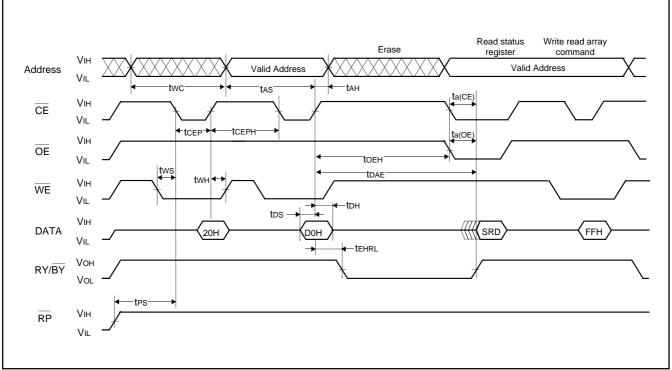


Fig. 74 AC wave for erase operation (CE control)

NOTES ON PROGRAMMING Processor Status Register

The contents of the processor status register (PS) after a reset are undefined, except for the interrupt disable flag (I) which is "1". After a reset, initialize flags which affect program execution. In particular, it is essential to initialize the index X mode (T) and the decimal mode (D) flags because of their effect on calculations.

Interrupts

The contents of the interrupt request bits do not change immediately after they have been written. After writing to an interrupt request register, execute at least one instruction before performing a BBC or BBS instruction.

Decimal Calculations

- To calculate in decimal notation, set the decimal mode flag (D) to "1", then execute an ADC or SBC instruction. After executing an ADC or SBC instruction, execute at least one instruction before executing a SEC, CLC, or CLD instruction.
- In decimal mode, the values of the negative (N), overflow (V), and zero (Z) flags are invalid.

Timers

If a value n (between 0 and 255) is written to a timer latch, the frequency division ratio is 1/(n+1).

Multiplication and Division Instructions

- The index X mode (T) and the decimal mode (D) flags do not affect the MUL and DIV instruction.
- The execution of these instructions does not change the contents of the processor status register.

Ports

The contents of the port direction registers cannot be read. The following cannot be used:

- The data transfer instruction (LDA, etc.)
- The operation instruction when the index X mode flag (T) is "1"
- The addressing mode which uses the value of a direction register as an index
- The bit-test instruction (BBC or BBS, etc.) to a direction register
- The read-modify-write instructions (ROR, CLB, or SEB, etc.) to a direction register.

Use instructions such as LDM and STA, etc., to set the port direction registers.

Serial I/O

In clock synchronous serial I/O, if the receive side is using an external clock and it is to output the $\overline{SRDY1}$ signal, set the transmit enable bit, the receive enable bit, and the $\overline{SRDY1}$ output enable bit to "1".

Serial I/O1 continues to output the final bit from the TxD pin after transmission is completed.

SOUT2 pin for serial I/O2 goes to high impedance after transmission is completed.

When an external clock is used as synchronous clock in serial I/O1 or serial I/O2, write transmission data to the transmit buffer register or serial I/O2 register while the transfer clock is "H".

A-D Converter

The comparator uses capacitive coupling amplifier whose charge will be lost if the clock frequency is too low.

Therefore, make sure that $f(X_{IN})$ in the middle/high-speed mode is at least on 500 kHz during an A-D conversion.

Do not execute the STP instruction during an A-D conversion.

Instruction Execution Time

The instruction execution time is obtained by multiplying the frequency of the internal clock ϕ by the number of cycles needed to execute an instruction.

The number of cycles required to execute an instruction is shown in the list of machine instructions.

The frequency of the internal clock ϕ is half of the XIN frequency in high-speed mode.

NOTES ON USAGE Differences between 3850 group (standard) and 3850 group (spec. H)

- (1) The absolute maximum ratings of 3850 group (spec. H) is smaller than that of 3850 group (standard).
 •Power source voltage Vcc = -0.3 to 6.5 V
 - •CNVss input voltage VI = -0.3 to Vcc +0.3 V
- (2) The oscillation circuit constants of XIN-XOUT, XCIN-XCOUT may be some differences between 3850 group (standard) and 3850 group (spec. H).
- (3) Do not write any data to the reserved area and the reserved bit. (Do not change the contents after rest.)
- (4) Fix bit 3 of the CPU mode register to "1".
- (5) Be sure to perform the termination of unused pins.

Handling of Source Pins

In order to avoid a latch-up occurrence, connect a capacitor suitable for high frequencies as bypass capacitor between power source pin (Vcc pin) and GND pin (Vss pin) and between power source pin (Vcc pin) and analog power source input pin (AVss pin). Besides, connect the capacitor to as close as possible. For bypass capacitor which should not be located too far from the pins to be connected, a ceramic capacitor of 0.01 $\mu\text{F}{-}0.1\mu\text{F}$ is recommended.

EPROM Version/One Time PROM Version/ Flash Memory Version

The CNVss pin is connected to the internal memory circuit block by a low-ohmic resistance, since it has the multiplexed function to be a programmable power source pin (VPP pin) as well.

To improve the noise reduction, connect a track between CNVss pin and Vss pin or Vcc pin with 1 to 10 $k\Omega$ resistance.

The mask ROM version track of CNVss pin has no operational interference even if it is connected to Vss pin or Vcc pin via a resistor.

DATA REQUIRED FOR MASK ORDERS/DATA REQUIRED FOR One Time PROM PROGRAMMING ORDERS/ROM PROGRAMMING METHOD

DATA REQUIRED FOR MASK ORDERS

The following are necessary when ordering a mask ROM production:

- 1. Mask ROM Order Confirmation Form*
- 2. Mark Specification Form*
- 3. Data to be written to ROM, in EPROM form (three identical copies) or one floppy disk.

DATA REQUIRED FOR One Time PROM PROGRAMMING ORDERS

The following are necessary when ordering a PROM programming service:

- 1. ROM Programming Confirmation Form*
- 2. Mark Specification Form* (only special mark with customer's trade mark logo)
- 3. Data to be programmed to PROM, in EPROM form (three identical copies) or one floppy disk.
- *For the mask ROM confirmation, the ROM programming confirmation form, and the mark specifications, refer to the "Renesas Technology Corp." Homepage (http://www.renesas.com/en/rom).

ROM PROGRAMMING METHOD

The built-in PROM of the blank One Time PROM version and buitin EPROM version can be read or programmed with a general-purpose PROM programmer using a special programming adapter. Set the address of PROM programmer in the user ROM area.

Table 27 Programming adapter

Package	Name of Programming Adapter
42P4B, 42S1B	PCA4738S-42A
42P2R-A/E	PCA4738F-42A

The PROM of the blank One Time PROM version is not tested or screened in the assembly process and following processes. To ensure proper operation after programming, the procedure shown in Figure 75 is recommended to verify programming.

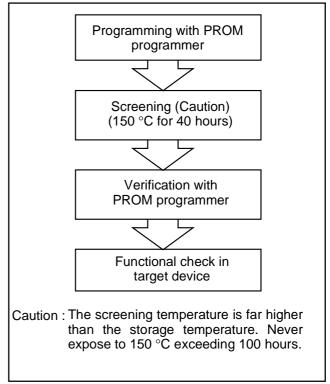


Fig. 75 Programming and testing of One Time PROM version

FUNCTIONAL DESCRIPTION SUPPLEMENT A-D Converter

A-D conversion is started by setting AD conversion completion bit to "0". During A-D conversion, internal operations are performed as follows.

- 1. After the start of A-D conversion, A-D conversion register goes to "0016".
- 2. The highest-order bit of A-D conversion register is set to "1", and the comparison voltage Vref is input to the comparator. Then, Vref is compared with analog input voltage VIN.
- 3. As a result of comparison, when Vref < VIN, the highest-order bit of A-D conversion register becomes "1". When Vref > VIN, the highest-order bit becomes "0".

Table 28 Relative formula for a reference voltage VREF of A-D converter and Vref

When $n = 0$	Vref = 0	
When n = 1 to 1023	$V_{ref} = \frac{V_{REF}}{1024} \times n$	

n: Value of A-D converter (decimal numeral)

Table 29 Change of A-D conversion register during A-D conversion

By repeating the above operations up to the lowest-order bit of the A-D conversion register, an analog value converts into a digital value.

A-D conversion completes at 61 clock cycles (15.25 μ s at f(XIN) = 8 MHz) after it is started, and the result of the conversion is stored into the A-D conversion register.

Concurrently with the completion of A-D conversion, A-D conversion interrupt request occurs, so that the AD conversion interrupt request bit is set to "1".

	Change of A-D conversion register	Value of comparison voltage (Vref)	
At start of conversion	0 0 0 0 0 0 0 0 0 0	0	
First comparison	1 0 0 0 0 0 0 0 0 0 0	VREF 2	
Second comparison	*1 1 0 0 0 0 0 0 0 0 0	$\frac{VREF}{2} \pm \frac{VREF}{4}$	
Third comparison	*1 *2 1 0 0 0 0 0 0 0	$\frac{VREF}{2} \pm \frac{VREF}{4} \pm \frac{VREF}{8}$	
~	\approx		
After completion of tenth comparison	A result of A-D conversion *1 *2 *3 *4 *5 *6 *7 *8 *9 *10	$\frac{VREF}{2} \pm \frac{VREF}{4} \pm \bullet \bullet \bullet \bullet \pm \frac{VREF}{1024}$	

***1-*10:** A result of the first comparison to the tenth comparison

FUNCTIONAL DESCRIPTION SUPPLEMENT

Figures 76 shows the A-D conversion equivalent circuit, and Figure 77 shows the A-D conversion timing chart.

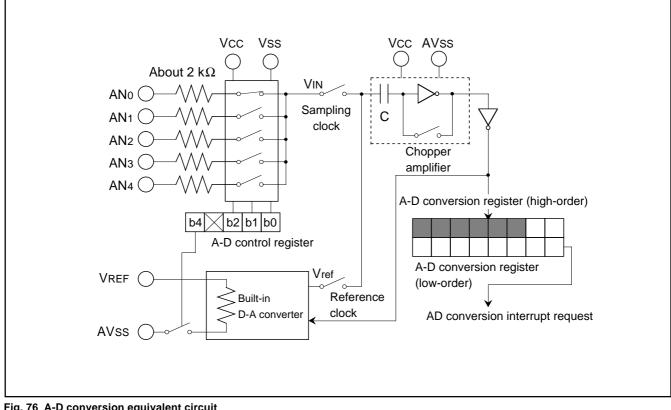


Fig. 76 A-D conversion equivalent circuit

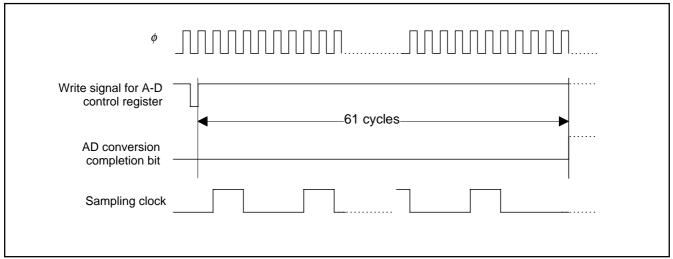


Fig. 77 A-D conversion timing chart

CHAPTER 2 APPLICATION

- 2.1 I/O port
- 2.2 Interrupt
- 2.3 Timer
- 2.4 Serial I/O
- 2.5 PWM
- 2.6 A-D converter
- 2.7 Watchdog timer
- 2.8 Reset
- 2.9 Clock generating circuit
- 2.10 Standby function
- 2.11 Flash memory mode

2.1 I/O port

2.1 I/O port

This paragraph describes the setting method of I/O port relevant registers, notes etc.

2.1.1 Memory map

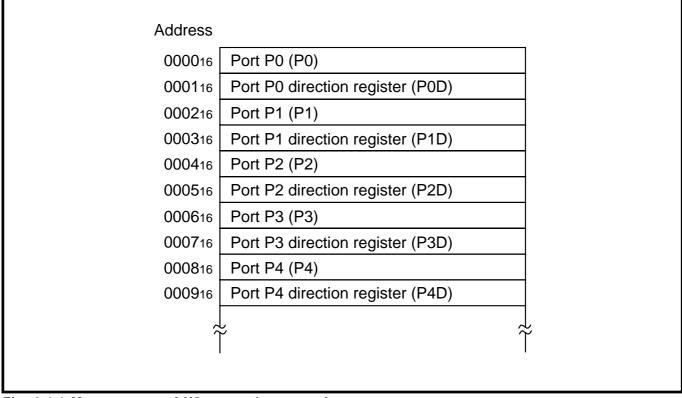


Fig. 2.1.1 Memory map of I/O port relevant registers

2.1.2 Relevant registers

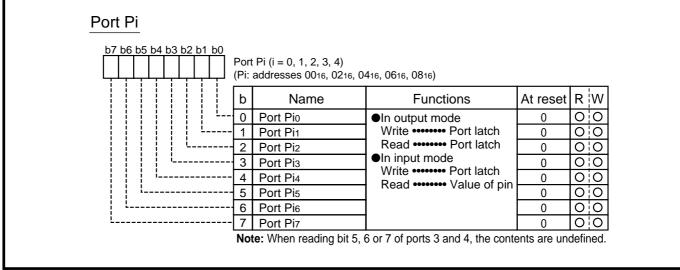
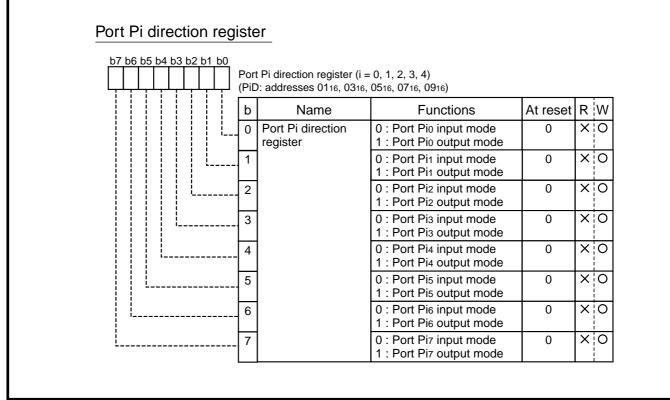


Fig. 2.1.2 Structure of Port Pi (i = 0, 1, 2, 3, 4)

2.1 I/O port





2.1.3 Terminate unused pins

Table	2.1.1	Termination	of	unused	pins
-------	-------	-------------	----	--------	------

Pins/Ports name	Termination		
P0, P1, P2, P3, P4	\bullet Set to the input mode and connect each to Vcc or Vss through a resistor of 1 k Ω		
	to 10 kΩ.		
	• Set to the output mode and open at "L" or "H" output state.		
Vref	Connect to Vss (GND).		
AVss	Connect to Vss (GND).		
Хоит	Open (only when using external clock)		

2.1 I/O port

2.1.4 Notes on I/O port

(1) Notes in standby state

In standby state^{*1}, do not make input levels of an I/O port "undefined", especially for I/O ports of the N-channel open-drain. When setting the N-channel open-drain port as an output, do not make input levels of an I/O port "undefined", too.

Pull-up (connect the port to Vcc) or pull-down (connect the port to Vss) these ports through a resistor.

When determining a resistance value, note the following points:

- External circuit
- Variation of output levels during the ordinary operation

Reason

When setting as an input port with its direction register, the transistor becomes the OFF state, which causes the ports to be the high-impedance state.

Accordingly, the potential which is input to the input buffer in a microcomputer is unstable in the state that input levels of an I/O port are "undefined". This may cause power source current.

In I/O ports of N-channel open-drain, when the contents of the port latch are "1", even if it is set as an output port with its direction register, it becomes the same phenomenon as the case of an input port.

 $\boldsymbol{\ast}1$ standby state: stop mode by executing \boldsymbol{STP} instruction

wait mode by executing WIT instruction

(2) Modifying output data with bit managing instruction

When the port latch of an I/O port is modified with the bit managing instruction*², the value of the unspecified bit may be changed.

Reason

The bit managing instructions are read-modify-write form instructions for reading and writing data by a byte unit. Accordingly, when these instructions are executed on a bit of the port latch of an I/O port, the following is executed to all bits of the port latch.

•As for bit which is set for input port:

The pin state is read in the CPU, and is written to this bit after bit managing.

•As for bit which is set for output port:

The bit value is read in the CPU, and is written to this bit after bit managing.

Note the following:

- •Even when a port which is set as an output port is changed for an input port, its port latch holds the output data.
- •As for a bit of which is set for an input port, its value may be changed even when not specified with a bit managing instruction in case where the pin state differs from its port latch contents.

 $\mathbf{\ast}\mathbf{2}$ Bit managing instructions: **SEB** and **CLB** instructions

2.1 I/O port

2.1.5 Termination of unused pins

(1) Terminate unused pins

- ① Output ports : Open
- 2 Input ports :

Connect each pin to Vcc or Vss through each resistor of 1 k Ω to 10 k Ω .

A for pins whose potential affects to operation modes such as pins CNV_{SS} , INT or others, select the V_{CC} pin or the V_{SS} pin according to their operation mode.

③ I/O ports :

• Set the I/O ports for the input mode and connect them to Vcc or Vss through each resistor of 1 k Ω to 10 k $\Omega.$

Set the I/O ports for the output mode and open them at "L" or "H".

- When opening them in the output mode, the input mode of the initial status remains until the mode of the ports is switched over to the output mode by the program after reset. Thus, the potential at these pins is undefined and the power source current may increase in the input mode. With regard to an effects on the system, thoroughly perform system evaluation on the user side.
- Since the direction register setup may be changed because of a program runaway or noise, set direction registers by program periodically to increase the reliability of program.
- ④ The AVss pin when not using the A-D converter :
 - When not using the A-D converter, handle a power source pin for the A-D converter, AVss pin as follows:

AVss: Connect to the Vss pin.

(2) Termination remarks

- ① Input ports and I/O ports :
 - Do not open in the input mode.

Reason

- The power source current may increase depending on the first-stage circuit.
- An effect due to noise may be easily produced as compared with proper termination 2 and 3 shown on the above.

2 I/O ports :

When setting for the input mode, do not connect to Vcc or Vss directly.

Reason

If the direction register setup changes for the output mode because of a program runaway or noise, a short circuit may occur between a port and Vcc (or Vss).

③ I/O ports:

When setting for the input mode, do not connect multiple ports in a lump to Vcc or Vss through a resistor.

Reason

If the direction register setup changes for the output mode because of a program runaway or noise, a short circuit may occur between ports.

• At the termination of unused pins, perform wiring at the shortest possible distance (20 mm or less) from microcomputer pins.

2.2 Interrupt

2.2 Interrupt

This paragraph explains the registers setting method and the notes relevant to the interrupt.

2.2.1 Memory map

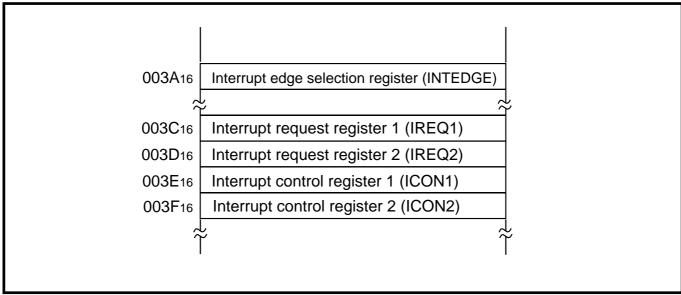


Fig. 2.2.1 Memory map of registers relevant to interrupt

2.2 Interrupt

2.2.2 Relevant registers

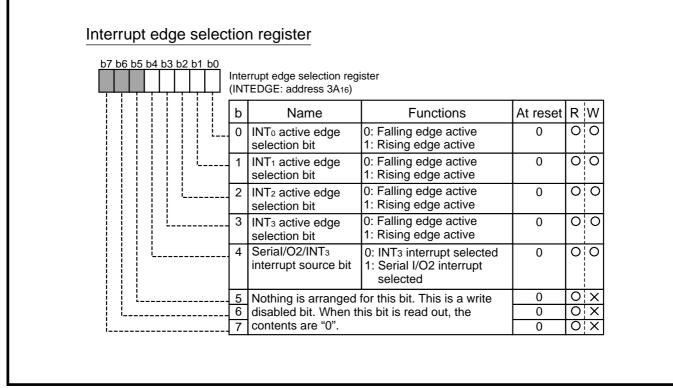


Fig. 2.2.2 Structure of Interrupt edge selection register

2.2 Interrupt

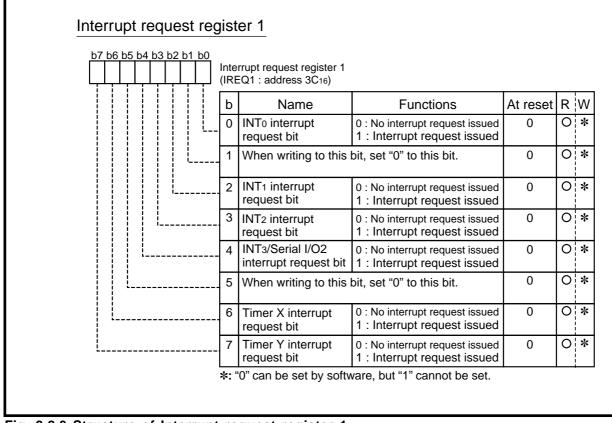


Fig. 2.2.3 Structure of Interrupt request register 1

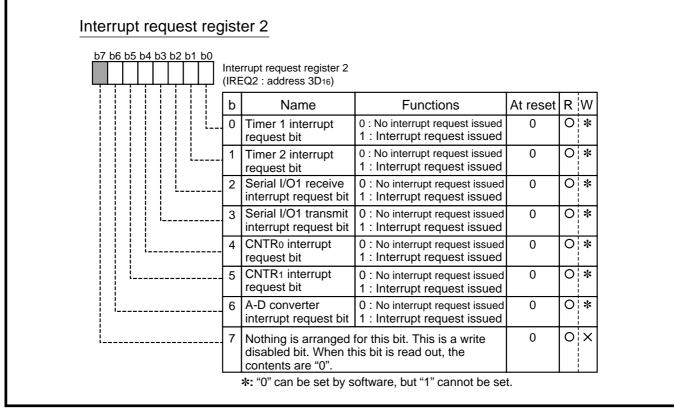


Fig. 2.2.4 Structure of Interrupt request register 2

2.2 Interrupt

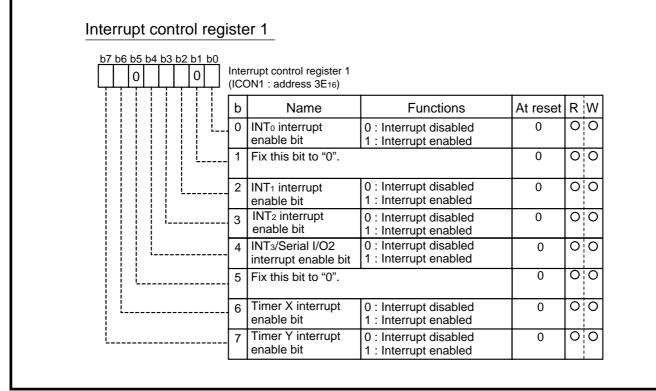
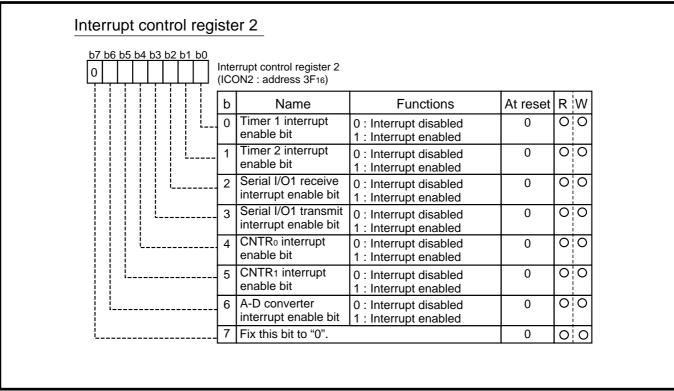


Fig. 2.2.5 Structure of Interrupt control register 1





2.2 Interrupt

2.2.3 Interrupt source

The 3850 group permits interrupts of 15 sources. These are vector interrupts with a fixed priority system. Accordingly, when two or more interrupt requests occur during the same sampling, the higher-priority interrupt is accepted first. This priority is determined by hardware, but a variety of priority processing can be performed by software, using an interrupt enable bit and an interrupt disable flag. For interrupt sources, vector addresses and interrupt priority, refer to Table 2.2.1.

	D · · ·	Vector Addresses (Note 1)		Interrupt Request	Remarks	
Interrupt Source	Priority	High	Low	Generating Conditions	Remarks	
Reset (Note 2)	1	FFFD16	FFFC16	At reset	Non-maskable	
ΙΝΤο	2	FFFB16	FFFA16	At detection of either rising or falling edge of INTo input	External interrupt (active edge selectable)	
Reserved	3	FFF916	FFF816	Reserved		
INT1	4	FFF716	FFF616	At detection of either rising or falling edge of INT1 input	External interrupt (active edge selectable)	
INT2	5	FFF516	FFF416	At detection of either rising or falling edge of INT2 input	External interrupt (active edge selectable)	
INT3	6	FFF316	FFF216	At detection of either rising or falling edge of INT3 input	External interrupt (active edge selectable)	
Serial I/O2				At completion of serial I/O2 data transfer	Switch by Serial I/O2/INT3 interrupt source bit	
Reserved	7	FFF116	FFF016	Reserved		
Timer X	8	FFEF16	FFEE16	At timer X underflow		
Timer Y	9	FFED16	FFEC16	At timer Y underflow		
Timer 1	10	FFEB16	FFEA16	At timer 1 underflow	STP release timer underflow	
Timer 2	11	FFE916	FFE816	At timer 2 underflow		
Serial I/O1 received	12	FFE716	FFE616	At completion of serial I/O1 data reception	Valid when serial I/O1 is selected	
Serial I/O1 transmit	13	FFE516	FFE416	At completion of serial I/O1 transfer shift or when transmis- sion buffer is empty	Valid when serial I/O1 is selected	
CNTR ₀	14	FFE316	FFE216	At detection of either rising or falling edge of CNTR0 input	External interrupt (active edge selectable)	
CNTR1	15	FFE116	FFE016	At detection of either rising or falling edge of CNTR1 input	External interrupt (active edge selectable)	
A-D converter	16	FFDF16	FFDE16	At completion of A-D conversion		
BRK instruction	17	FFDD16	FFDC16	At BRK instruction execution	Non-maskable software interrupt	

Table 2.2.1 Interrupt sources	s, vector addresses	and priority of 3850 group
-------------------------------	---------------------	----------------------------

Notes 1: Vector addresses contain interrupt jump destination addresses.

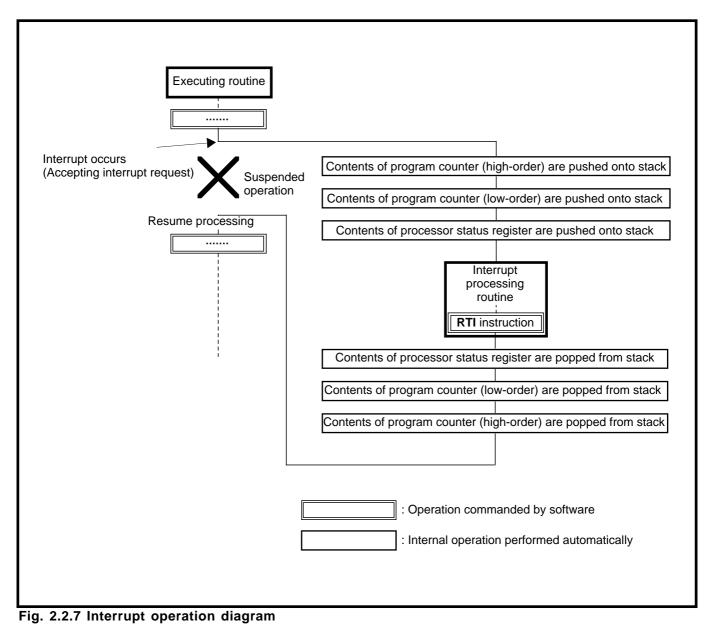
2: Reset function in the same way as an interrupt with the highest priority.

2.2.4 Interrupt operation

When an interrupt request is accepted, the contents of the following registers just before acceptance of the interrupt requests are automatically pushed onto the stack area in the order of (1, 2) and (3).

①High-order contents of program counter (PC_H)
 ②Low-order contents of program counter (PC_L)
 ③Contents of processor status register (PS)

After the contents of the above registers are pushed onto the stack area, the accepted interrupt vector address enters the program counter and consequently the interrupt processing routine is executed. When the RTI instruction is executed at the end of the interrupt processing routine, the contents of the above registers pushed onto the stack area are restored to the respective registers in the order of (3), (2) and (1); and the microcomputer resumes the processing executed just before acceptance of the interrupts. Figure 2.2.7 shows an interrupt operation diagram.



2.2 Interrupt

(1) Processing upon acceptance of interrupt request

Upon acceptance of an interrupt request, the following operations are automatically performed. ^①The processing being executed is stopped.

- The contents of the program counter and the processor status register are pushed onto the stack area. Figure 2.2.8 shows the changes of the stack pointer and the program counter upon acceptance of an interrupt request.
- ③Concurrently with the push operation, the jump destination address (the beginning address of the interrupt processing routine) of the occurring interrupt stored in the vector address is set in the program counter, then the interrupt processing routine is executed.
- ④After the interrupt processing routine is started, the corresponding interrupt request bit is automatically cleared to "0". The interrupt disable flag is set to "1" so that multiple interrupts are disabled.

Accordingly, for executing the interrupt processing routine, it is necessary to set the jump destination address in the vector area corresponding to each interrupt.

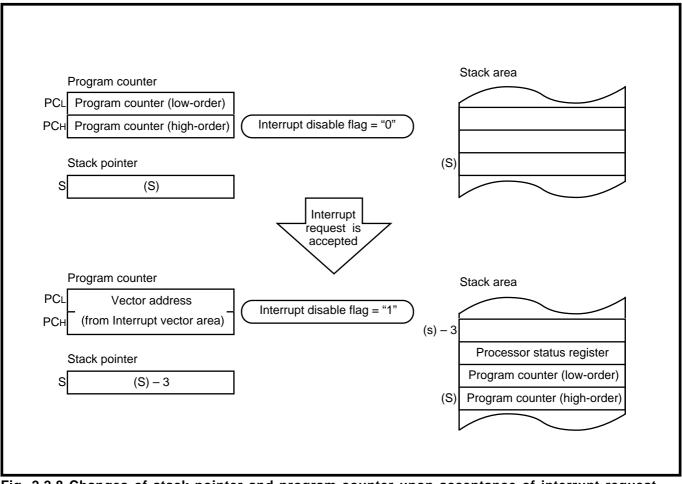


Fig. 2.2.8 Changes of stack pointer and program counter upon acceptance of interrupt request

2.2 Interrupt

(2) Timing after acceptance of interrupt request

The interrupt processing routine begins with the machine cycle following the completion of the instruction that is currently being executed.

Figure 2.2.9 shows the time up to execution of interrupt processing routine and Figure 2.2.10 shows the timing chart after acceptance of interrupt request.

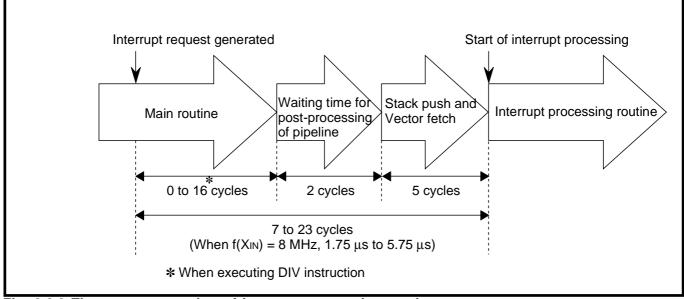
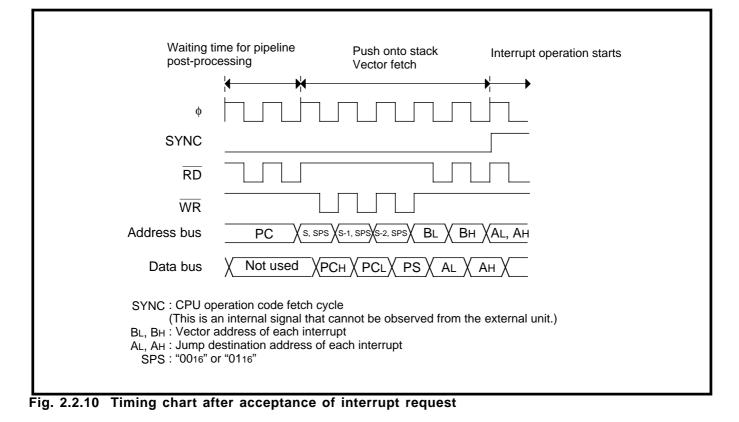


Fig. 2.2.9 Time up to execution of interrupt processing routine



2.2 Interrupt

2.2.5 Interrupt control

The acceptance of all interrupts, excluding the BRK instruction interrupt, can be controlled by the interrupt request bit, interrupt enable bit, and an interrupt disable flag, as described in detail below. Figure 2.2.11 shows an interrupt control diagram.

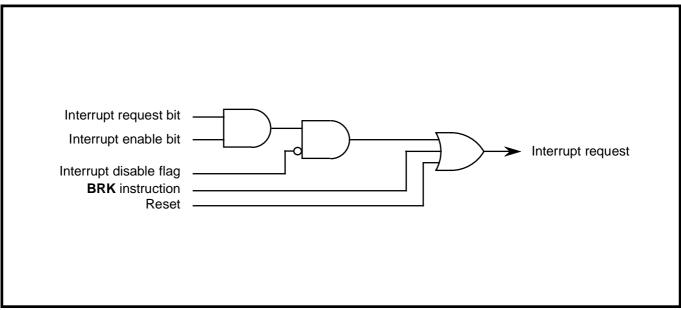


Fig. 2.2.11 Interrupt control diagram

The interrupt request bit, interrupt enable bit and interrupt disable flag function independently and do not affect each other. An interrupt is accepted when all the following conditions are satisfied.

•Interrupt request bit "1"

•Interrupt enable bit "1"

•Interrupt disable flag "0"

Though the interrupt priority is determined by hardware, a variety of priority processing can be performed by software using the above bits and flag. Table 2.2.2 shows a list of interrupt control bits according to the interrupt source.

(1) Interrupt request bits

The interrupt request bits are allocated to the interrupt request register 1 (address $3C_{16}$) and interrupt request register 2 (address $3D_{16}$).

The occurrence of an interrupt request causes the corresponding interrupt request bit to be set to "1". The interrupt request bit is held in the "1" state until the interrupt is accepted. When the interrupt is accepted, this bit is automatically cleared to "0".

Each interrupt request bit can be set to "0", but cannot be set to "1", by software.

(2) Interrupt enable bits

The interrupt enable bits are allocated to the interrupt control register 1 (address $003E_{16}$) and the interrupt control register 2 (address $3F_{16}$).

The interrupt enable bits control the acceptance of the corresponding interrupt request.

When an interrupt enable bit is "0", the corresponding interrupt request is disabled. If an interrupt request occurs when this bit is "0", the corresponding interrupt request bit is set to "1" but the interrupt is not accepted. In this case, unless the interrupt request bit is set to "0" by software, the interrupt request bit remains in the "1" state.

When an interrupt enable bit is "1", the corresponding interrupt is enabled. If an interrupt request occurs when this bit is "1", the interrupt is accepted (when interrupt disable flag = "0"). Each interrupt enable bit can be set to "0" or "1" by software.

2.2 Interrupt

(3) Interrupt disable flag

The interrupt disable flag is allocated to bit 2 of the processor status register. The interrupt disable flag controls the acceptance of interrupt request except BRK instruction.

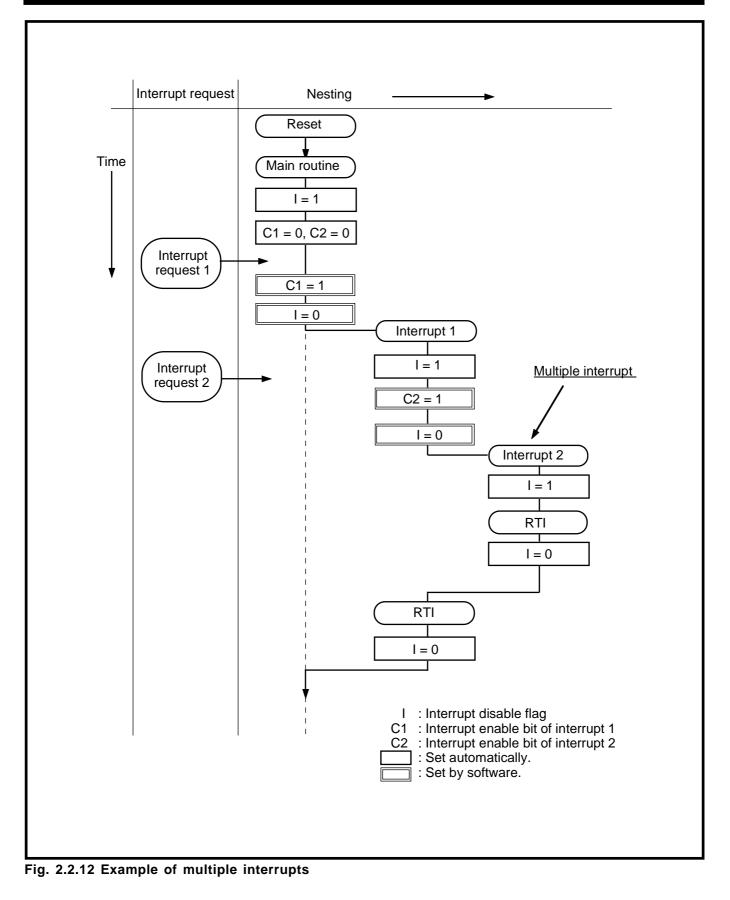
When this flag is "1", the acceptance of interrupt requests is disabled. When the flag is "0", the acceptance of interrupt requests is enabled. This flag is set to "1" with the SEI instruction and is set to "0" with the CLI instruction.

When a main routine branches to an interrupt processing routine, this flag is automatically set to "1", so that multiple interrupts are disabled. To use multiple interrupts, set this flag to "0" with the CLI instruction within the interrupt processing routine. Figure 2.2.12 shows an example of multiple interrupts.

Interrupt course	Interrupt enable bit		Interrupt request bit		
Interrupt source	Address	Bit	Address	Bit	
INT ₀	003E16	b0	003C ₁₆	b0	
INT ₁	003E16	b2	003C ₁₆	b2	
INT ₂	003E16	b3	003C ₁₆	b3	
INT ₃ /Serial I/O2	003E16	b4	003C ₁₆	b4	
Timer X	003E16	b6	003C ₁₆	b6	
Timer Y	003E16	b7	003C ₁₆	b7	
Timer 1	003F16	b0	003D ₁₆	b0	
Timer 2	003F ₁₆	b1	003D ₁₆	b1	
Serial I/O1 reception	003F16	b2	003D ₁₆	b2	
Serial I/O1 transmission	003F ₁₆	b3	003D ₁₆	b3	
CNTR ₀	003F ₁₆	b4	003D ₁₆	b4	
CNTR ₁	003F ₁₆	b5	003D ₁₆	b5	
A-D converter	003F ₁₆	b6	003D ₁₆	b6	

Table 2.2.2 List of interrupt bits according to interrupt source

2.2 Interrupt



2.2 Interrupt

2.2.6 INT interrupt

The INT interrupt requests is generated when the microcomputer detects a level change of each INT pin (INT₀–INT₃).

(1) Active edge selection

 INT_0-INT_3 can be selected from either a falling edge or rising edge detection as an active edge by the interrupt edge selection register. In the "0" state, the falling edge of the corresponding pin is detected. In the "1" state, the rising edge of the corresponding pin is detected.

(2) INT₃ interrupt source selection

Which of interrupt source of the serial $I/O2/INT_3$ interrupt source can be selected by the serial $I/O2/INT_3$ interrupt source bit (bit 4 of address $3A_{16}$). (Set this bit to "0" when using INT₃.)

2.2 Interrupt

2.2.7 Notes on interrupts

(1) Change of relevant register settings

When the setting of the following registers or bits is changed, the interrupt request bit may be set to "1". When not requiring the interrupt occurrence synchronized with these setting, take the following sequence.

•Interrupt edge selection register (address 3A₁₆)

•Timer XY mode register (address 2316)

Set the above listed registers or bits as the following sequence.

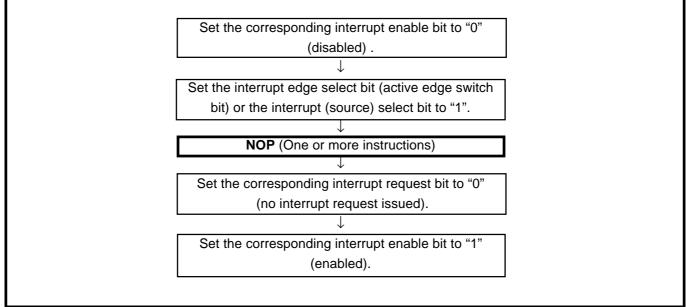


Fig. 2.2.13 Sequence of changing relevant register

Reason

When setting the followings, the interrupt request bit may be set to "1". •When setting external interrupt active edge

Concerned register: Interrupt edge selection register (address 3A₁₆)

Timer XY mode register (address 23₁₆)

•When switching interrupt sources of an interrupt vector address where two or more interrupt sources are allocated.

Concerned register: Interrupt edge selection register (address 3A16)

2.2 Interrupt

(2) Check of interrupt request bit

• When executing the **BBC** or **BBS** instruction to an interrupt request bit of an interrupt request register immediately after this bit is set to "0" by using a data transfer instruction, execute one or more instructions before executing the **BBC** or **BBS** instruction.

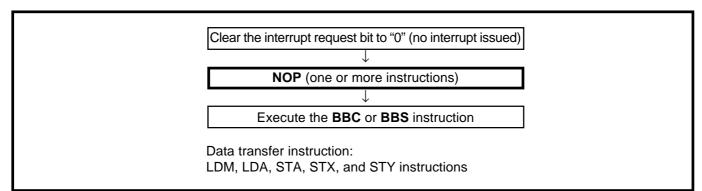


Fig. 2.2.14 Sequence of check of interrupt request bit

Reason

If the BBC or BBS instruction is executed immediately after an interrupt request bit of an interrupt request register is cleared to "0", the value of the interrupt request bit before being cleared to "0" is read.

2.3 Timer

2.3 Timer

This paragraph explains the registers setting method and the notes relevant to the timers.

2.3.1 Memory map

Address	${\approx}$
002016	Prescaler 12 (PRE12)
002116	Timer 1 (T1)
002216	Timer 2 (T2)
002316	Timer XY mode register (TM)
002416	Prescaler X (PREX)
002516	Timer X (TX)
002616	Prescaler Y (PREY)
002716	Timer Y (TY)
002816	Timer count source selection register (TCSS)
Ŕ	\div
003C16	Interrupt request register 1 (IREQ1)
003D16	Interrupt request register 2 (IREQ2)
003E16	Interrupt control register 1 (ICON1)
003F16	Interrupt control register 2 (ICON2)

Fig. 2.3.1 Memory map of registers relevant to timers

2.3.2 Relevant registers

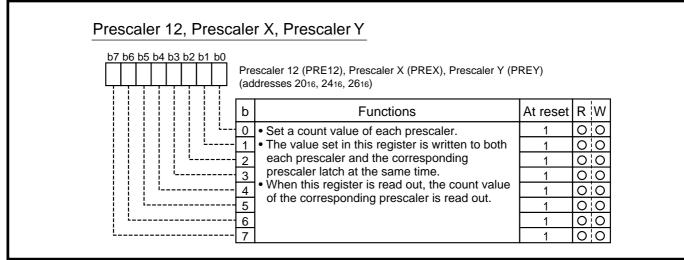


Fig. 2.3.2 Structure of Prescaler 12, Prescaler X, Prescaler Y

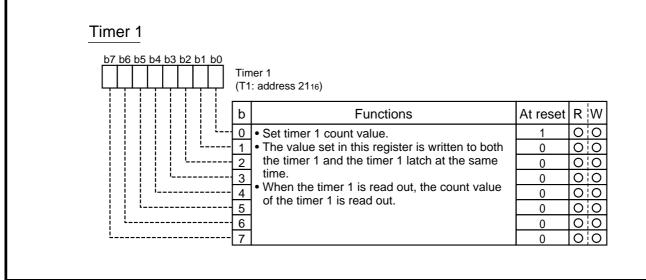
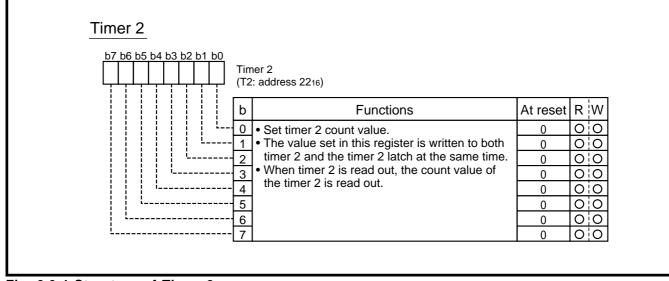
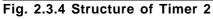


Fig. 2.3.3 Structure of Timer 1





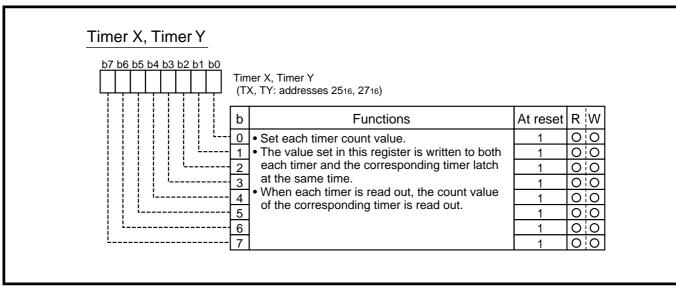


Fig. 2.3.5 Structure of Timer X, Timer Y

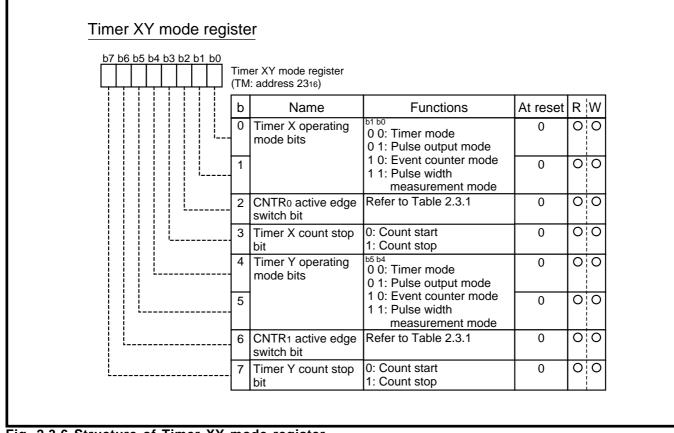


Fig. 2.3.6 Structure of Timer XY mode register

Table 2.3.1 CNTR₀/CNTR₁ active edge switch bit function

Timer X /Timer Y	Set	Timer function	CNTR ₀ / CNTR ₁ interrupt request
operation modes	value		occurrence source
Timer mode	"0"	No influence to timer count	CNTR ₀ /CNTR ₁ input signal falling edge
	"1"	No influence to timer count	CNTR ₀ /CNTR ₁ input signal rising edge
Pulse output	"0"	Pulse output start: Beginning	Output signal falling edge count
mode		at "H" level	
	"1"	Pulse output start: Beginning	Output signal rising edge count
		at "L" level	
Event counter	"0"	Rising edge count	Input signal falling edge count
mode	"1"	Falling edge count	Input signal rising edge count
Pulse width	"0"	"H" level width measurement	Input signal falling edge count
measurement mode	"1"	"L" level width measurement	Input signal rising edge count

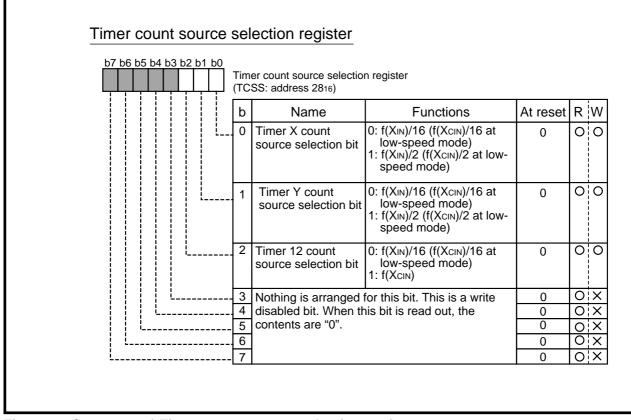


Fig. 2.3.7 Structure of Timer count source selection register

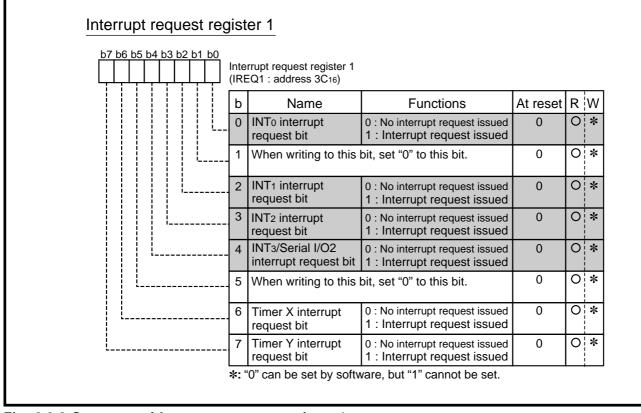


Fig. 2.3.8 Structure of Interrupt request register 1

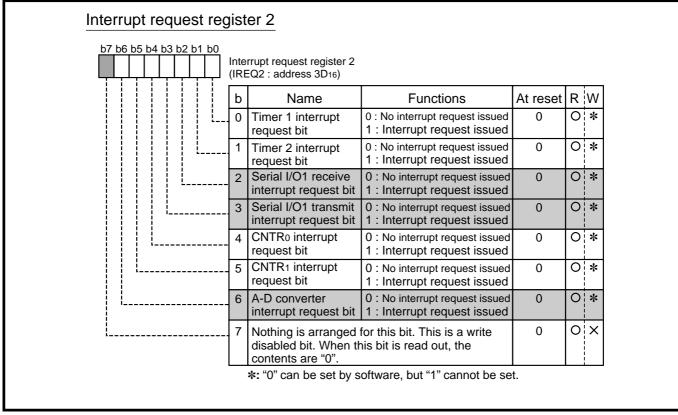


Fig. 2.3.9 Structure of Interrupt request register 2

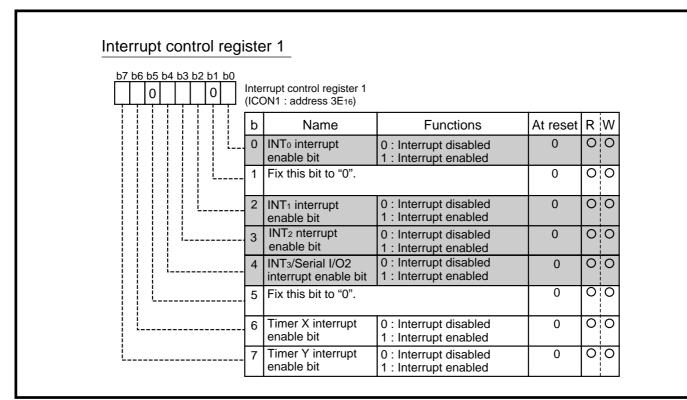


Fig. 2.3.10 Structure of Interrupt control register 1

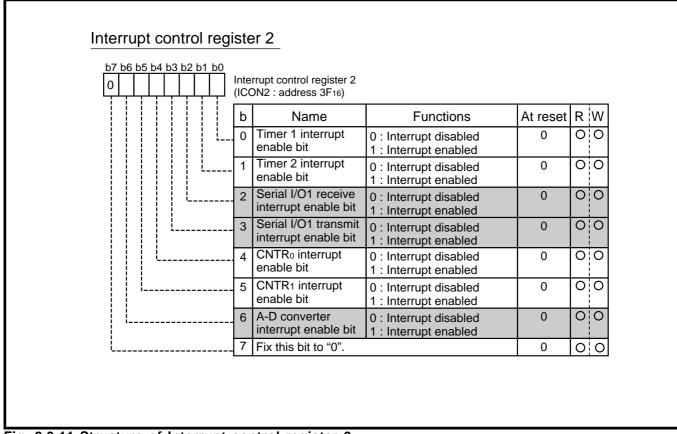


Fig. 2.3.11 Structure of Interrupt control register 2

2.3.3 Timer application examples

(1) Basic functions and uses

[Function 1] Control of Event interval (Timer X, Timer Y, Timer 1, Timer 2)

When a certain time, by setting a count value to each timer, has passed, the timer interrupt request occurs.

<Use>

•Generation of an output signal timing

•Generation of a wait time

[Function 2] Control of Cyclic operation (Timer X, Timer Y, Timer 1, Timer 2)

The value of the timer latch is automatically written to the corresponding timer each time the timer underflows, and each timer interrupt request occurs in cycles.

<Use>

•Generation of cyclic interrupts

•Clock function (measurement of 250 ms); see Application example 1

•Control of a main routine cycle

[Function 3] Output of Rectangular waveform (Timer X, Timer Y)

The output level of the CNTR pin is inverted each time the timer underflows (in the pulse output mode).

<Use>

•Piezoelectric buzzer output; see Application example 2

•Generation of the remote control carrier waveforms

[Function 4] Count of External pulses (Timer X, Timer Y)

External pulses input to the CNTR pin are counted as the timer count source (in the event counter mode).

<Use>

•Frequency measurement; see Application example 3

•Division of external pulses

•Generation of interrupts due to a cycle using external pulses as the count source; count of a reel pulse

[Function 5] Measurement of External pulse width (Timer X, Timer Y)

The "H" or "L" level width of external pulses input to CNTR pin is measured (in the pulse width measurement mode).

<Use>

•Measurement of external pulse frequency (measurement of pulse width of FG pulse* for a motor); see Application example 4

•Measurement of external pulse duty (when the frequency is fixed)

FG pulse*: Pulse used for detecting the motor speed to control the motor speed.

2.3 Timer

(2) Timer application example 1: Clock function (measurement of 250 ms)

Outline: The input clock is divided by the timer so that the clock can count up at 250 ms intervals. **Specifications**: •The clock $f(X_{IN}) = 4.19$ MHz (2²² Hz) is divided by the timer X.

•The clock is counted up in the process routine of the timer X interrupt which occurs at 250 ms intervals.

Figure 2.3.12 shows the timers connection and setting of division ratios; Figure 2.3.13 shows the relevant registers setting; Figure 2.3.14 shows the control procedure.

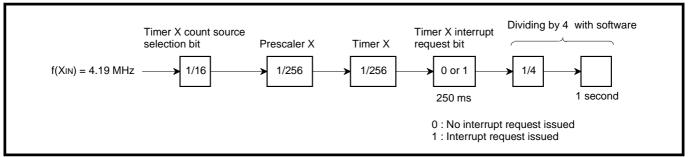
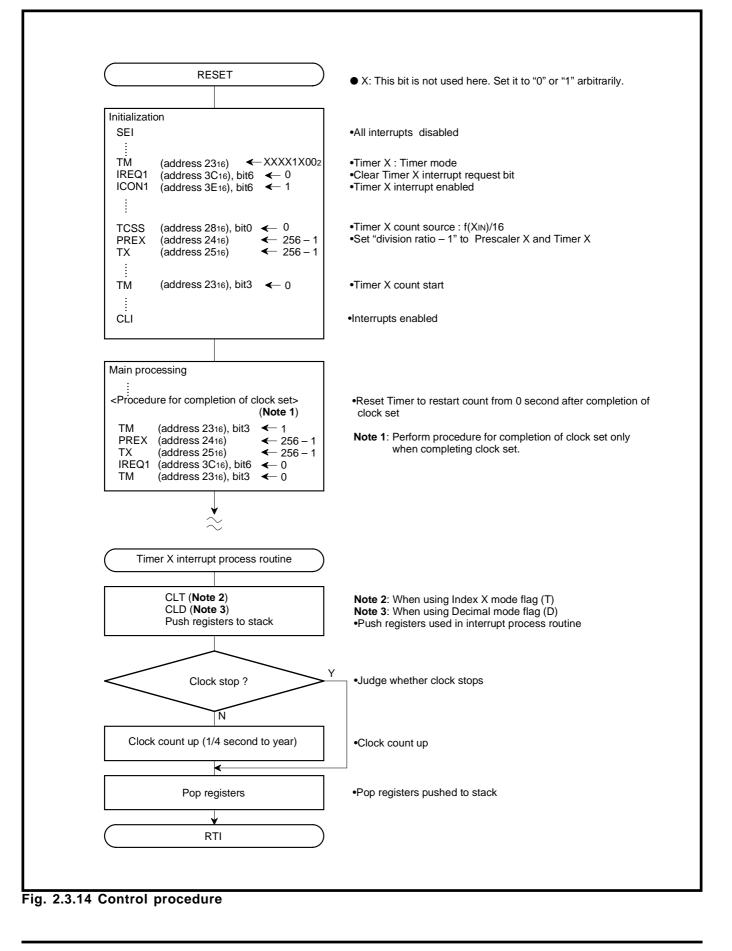


Fig. 2.3.12 Timers connection and setting of division ratios

TCSS	Timer count source selection register (address 28 ₁₆) b7 b0 b0 Timer X count source : f(XIN)/16
ТМ	Timer XY mode register (address 23 ₁₆) b7 b0 1 0 0 Timer X operating mode: Timer mode Timer X count: Stop Clear to "0" when starting count.
PREX	Prescaler X (address 24 ₁₆) b7 b0 255
тх	Timer X (address 2516) Set "division ratio – 1" b7 b0 255
IREQ1	Interrupt request register 1 (address 3C ₁₆)
ICON1	Interrupt control register 1 (address 3E ₁₆) b7 1 1 Timer X interrupt: Enabled

Fig. 2.3.13 Relevant registers setting



2.3 Timer

(3) Timer application example 2: Piezoelectric buzzer output

Outline: The rectangular waveform output function of the timer is applied for a piezoelectric buzzer output.

Specifications: •The rectangular waveform, dividing the clock $f(X_{IN}) = 4.19$ MHz (2²² Hz) into about 2 kHz (2048 Hz), is output from the P27/CNTR₀ pin.

•The level of the P27/CNTR₀ pin is fixed to "H" while a piezoelectric buzzer output stops.

Figure 2.3.15 shows a peripheral circuit example, and Figure 2.3.16 shows the timers connection and setting of division ratios. Figure 2.3.17 shows the relevant registers setting, and Figure 2.3.18 shows the control procedure.

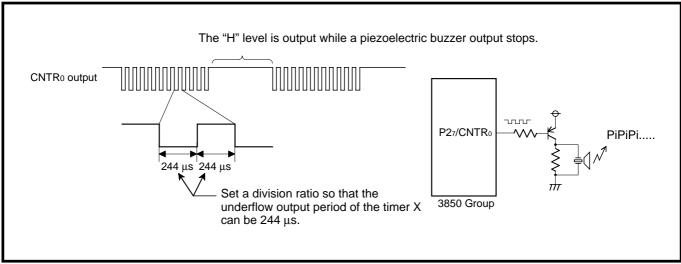


Fig. 2.3.15 Peripheral circuit example

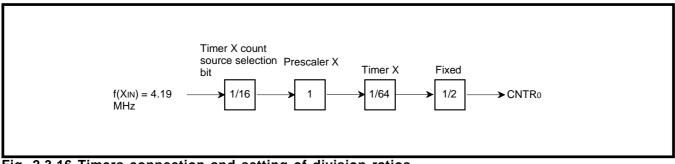
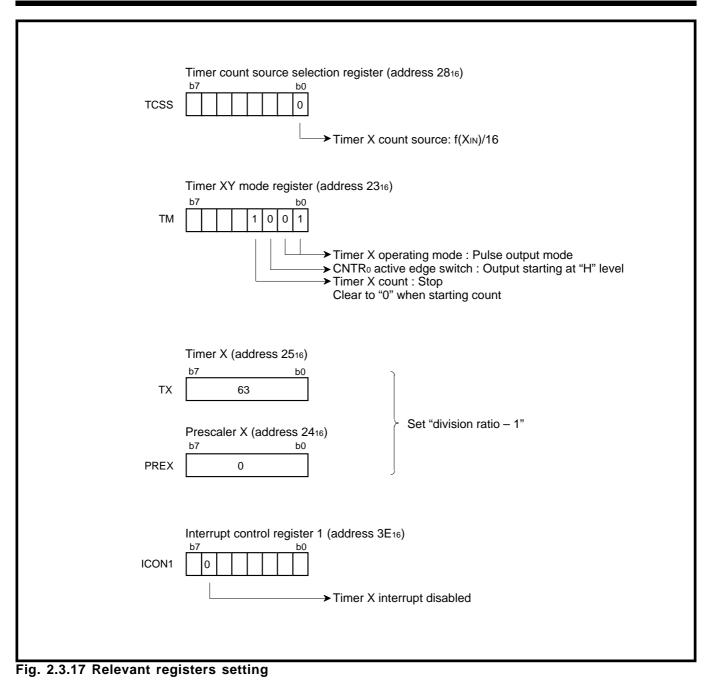
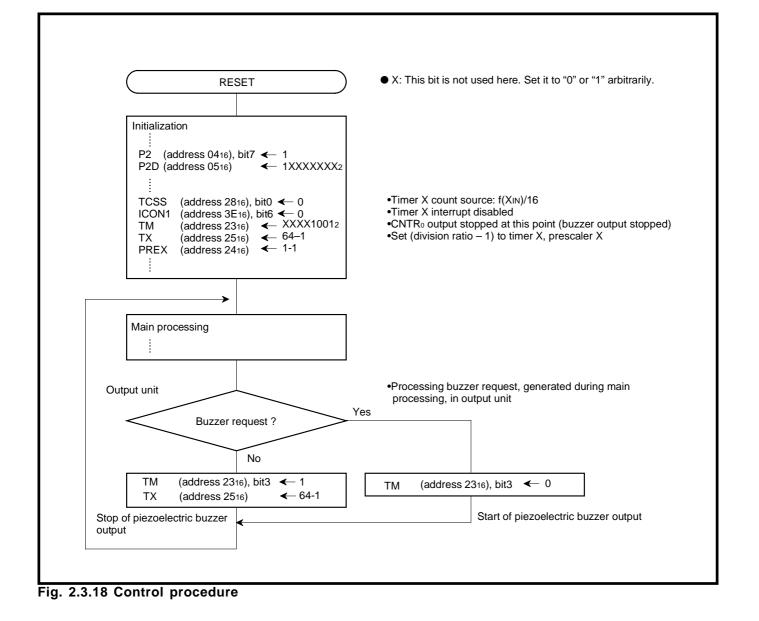


Fig. 2.3.16 Timers connection and setting of division ratios





2.3 Timer

(4) Timer application example 3: Frequency measurement

Outline: The following two values are compared to judge whether the frequency is within a valid range.

•A value by counting pulses input to P4₀/CNTR₁ pin with the timer.

A reference value

Specifications: •Clock $f(X_{IN}) = 4.19 \text{ MHz} (2^{22} \text{ Hz})$

- The pulse is input to the P4₀/CNTR₁ pin and counted by the timer Y.
 A count value is read out at about 2 ms intervals, the timer 1 interrupt interval. When the count value is 28 to 40, it is judged that the input pulse is valid.
 Because the timer is a down-counter, the count value is compared with 227 to 215 (Note).
- Note: 227 to $215 = \{255 \text{ (initial value of counter)} 28\}$ to $\{255 40\}$; 28 to 40 means the number of valid value.

Figure 2.3.19 shows the judgment method of valid/invalid of input pulses; Figure 2.3.20 shows the relevant registers setting; Figure 2.3.21 shows the control procedure.

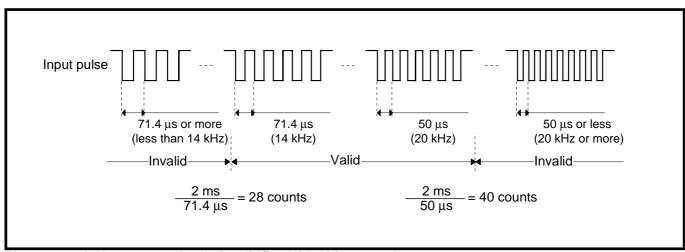
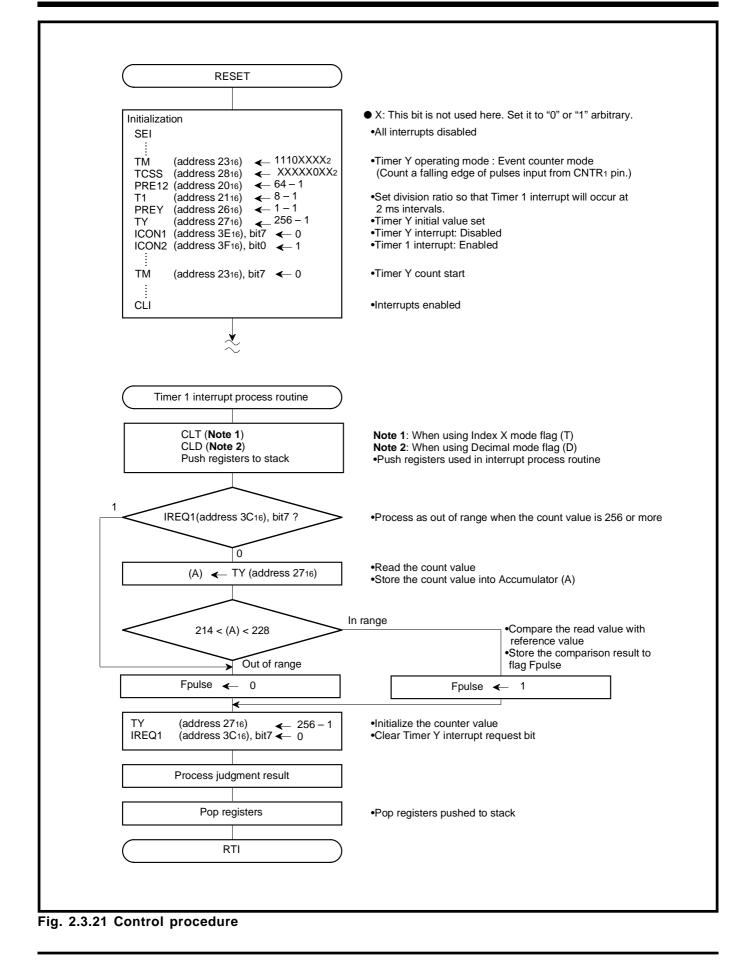


Fig. 2.3.19 Judgment method of valid/invalid of input pulses

ТМ	Timer XY mode register (address 23 ₁₆) ^{b7} b0 1 1 1 0 Timer Y operating mode: Event counter mode CNTR1 active edge switch: Falling edge count Timer Y count: Stop Clear to "0" when starting count
TCSS	Timer count source selection register (address 28 ₁₆) b7 $b0$ 0 $b0$ Timer 1 count source: f(X _{IN})/16
PRE12	Prescaler 12 (address 20 ₁₆) b7 b0 63
T1	Timer 1 (address 2116) b7 b0 7 Set "division ratio – 1"
PREY	Prescaler Y (address 2616) b7 b0 0
TY	Timer Y (address 2716)b7b0255Set 255 just before counting pulses (After a certain time has passed, the number of input pulses is decreased from this value.)
ICON1	Interrupt control register 1 (address 3E ₁₆) b7 b0 0 b0 Timer Y interrupt: Disabled
ICON2	Interrupt control register 2 (address 3F ₁₆) b7 b0 1 Timer 1 interrupt: Enabled
IREQ1 Fig. 2.3.20 Releva	Interrupt request register 1 (address 3C ₁₆) b7 b0 0 J J J J J J J J J J J J J J J J J J J



2.3 Timer

(5) Timer application example 4: Measurement of FG pulse width for motor

Outline: The timer X counts the "H" level width of the pulses input to the CNTR₀ pin. An underflow is detected by the timer X interrupt and an end of the input pulse "H" level is detected by the CNTR₀ interrupt.

Specifications: •The timer X counts the "H" level width of the FG pulse input to the CNTR₀ pin.

<Example>

When the clock frequency is 4.19 MHz, the count source is 3.8 μ s, which is obtained by dividing the clock frequency by 16. Measurement can be made up to 250 ms in the range of FFFF₁₆ to 0000₁₆.

Figure 2.3.22 shows the timers connection and setting of division ratio; Figure 2.3.23 shows the relevant registers setting; Figure 2.3.24 shows the control procedure.

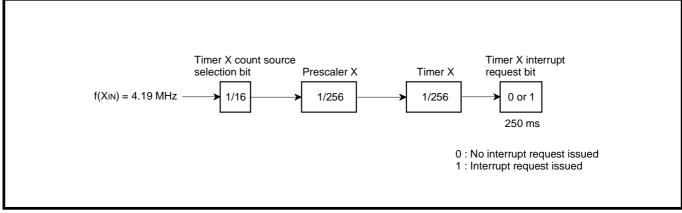


Fig. 2.3.22 Timers connection and setting of division ratios

тм	Timer XY mode register (address 23 ₁₆) b7 b0 1 0 1 1 Timer X operating mode: Pulse width measurement mode CNTRo active edge switch: "H" level width measurement Timer X count: Stop
PREX	Clear to "0" when starting count Prescaler X (address 24 ₁₆) b7 b0 255
тх	Timer X (address 2516) Set "division ratio – 1" b7 b0 255
IREQ1	Interrupt request register 1 (address 3C ₁₆) b7 b0 0 Timer X interrupt request (Set to "1" automatically when Timer X underflows)
ICON1	Interrupt control register 1 (address 3E ₁₆) b7 b0 1 Timer X interrupt: Enabled
IREQ2	Interrupt request register 2 (address 3D ₁₆) b7 b0 0 CNTR0 interrupt request (Set to "1" automatically when "H" level input came to the end)
ICON2	Interrupt control register 2 (address 3F ₁₆) b7 b0 1 1 CNTR0 interrupt: Enabled

Fig. 2.3.23 Relevant registers setting

RESET	● X: This bit is not used here. Set it to "0" or "1" arbitrarily.
Initialization	•All interrupts disabled
SEI $\begin{array}{c} & \\ & \\ & \\ & \\ & \\ & \\ & \\ & \\ & \\ & $	 All Interrupts disabled Timer X: Pulse width measurement mode (Measure "H" level of pulses input from CNTRo pin.) Set division ratio so that Timer X interrupt will occur at 250 ms intervals. Clear Timer X interrupt request bit Timer X interrupt enabled Clear CNTRo interrupt request bit CNTRo interrupt enabled
TM (address 2316), bit3 ← 0	•Timer X count start
CLI	•Interrupts enabled
\sim Timer X interrupt process routine	
Error processing	•Error occured
RTI	
RTI CNTRo interrupt process routine	
	Note 1: When using Index X mode flag (T) Note 2: When using Decimal mode flag (D) •Push registers used in interrupt process routine
CNTR₀ interrupt process routine CLT (Note 1) CLD (Note 2) Push registers to stack (A) ← PREX Low-order 8-bit result of ← Inverted (A) pulse width measurement (A) ← TX High-order 8-bit result of ← Inverted (A)	Note 2: When using Decimal mode flag (D)
CNTR₀ interrupt process routine CLT (Note 1) CLD (Note 2) Push registers to stack (A) ← PREX Low-order 8-bit result of ← Inverted (A) pulse width measurement (A) ← TX	Note 2: When using Decimal mode flag (D) •Push registers used in interrupt process routine
CNTR₀ interrupt process routine CLT (Note 1) CLD (Note 2) Push registers to stack (A) ← PREX Low-order 8-bit result of ← Inverted (A) pulse width measurement (A) ← TX High-order 8-bit result of ← Inverted (A) pulse width measurement PREX (address 2416) ← 256–1	 Note 2: When using Decimal mode flag (D) Push registers used in interrupt process routine Read the count value and store it to RAM Division ratio set so that Timer X interrupt will
CNTR ₀ interrupt process routine CLT (Note 1) CLD (Note 2) Push registers to stack (A) \leftarrow PREX Low-order 8-bit result of \leftarrow Inverted (A) pulse width measurement (A) \leftarrow TX High-order 8-bit result of \leftarrow Inverted (A) pulse width measurement PREX (address 2416) \leftarrow 256–1 TX (address 2516) \leftarrow 256–1	 Note 2: When using Decimal mode flag (D) Push registers used in interrupt process routine Read the count value and store it to RAM Division ratio set so that Timer X interrupt will occur at 250 ms intervals.

2.3.4 Notes on timer

- If a value n (between 0 and 255) is written to a timer latch, the frequency division ratio is 1/(n+1).
- When switching the count source by the timer 12, X and Y count source selection bits, the value of timer count is altered in unconsiderable amount owing to generating of thin pulses in the count input signals.

Therefore, select the timer count source before set the value to the prescaler and the timer.

2.4 Serial I/O

2.4 Serial I/O

This paragraph explains the registers setting method and the notes relevant to the Serial I/O.

2.4.1 Memory map

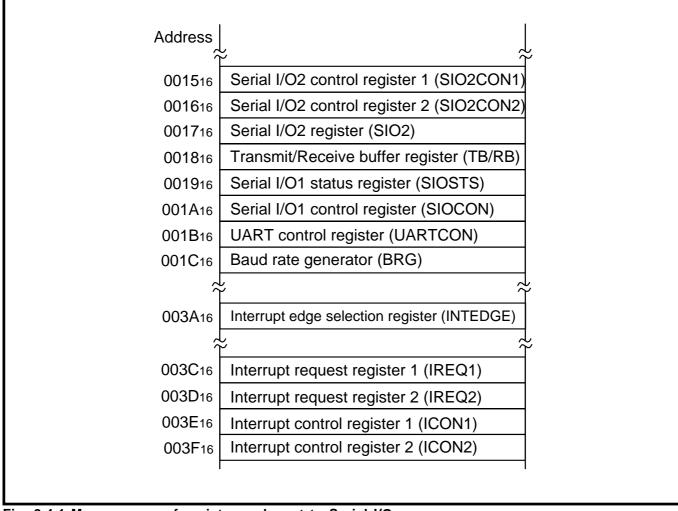


Fig. 2.4.1 Memory map of registers relevant to Serial I/O

2.4 Serial I/O

2.4.2 Relevant registers

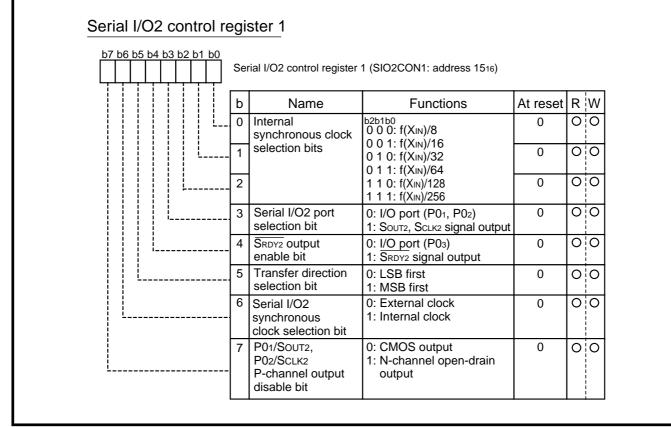


Fig. 2.4.2 Structure of Serial I/O2 control register 1

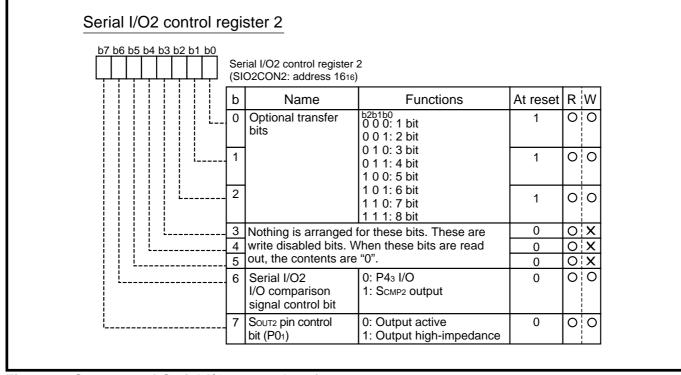


Fig. 2.4.3 Structure of Serial I/O2 control register 2

2.4 Serial I/O

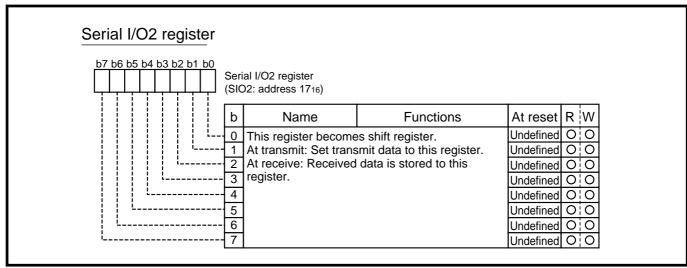


Fig. 2.4.4 Structure of Serial I/O2 register

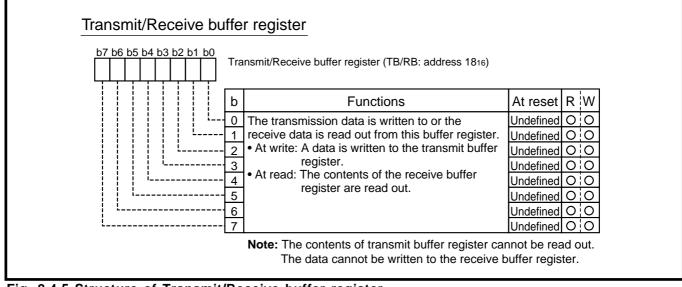


Fig. 2.4.5 Structure of Transmit/Receive buffer register

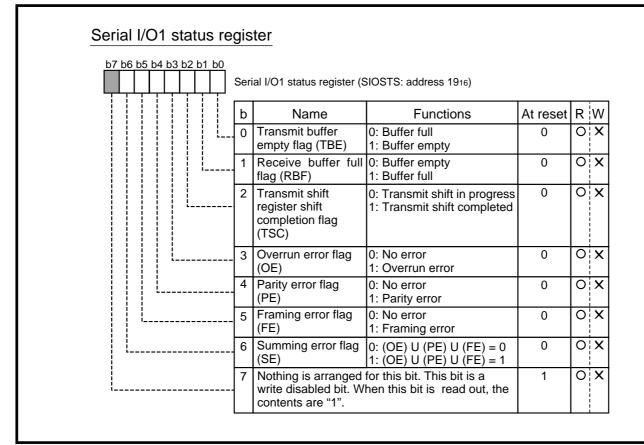


Fig. 2.4.6 Structure of Serial I/O1 status register

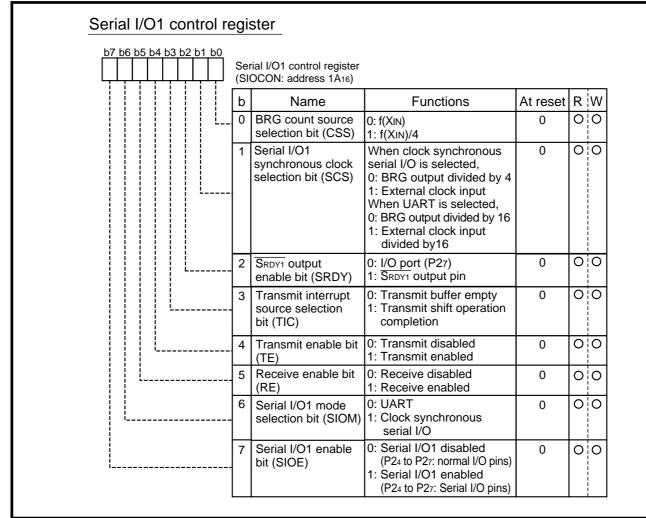
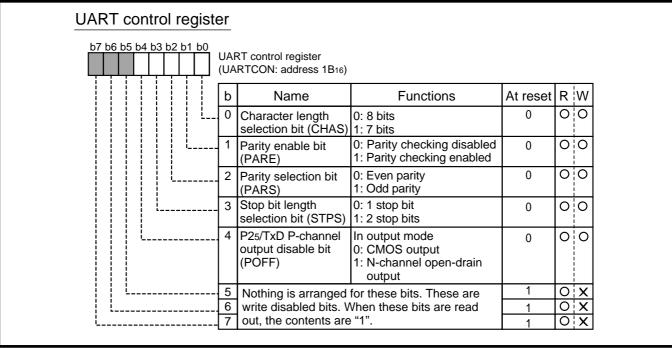
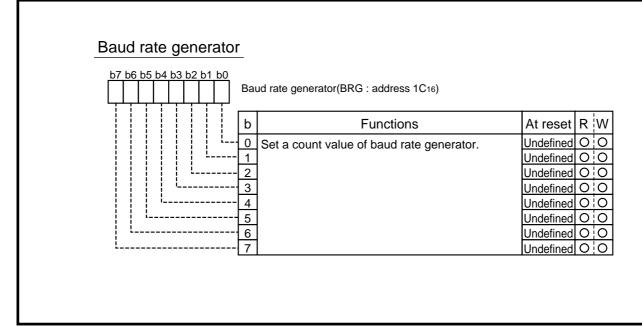
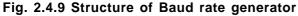


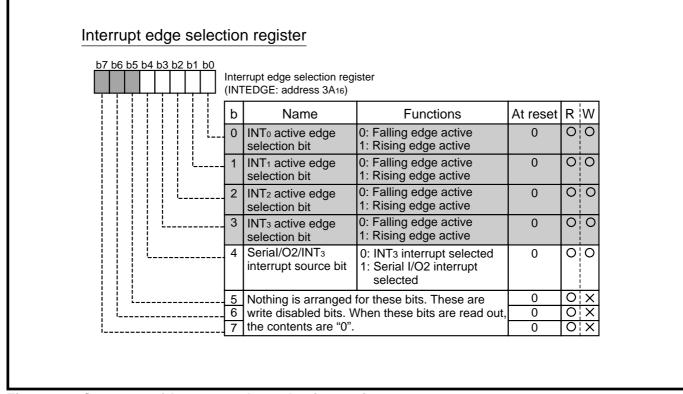
Fig. 2.4.7 Structure of Serial I/O1 control register	Fi
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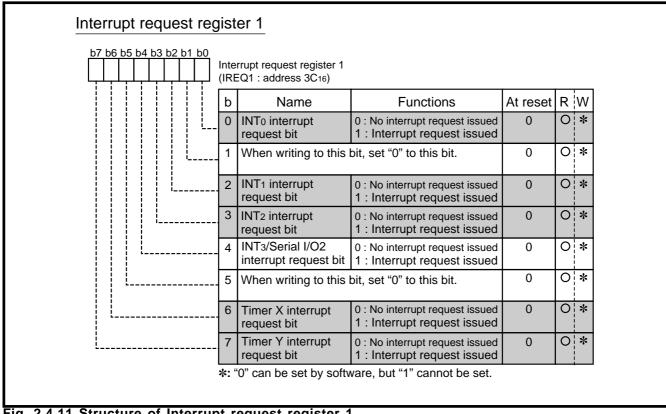


Fig. 2.4.11 Structure of Interrupt request register 1

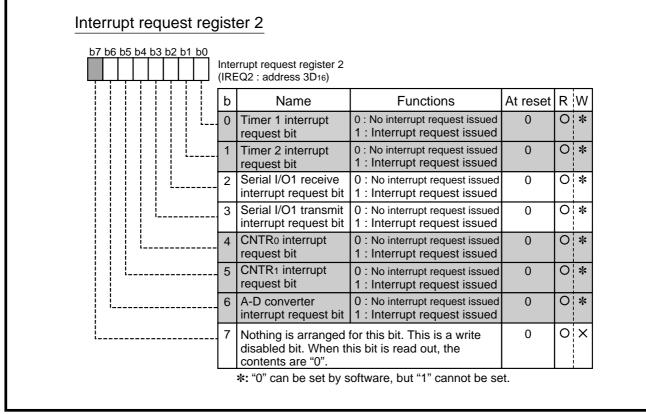


Fig. 2.4.12 Structure of Interrupt request register 2

b7 b6 b5 b4 b3 b2 b1 b0						
		rrupt control register 1 DN1 : address 3E16)				
	b	Name	Functions	At reset	R	W
	0	INT ₀ interrupt enable bit	0 : Interrupt disabled 1 : Interrupt enabled	0	0	0
	1	Fix this bit to "0".		0	0	0
	2	INT₁ interrupt enable bit	0 : Interrupt disabled 1 : Interrupt enabled	0	0	0
	. 3	INT ₂ nterrupt enable bit	0 : Interrupt disabled 1 : Interrupt enabled	0	0	0
	4	INT ₃ /Serial I/O2 interrupt enable bit	0 : Interrupt disabled 1 : Interrupt enabled	0	0	0
	5	Fix this bit to "0".		0	0	0
	6	Timer X interrupt enable bit	0 : Interrupt disabled 1 : Interrupt enabled	0	0	0
l	7	Timer Y interrupt enable bit	0 : Interrupt disabled 1 : Interrupt enabled	0	0	0

Fig. 2.4.13 Structure of Interrupt control register 1

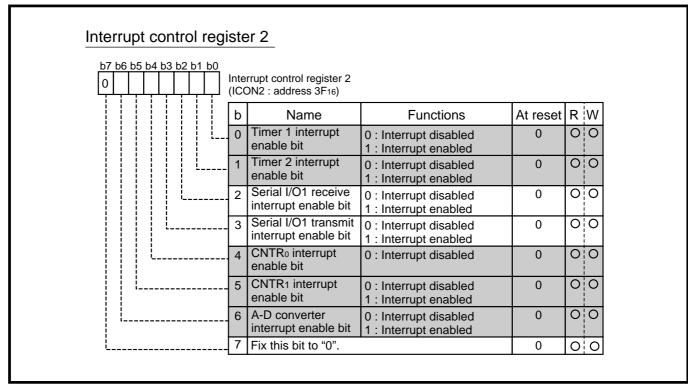


Fig. 2.4.14 Structure of Interrupt control register 2

2.4 Serial I/O

2.4.3 Serial I/O connection examples

(1) Control of peripheral IC equipped with CS pin

Figure 2.4.15 shows connection examples of a peripheral IC equipped with the CS pin. There are connection examples using a clock synchronous serial I/O mode.

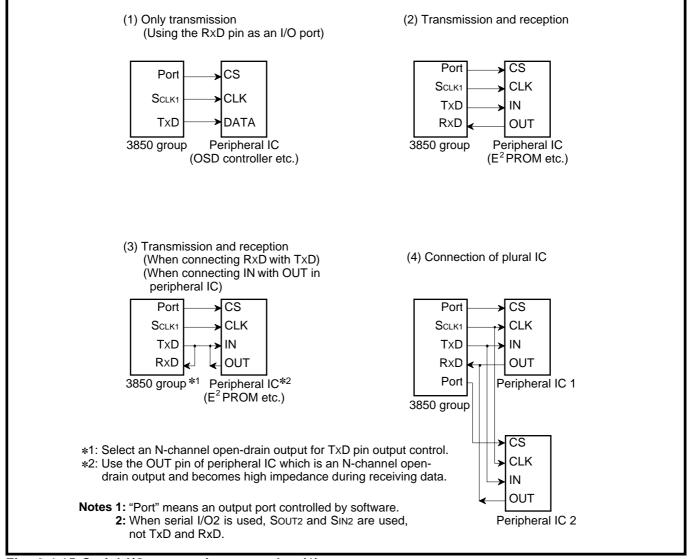
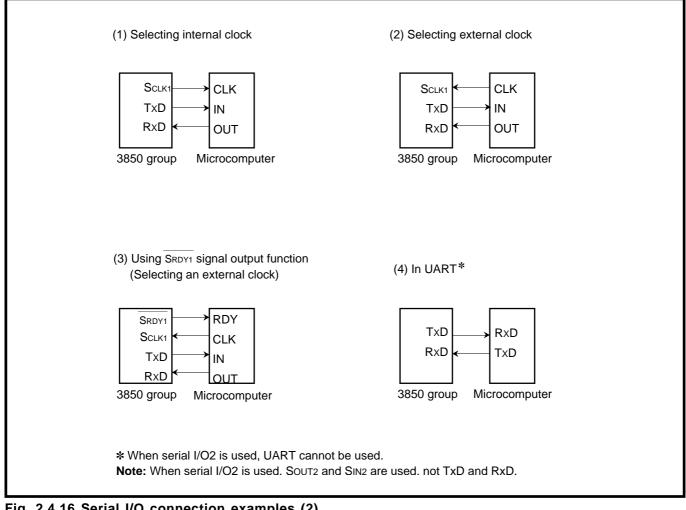


Fig. 2.4.15 Serial I/O connection examples (1)

2.4 Serial I/O

(2) Connection with microcomputer

Figure 2.4.16 shows connection examples with another microcomputer.





2.4 Serial I/O

2.4.4 Setting of serial I/O transfer data format

A clock synchronous or clock asynchronous (UART) can be selected as a data format of Serial I/O1. A clock synchronous is used as a data format of Serial I/O2.

Figure 2.4.17 shows the serial I/O transfer data format.

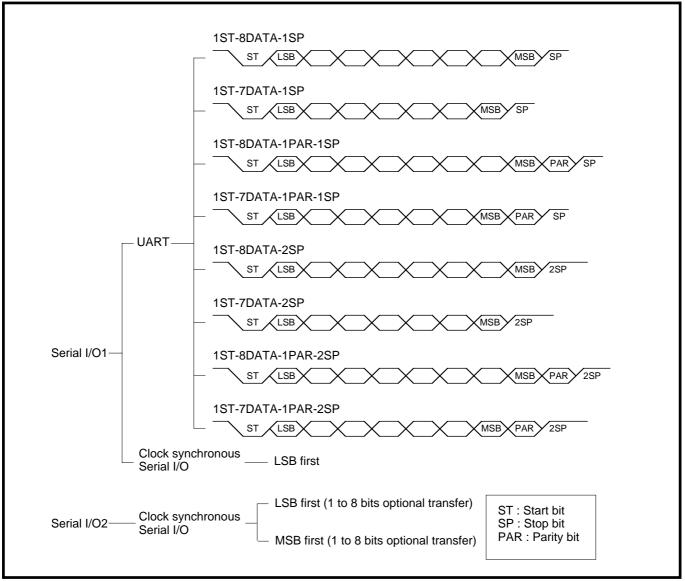


Fig. 2.4.17 Serial I/O transfer data format

2.4 Serial I/O

2.4.5 Serial I/O application examples

(1) Communication using clock synchronous serial I/O (transmit/receive)

Outline : 2-byte data is transmitted and received, using the clock synchronous serial I/O. The $\overline{S_{RDY1}}$ signal is used for communication control.

Figure 2.4.18 shows a connection diagram, and Figure 2.4.19 shows a timing chart. Figure 2.4.20 shows a registers setting relevant to the transmitting side, and Figure 2.4.21 shows registers setting relevant to the receiving side.

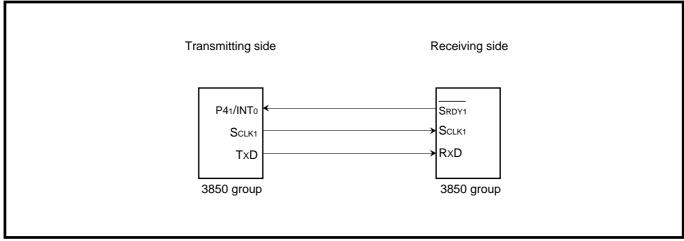


Fig. 2.4.18 Connection diagram

Specifications : • The Serial I/O is used (clock synchronous serial I/O is selected.)

- Synchronous clock frequency : 125 kHz ($f(X_{IN}) = 4$ MHz is divided by 32)
- The $\overline{S_{RDY1}}$ (receivable signal) is used.
- The receiving side outputs the $\overline{S_{RDY1}}$ signal at intervals of 2 ms (generated by timer), and 2-byte data is transferred from the transmitting side to the receiving side.

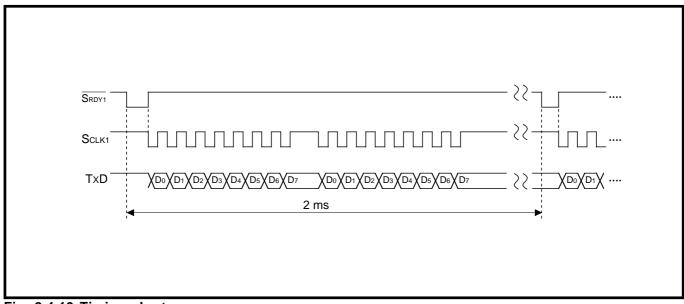
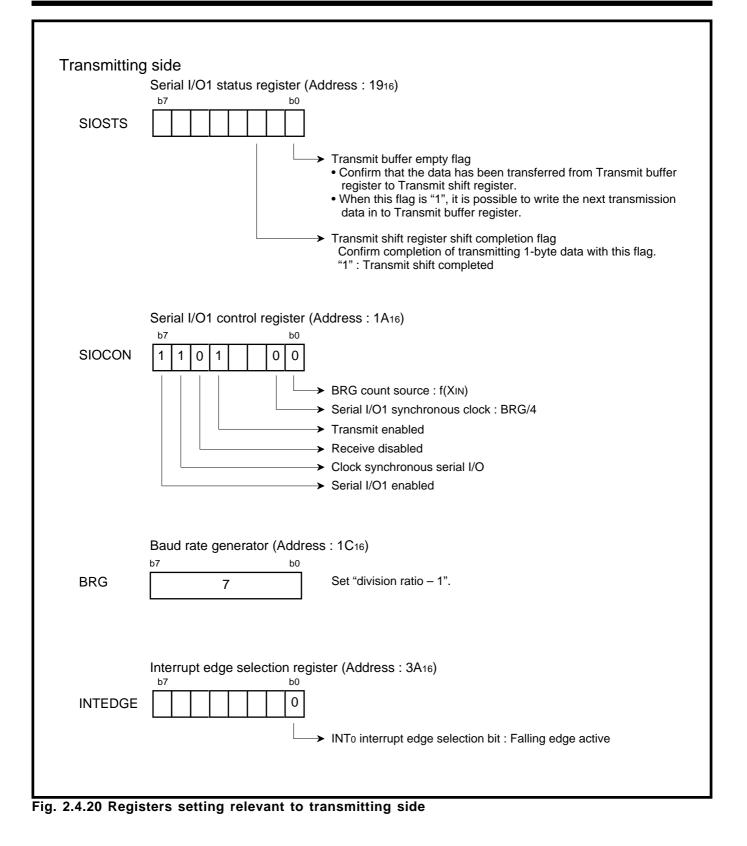


Fig. 2.4.19 Timing chart

2.4 Serial I/O



2.4 Serial I/O

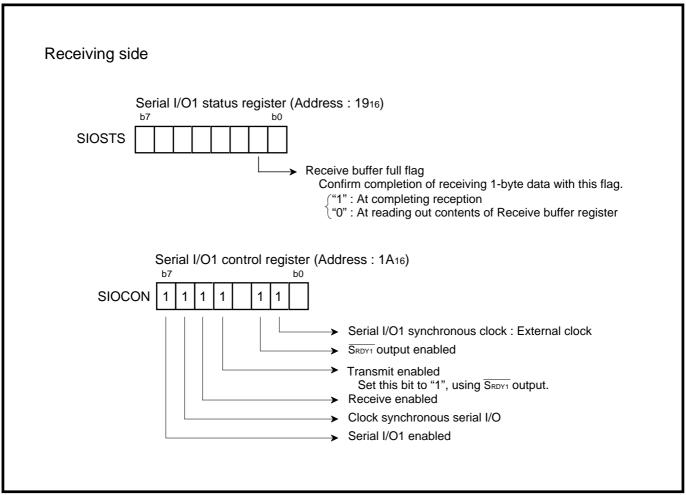


Fig. 2.4.21 Registers setting relevant to receiving side

2.4 Serial I/O

Figure 2.4.22 shows a control procedure of the transmitting side, and Figure 2.4.23 shows a control procedure of the receiving side.

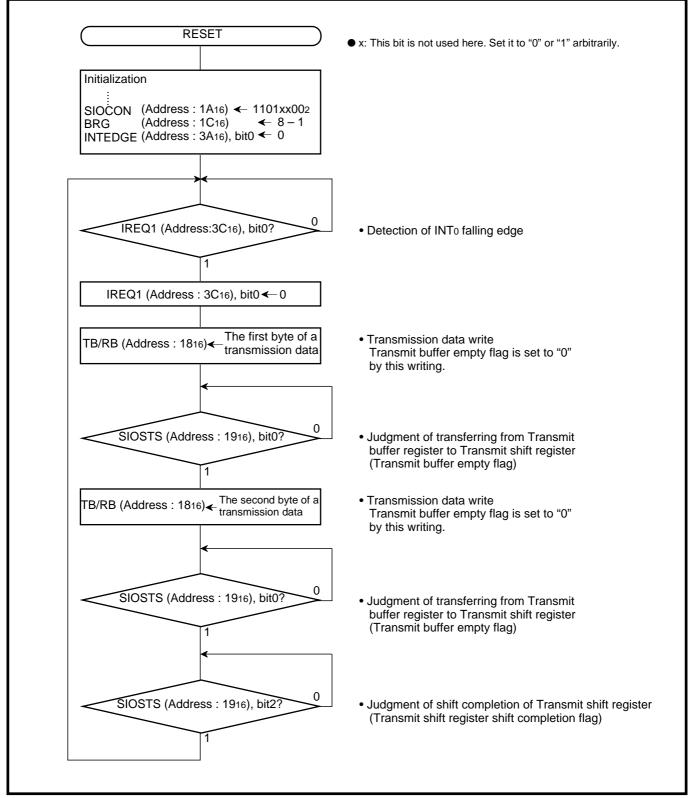


Fig. 2.4.22 Control procedure of transmitting side

2.4 Serial I/O

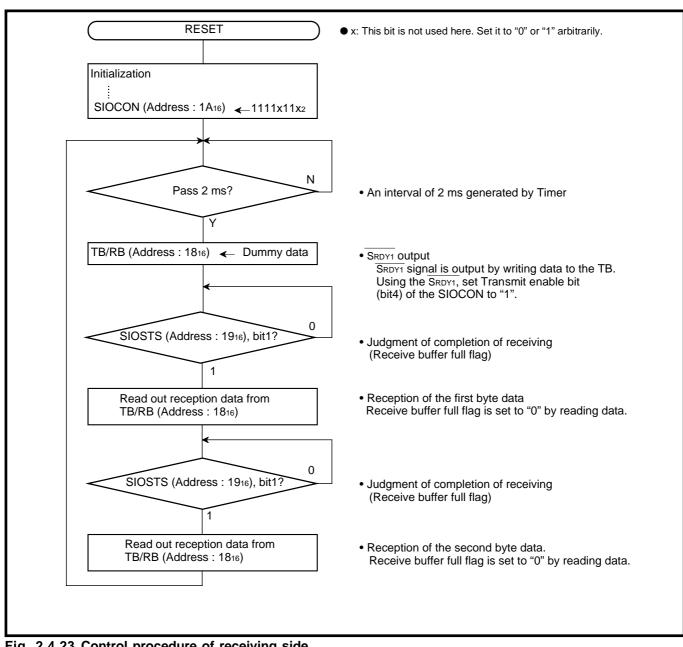


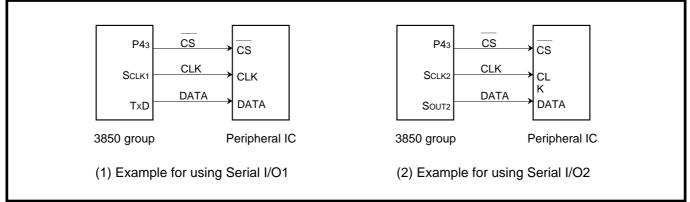
Fig. 2.4.23 Control procedure of receiving side

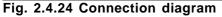
2.4 Serial I/O

(2) Output of serial data (control of peripheral IC)

Outline : 4-byte data is transmitted and received, using the clock synchronous serial I/O. The CS signal is output to a peripheral IC through port P4₃.

Figure 2.4.24 shows a connection diagram, and Figure 2.4.25 shows a timing chart.





Specifications : • The Serial I/O is used (clock synchronous serial I/O is selected.)

- Synchronous clock frequency : 125 kHz ($f(X_{IN}) = 4$ MHz is divided by 32)
- Transfer direction : LSB first
- The Serial I/O interrupt is not used.
- Port P4₃ is connected to the \overline{CS} pin ("L" active) of the peripheral IC for transmission control; the output level of port P4₃ is controlled by software.

CS	
CLK	
DATA	$X DO_0 X X DO_1 X DO_2 X X DO_3 X$
	Note: When serial I/O1 is used, the level of the last bit is held. When serial I/O2 is used, SOUT2 pin is in the high-impedance state after the transfer is completed.

Fig. 2.4.25 Timing chart (Serial I/O1)

2.4 Serial I/O

Figure 2.4.26 shows registers setting relevant to Serial I/O1, and Figure 2.4.27 shows a setting of serial I/O1 transmission data.

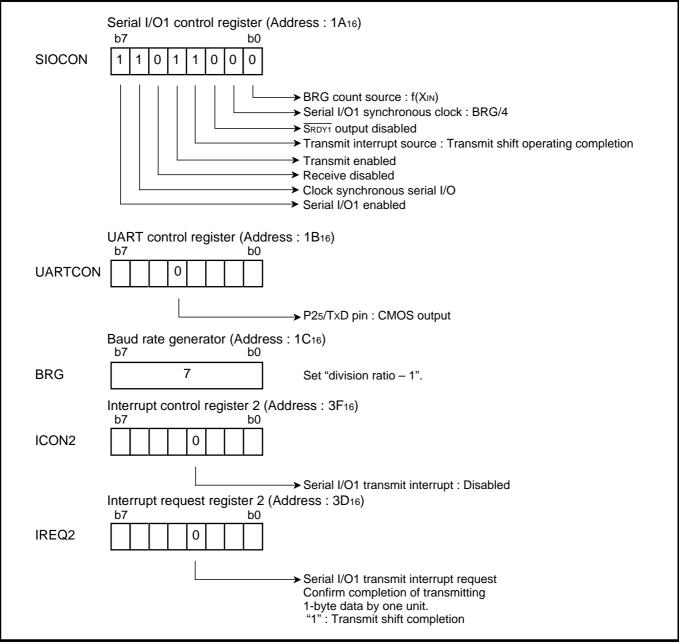
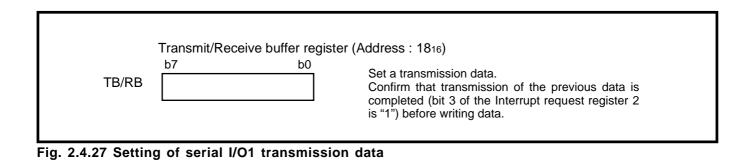


Fig. 2.4.26 Registers setting relevant to Serial I/O1



2.4 Serial I/O

Example for using Serial I/O1

When the registers are set as shown in Figure 2.4.26, the Serial I/O1 can transmit 1-byte data by writing data to the transmit buffer register.

Thus, after setting the \overline{CS} signal to "L", write the transmission data to the transmit buffer register by each 1 byte, and return the \overline{CS} signal to "H" when 4-byte data has been transmitted. Figure 2.4.28 shows a control procedure of Serial I/O1.

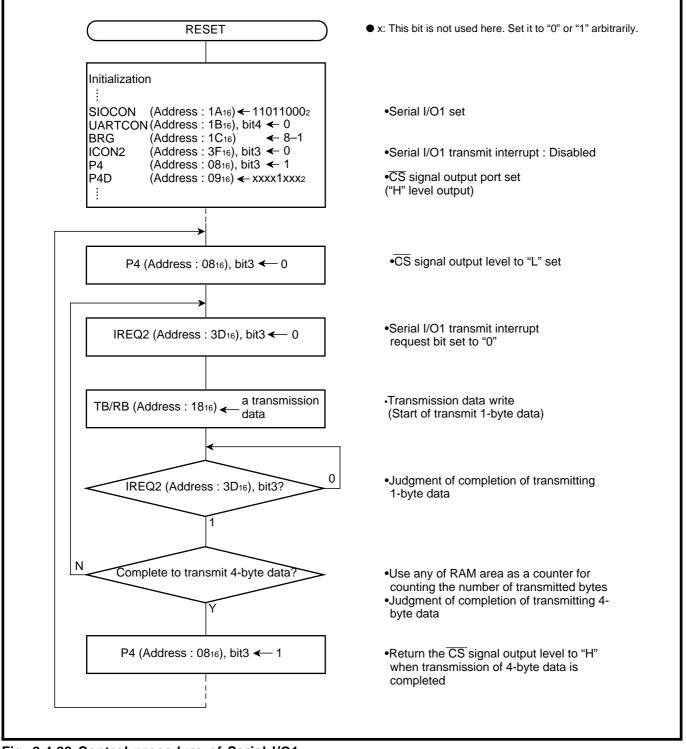
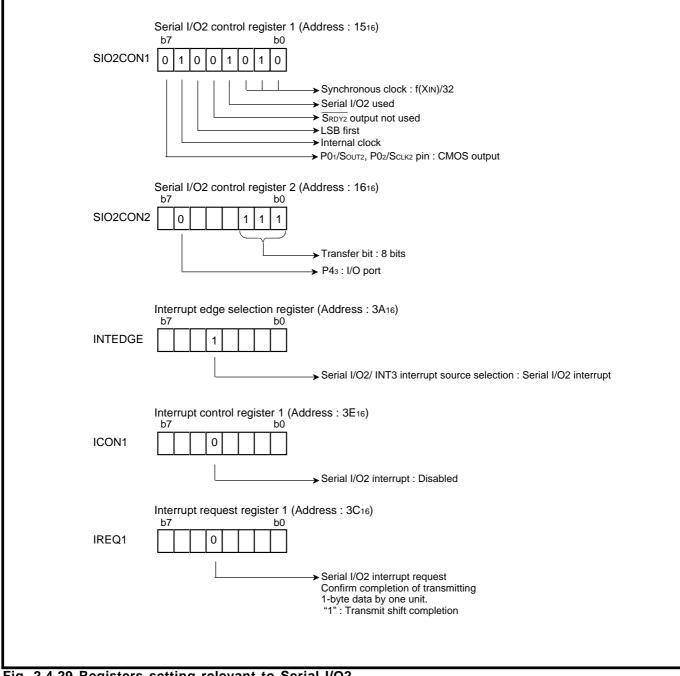
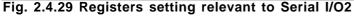


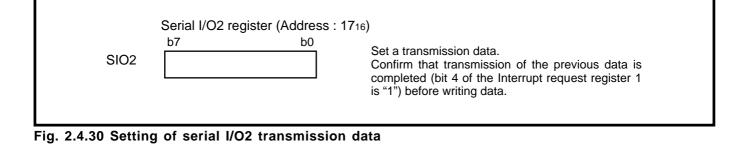
Fig. 2.4.28 Control procedure of Serial I/O1

2.4 Serial I/O

Figure 2.4.29 shows registers setting relevant to Serial I/O2, and Figure 2.4.30 shows a setting of serial I/O2 transmission data.







2.4 Serial I/O

Example for using Sreial I/O2

When the registers are set as shown in Fig. 2.4.29, the Serial I/O2 can transmit 1-byte data by writing data to the serial I/O2 register.

Thus, after setting the \overline{CS} signal to "L", write the transmission data to Serial I/O2 by each 1 byte, and return the \overline{CS} signal to "H" when 4-byte data has been transmitted. Figure 2.4.31 shows a control procedure of Serial I/O2.

 x: This bit is not used here. Set it to "0" or "1" arbitrarily. RESET Initialization Serial I/O2 control register set SIO2CON2(Address : 2616) - x0xxx1112 •Serial I/O2/INT3 interrupt source selection : Serial I/O2 interrupt •Serial I/O2 interrupt : Disabled ICON1 (Address : 3E16), bit4 ← 0 (Address : 0816), bit3 - 1 •CS signal output port set P4 P4D (Address : 0916) ← xxxx1xxx2 ("H" level output) ł •CS signal output level set to "L" P4 (Address : 08₁6), bit3 ← 0 Serial I/O2 interrupt request bit set to "0" IREQ1 (Address : 3C16), bit4 ← 0 •Transmission data write a transmission SIO2 (Address : 1716) 👞 (Start of transmit 1-byte data) data •Judgment of completion of transmitting 1-0 IREQ1 (Address : 3C16), bit4? byte data 1 •Use any of RAM area as a counter for counting the number of transmitted bytes Ν Complete to transmit 4-byte data? •Judgment of completion of transmitting 4byte data •Return the CS signal output level to "H" P4 (Address : 0816), bit3 - 1 when transmission of 4-byte data is completed

Fig. 2.4.31 Control procedure of Serial I/O2

(3) Cyclic transmission or reception of block data (data of specified number of bytes) between two microcomputers

Outline : When the clock synchronous serial I/O is used for communication, synchronization of the clock and the data between the transmitting and receiving sides may be lost because of noise included in the synchronous clock. It is necessary to correct that constantly, using "heading adjustment".

This "heading adjustment" is carried out by using the interval between blocks in this example.

Figure 2.4.32 shows a connection diagram.

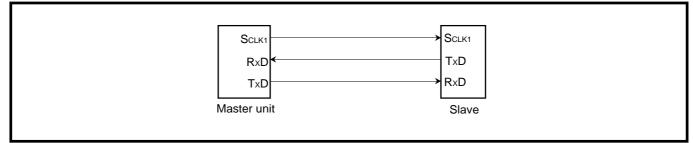


Fig. 2.4.32 Connection diagram

Specifications :

- The serial I/O is used (clock synchronous serial I/O is selected).
- Synchronous clock frequency : 131 kHz (f(X_{IN}) = 4.19 MHz is divided by 32)
- \bullet Byte cycle: 488 μs
- Number of bytes for transmission or reception : 8 byte/block
- Block transfer cycle : 16 ms
- Block transfer term : 3.5 ms
- Interval between blocks : 12.5 ms
- Heading adjustment time : 8 ms

Limitations of specifications :

- Reading of the reception data and setting of the next transmission data must be completed within the time obtained from "byte cycle – time for transferring 1-byte data" (in this example, the time taken from generating of the serial I/O1 receive interrupt request to input of the next synchronous clock is 431 μs).
- "Heading adjustment time < interval between blocks" must be satisfied.

2.4 Serial I/O

The communication is performed according to the timing shown in Figure 2.4.33. In the slave unit, when a synchronous clock is not input within a certain time (heading adjustment time), the next clock input is processed as the beginning (heading) of a block.

When a clock is input again after one block (8 byte) is received, the clock is ignored. Figure 2.4.34 shows relevant registers setting.

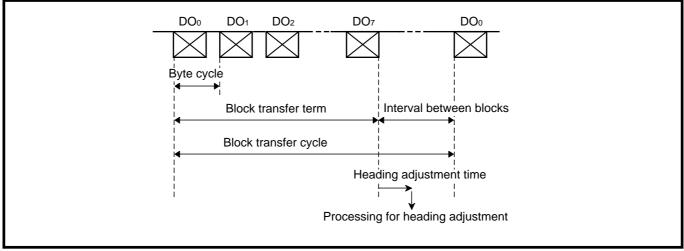


Fig. 2.4.33 Timing chart

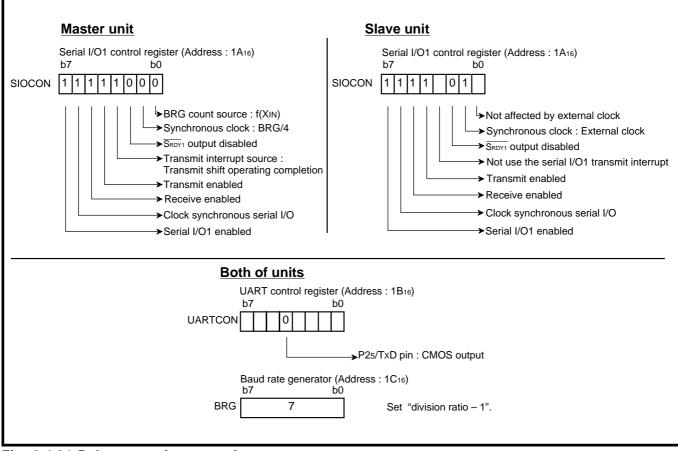


Fig. 2.4.34 Relevant registers setting

2.4 Serial I/O

Control procedure :

Control in the master unit

After setting the relevant registers shown in Figure 2.4.34, the master unit starts transmission or reception of 1-byte data by writing transmission data to the transmit buffer register.

To perform the communication in the timing shown in Figure 2.4.33, take the timing into account and write transmission data. Additionally, read out the reception data when the serial I/O1 transmit interrupt request bit is set to "1", or before the next transmission data is written to the transmit buffer register.

Figure 2.4.35 shows a control procedure of the master unit using timer interrupts.

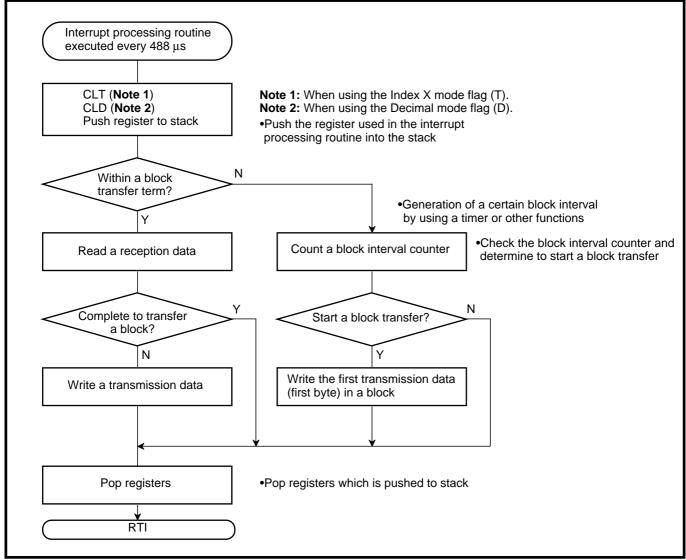


Fig. 2.4.35 Control procedure of master unit

2.4 Serial I/O

• Control in the slave unit

After setting the relevant registers as shown in Figure 2.4.34, the slave unit becomes the state where a synchronous clock can be received at any time, and the serial I/O receive interrupt request bit is set to "1" each time an 8-bit synchronous clock is received.

In the serial I/O receive interrupt processing routine, the data to be transmitted next is written to the transmit buffer register after the received data is read out.

However, if no serial I/O receive interrupt occurs for a certain time (heading adjustment time or more), the following processing will be performed.

1. The first 1-byte data of the transmission data in the block is written into the transmit buffer register.

2. The data to be received next is processed as the first 1 byte of the received data in the block. Figure 2.4.36 shows a control procedure of the slave unit using the serial I/O receive interrupt and any timer interrupt (for heading adjustment).

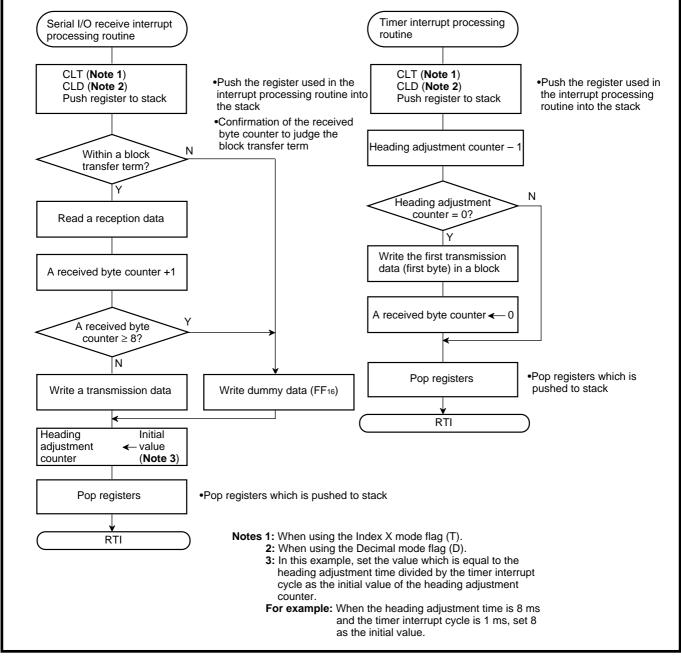


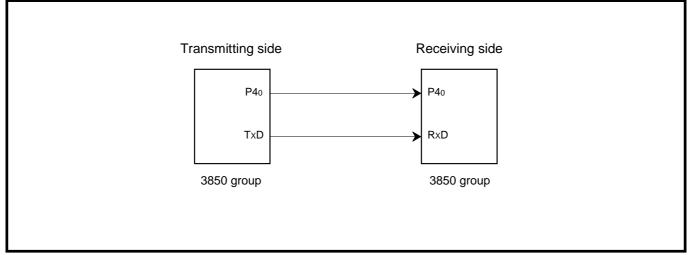
Fig. 2.4.36 Control procedure of slave unit

2.4 Serial I/O

(4) Communication (transmit/receive) using asynchronous serial I/O (UART)

Outline : 2-byte data is transmitted and received, using the asynchronous serial I/O. Port P4₀ is used for communication control.

Figure 2.4.37 shows a connection diagram, and Figure 2.4.38 shows a timing chart.





Specifications : • The Serial I/O1 is used (UART is selected).

- Transfer bit rate : 9600 bps ($f(X_{IN}) = 4.9152$ MHz is divided by 512)
- Communication control using port P40
- (The output level of port P4₀ is controlled by software.)
- 2-byte data is transferred from the transmitting side to the receiving side at intervals of 10 ms generated by the timer.

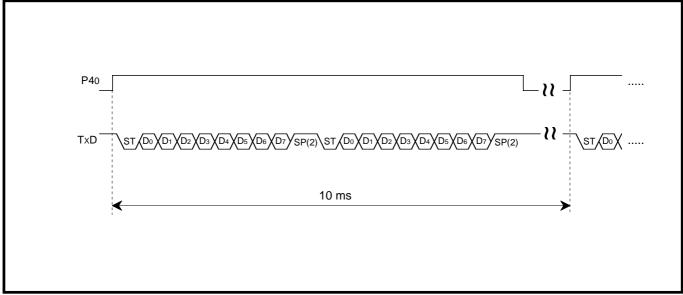


Fig. 2.4.38 Timing chart (using UART)

2.4 Serial I/O

Table 2.4.1 and Table 2.4.2 show setting examples of the baud rate generator (BRG) values and transfer bit rate values; Figure 2.4.39 shows registers setting relevant to the transmitting side; Figure 2.4.40 shows registers setting relevant to the receiving side.

BRG count source		Transfer bit rate	e (bps) (Note 2)
(Note 1)	BRG setting value	at f(XIN) = 4.9152 MHz	at f(XIN) = 8 MHz
f(XIN)/4	255(FF16)	300	488.28125
f(XIN)/4	127(7F16)	600	976.5625
f(XIN)/4	63(3F16)	1200	1953.125
f(XIN)/4	31(1F16)	2400	3906.25
f(XIN)/4	15(0F16)	4800	7812.5
f(XIN)/4	7(0716)	9600	15625
f(XIN)/4	3(0316)	19200	31250
f(XIN)/4	1(0116)	38400	62500
f(XIN)	3(0316)	76800	125000
f(XIN)	1(0116)	153600	250000
f(XIN)	0(0016)	307200	500000

Table 2.4.1 S	Setting exam	ples of Baud	rate generator	values and	transfer bit rate	values (1)
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Table 2.4.2 Setting examples of Baud rate generator values and transfer bit rate values (2)

BRG count source	BBC patting value	Transfer bit rate (bps) (Note 2)		
(Note 1)	BRG setting value	at f(XIN) = 7.9872 MHz		
f(XIN)/4	207(CF16)	600		
f(XIN)/4	103(6716)	1200		
f(XIN)/4	51(3316)	2400		
f(XIN)/4	25(1916)	4800		
f(XIN)/4	12(0C16)	9600		
f(XIN)	25(1916)	19200		
f(XIN)	12(0C16)	38400		

Notes 1: Select the BRG count source with bit 0 of the serial I/O1 control register (Address : 1A16).

2: Equation of transfer bit rate:

Transfer bit rate (bps) =
$$\frac{f(XIN)}{(BRG \text{ setting value + 1}) \times 16 \times \text{m}^*}$$

*m: When bit 0 of the serial I/O1 control register (Address : 1A16) is set to "0", a value of m is 1.

When bit 0 of the serial I/O1 control register (Address : 1A16) is set to "1", a value of m is 4.

2.4 Serial I/O

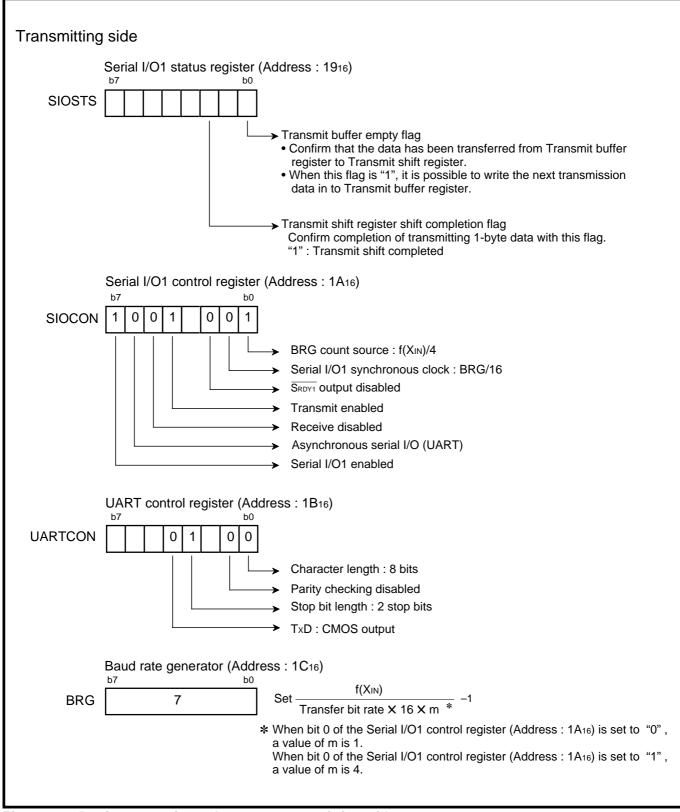
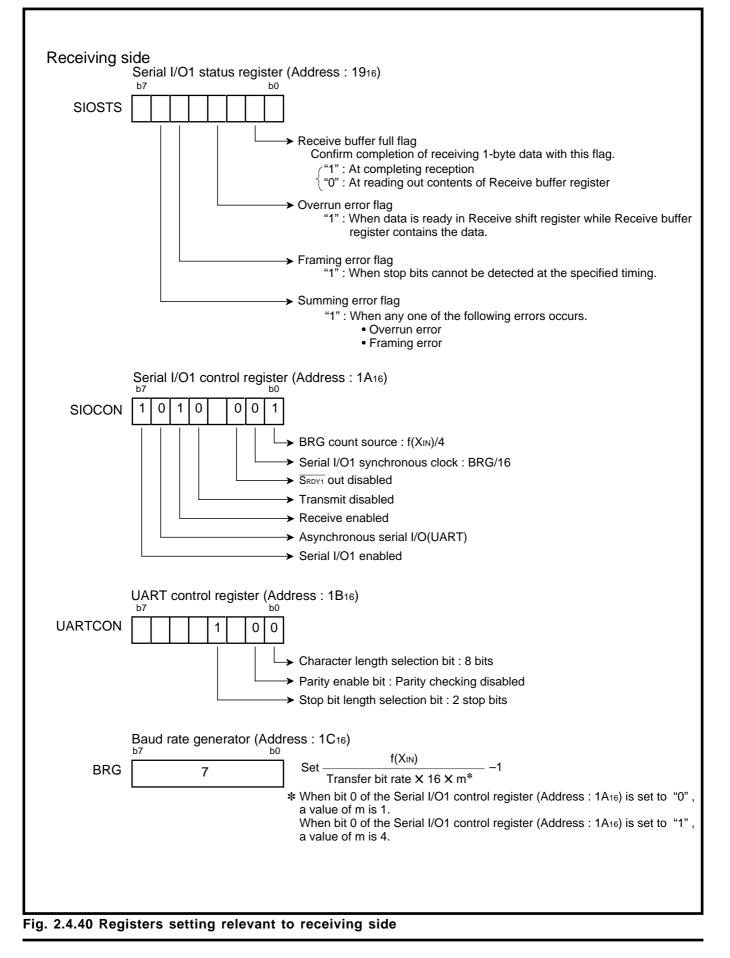


Fig. 2.4.39 Registers setting relevant to transmitting side

2.4 Serial I/O



2.4 Serial I/O

Figure 2.4.41 shows a control procedure of the transmitting side, and Figure 2.4.42 shows a control procedure of the receiving side.

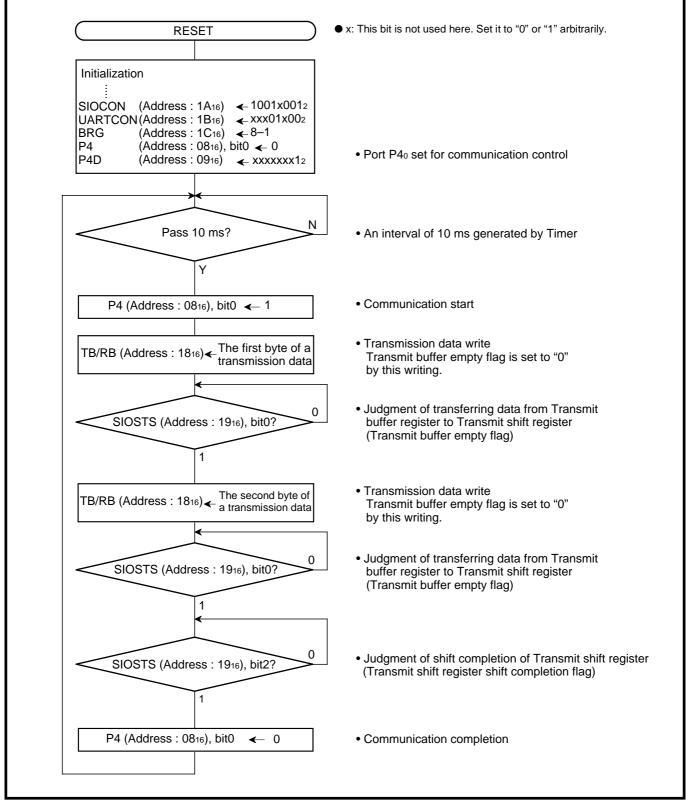


Fig. 2.4.41 Control procedure of transmitting side

2.4 Serial I/O

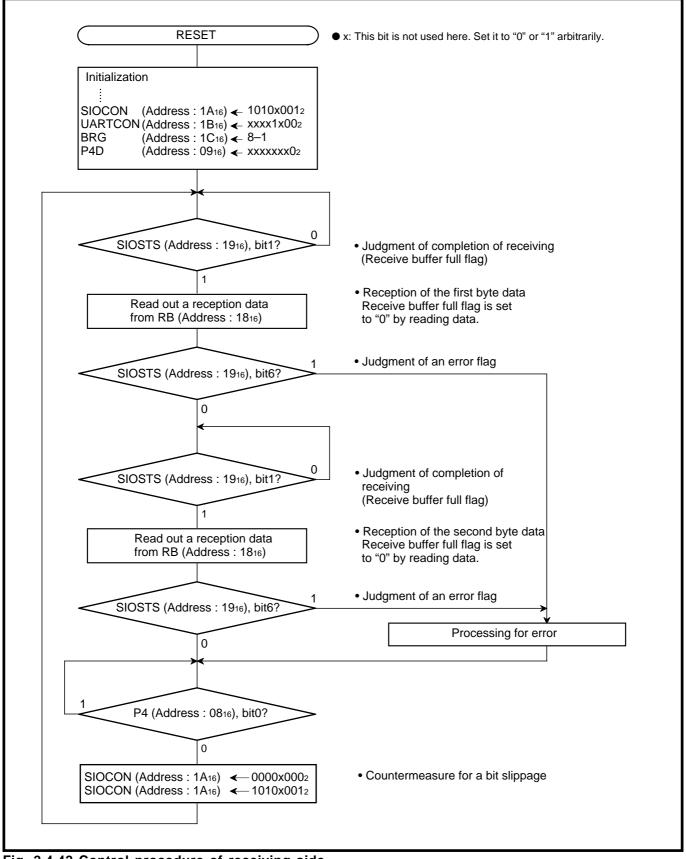


Fig. 2.4.42 Control procedure of receiving side

2.4 Serial I/O

2.4.6 Notes on serial I/O

(1) Notes when selecting clock synchronous serial I/O (Serial I/O1)

① Stop of transmission operation

Clear the serial I/O1 enable bit and the transmit enable bit to "0" (Serial I/O1 and transmit disabled).

Reason

Since transmission is not stopped and the transmission circuit is not initialized even if only the serial I/O1 enable bit is cleared to "0" (Serial I/O1 disabled), the internal transmission is running (in this case, since pins TxD, RxD, SCLK1, and SRDY1 function as I/O ports, the transmission data is not output). When data is written to the transmit buffer register in this state, data starts to be shifted to the transmit shift register. When the serial I/O1 enable bit is set to "1" at this time, the data during internally shifting is output to the TxD pin and an operation failure occurs.

② Stop of receive operation

Clear the receive enable bit to "0" (receive disabled), or clear the serial I/O1 enable bit to "0" (Serial I/O1 disabled).

③ Stop of transmit/receive operation

Clear the transmit enable bit and receive enable bit to "0" simultaneously (transmit and receive disabled).

(when data is transmitted and received in the clock synchronous serial I/O mode, any one of data transmission and reception cannot be stopped.)

Reason

In the clock synchronous serial I/O mode, the same clock is used for transmission and reception. If any one of transmission and reception is disabled, a bit error occurs because transmission and reception cannot be synchronized.

In this mode, the clock circuit of the transmission circuit also operates for data reception. Accordingly, the transmission circuit does not stop by clearing only the transmit enable bit to "0" (transmit disabled). Also, the transmission circuit is not initialized by clearing the serial I/O1 enable bit to "0" (Serial I/O1 disabled) (refer to (1) ①).

2.4 Serial I/O

(2) Notes when selecting clock asynchronous serial I/O (Serial I/O1)

① Stop of transmission operation

Clear the transmit enable bit to "0" (transmit disabled).

Reason

Since transmission is not stopped and the transmission circuit is not initialized even if only the serial I/O1 enable bit is cleared to "0" (Serial I/O1 disabled), the internal transmission is running (in this case, since pins TxD, RxD, SCLK1, and SRDY1 function as I/O ports, the transmission data is not output). When data is written to the transmit buffer register in this state, data starts to be shifted to the transmit shift register. When the serial I/O1 enable bit is set to "1" at this time, the data during internally shifting is output to the TxD pin and an operation failure occurs.

② Stop of receive operation

Clear the receive enable bit to "0" (receive disabled).

③ Stop of transmit/receive operation

Only transmission operation is stopped.

Clear the transmit enable bit to "0" (transmit disabled).

Reason

Since transmission is not stopped and the transmission circuit is not initialized even if only the serial I/O1 enable bit is cleared to "0" (Serial I/O1 disabled), the internal transmission is running (in this case, since pins TxD, RxD, SCLK1, and SRDY1 function as I/O ports, the transmission data is not output). When data is written to the transmit buffer register in this state, data starts to be shifted to the transmit shift register. When the serial I/O1 enable bit is set to "1" at this time, the data during internally shifting is output to the TxD pin and an operation failure occurs.

Only receive operation is stopped.

Clear the receive enable bit to "0" (receive disabled).

(3) SRDY1 output of reception side

When signals are output from the $\overline{SRDY1}$ pin on the reception side by using an external clock in the clock synchronous serial I/O mode, set all of the receive enable bit, the $\overline{SRDY1}$ output enable bit, and the transmit enable bit to "1" (transmit enabled).

(4) Setting serial I/O1 control register again (Serial I/O1)

Set the serial I/O1 control register again after the transmission and the reception circuits are reset by clearing both the transmit enable bit and the receive enable bit to "0".

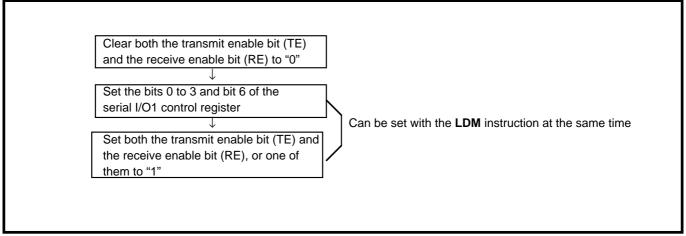


Fig. 2.4.43 Sequence of setting serial I/O1 control register again

2.4 Serial I/O

(5) Data transmission control with referring to transmit shift register completion flag (Serial I/O1) The transmit shift register completion flag changes from "1" to "0" with a delay of 0.5 to 1.5 shift clocks. When data transmission is controlled with referring to the flag after writing the data to the transmit buffer register, note the delay.

(6) Transmission control when external clock is selected (Serial I/O1)

When an external clock is used as the synchronous clock for data transmission, set the transmit enable bit to "1" at "H" of the SCLK1 input level. Also, write the transmit data to the transmit buffer register at "H" of the SCLK1 input level.

(7) Transmit interrupt request when transmit enable bit is set (Serial I/O1)

When the transmit interrupt is used, set the transmit interrupt enable bit to transmit enabled as shown in the following sequence.

- ① Set the interrupt enable bit to "0" (disabled) with CLB instruction.
- 2 Prepare serial I/O for transmission/reception.
- ③ Set the interrupt request bit to "0" with CLB instruction after 1 or more instruction has been executed.
- ④ Set the interrupt enable bit to "1" (enabled).

Reason

When the transmission enable bit is set to "1", the transmit buffer empty flag and transmit shift register completion flag are set to "1". The interrupt request is generated and the transmission interrupt bit is set regardless of which of the two timings listed below is selected as the timing for the transmission interrupt to be generated.

- Transmit buffer empty flag is set to "1"
- Transmit shift register completion flag is set to "1"

(8) Transmit data writing (Serial I/O2)

In the clock synchronous serial I/O, when selecting an external clock as synchronous clock, write the transmit data to the serial I/O2 register (serial I/O shift register) at "H" of the transfer clock input level.

2.5 PWM

2.5 PWM

This paragraph explains the registers setting method and the notes relevant to the PWM.

2.5.1 Memory map

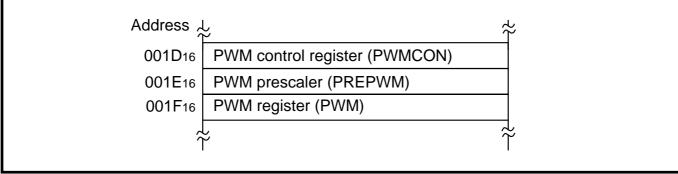


Fig. 2.5.1 Memory map of registers relevant to PWM

2.5.2 Relevant registers

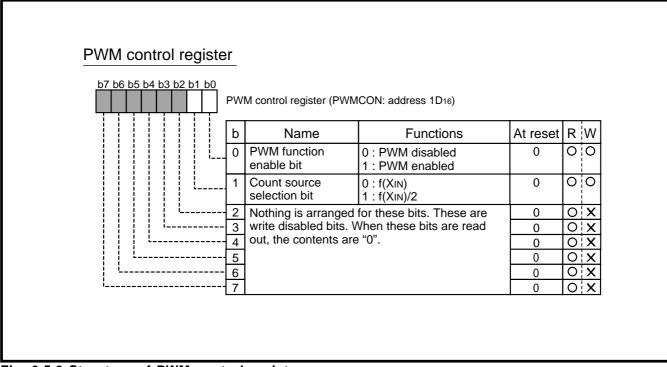


Fig. 2.5.2 Structure of PWM control register

2.5 PWM

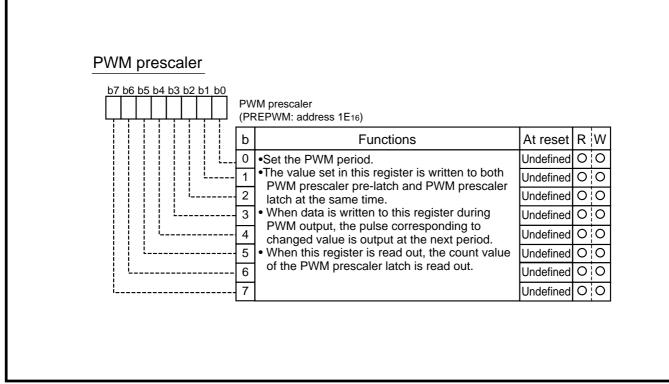
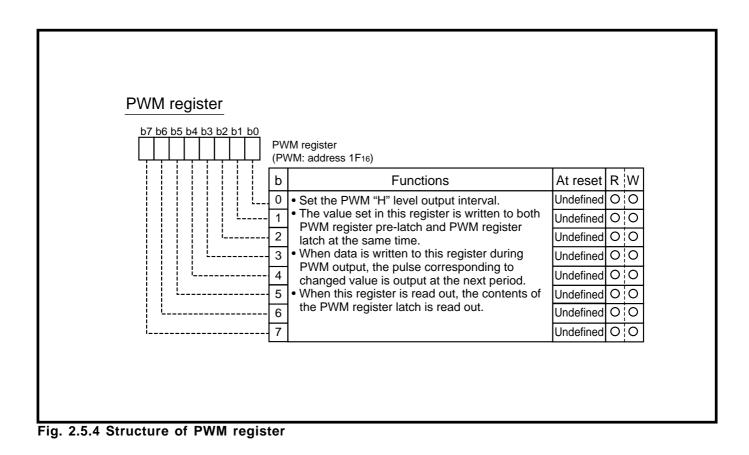


Fig. 2.5.3 Structure of PWM prescaler



2.5 PWM

2.5.3 PWM output circuit application example

<Motor control>

Outline : The rotation speed of the motor is controlled by using PWM (pulse width modulation) output.

Figure 2.5.5 shows a connection diagram ; Figures 2.5.6 shows PWM output timing, and Figure 2.5.7 shows a setting of the related registers.

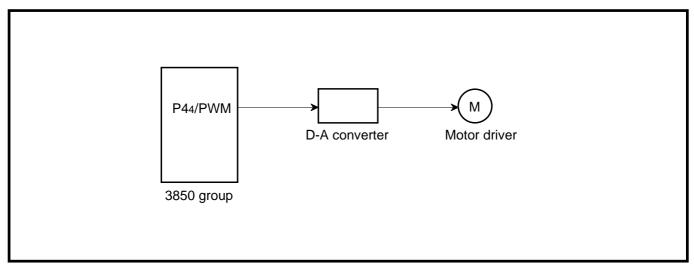


Fig. 2.5.5 Connection diagram

Specifications : • Motor is controlled by using the PWM output function of 8-bit resolution.

- Clock f(XIN) = 5.0 MHz
- "T", PWM cycle : 102 μs
- "t", "H" level width of output pulse : 40 μ s (Fixed speed)

* A motor speed can be changed by modifying the "H" level width of output pulse.

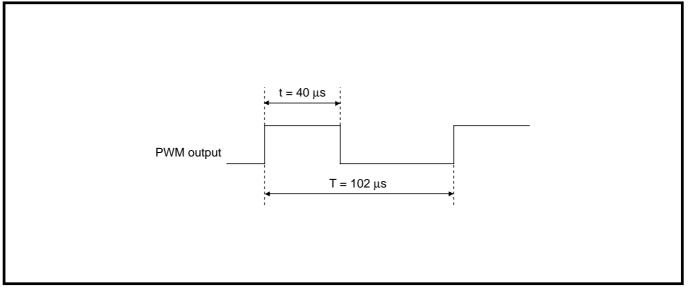


Fig. 2.5.6 PWM output timing

2.5 PWM

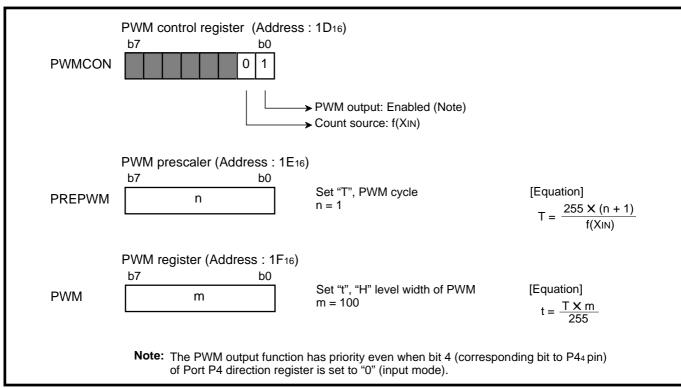


Fig. 2.5.7 Setting of relevant registers

<About PWM output>

1. Set the PWM function enable bit to "1" : The P44/PWM pin is used as the PWM pin.

The pulse beginning with "H" level pulse is output.

2. Set the PWM function enable bit to "0" : The P44/PWM pin is used as the port P44.

Thus, when fixing the output level, take the following procedure:

- (1) Write an output value to bit 6 of the port P4 register.
- (2) Write "000100002" to the port P4 direction register.
- 3. After data is set to the PWM prescaler and the PWM register, the PWM waveforms corresponding to updated data will be output from the next repetitive cycle.

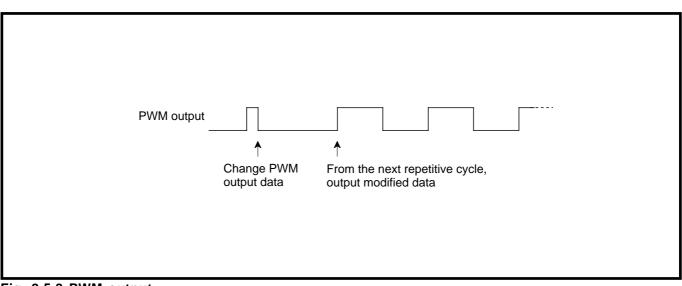


Fig. 2.5.8 PWM output

2.5 PWM

Control procedure : By setting the related registers as shown by Figure 2.5.7, PWM waveforms are output to the externals. This PWM output is integrated through the low pass filter, and that converted into DC signals is used for control of the motor.

Figure 2.5.9 shows control procedure.

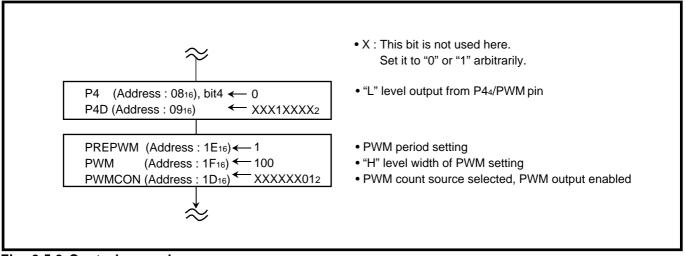


Fig. 2.5.9 Control procedure

2.5.4 Notes on PWM

The PWM starts after the PWM enable bit is set to enable and "L" level is output from the PWM pin. The length of this "L" level output is as follows:

 $\frac{n+1}{2 \cdot f(X_{IN})}$ sec. (Count source selection bit = 0, where n is the value set in the prescaler) $\frac{n+1}{2 \cdot f(X_{IN})}$ sec. (Count source selection bit = 1, where n is the value set in the prescaler)

2.6 A-D converter

This paragraph explains the registers setting method and the notes relevant to the A-D converter.

2.6.1 Memory map

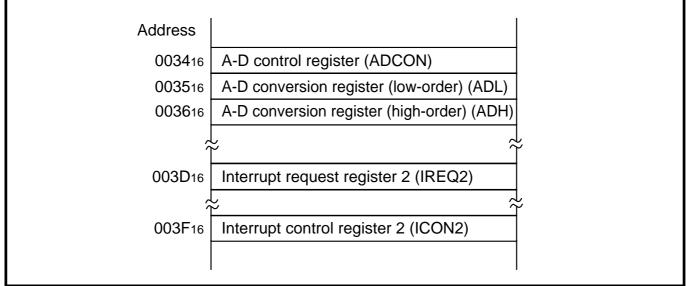
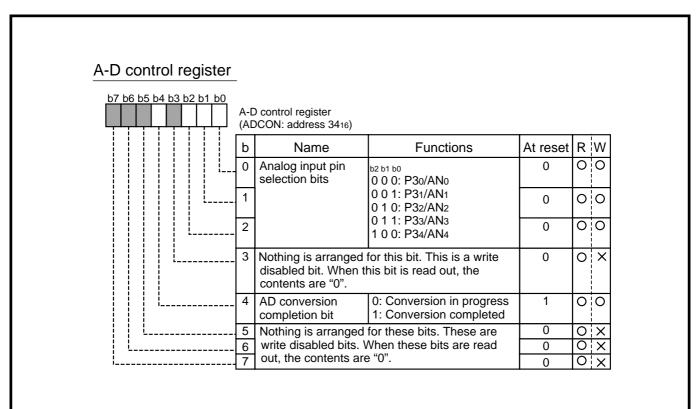
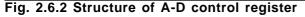


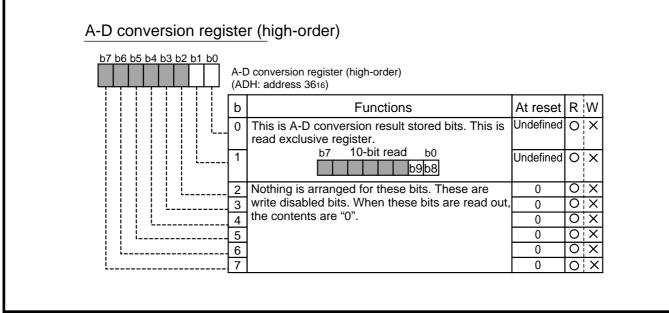
Fig. 2.6.1 Memory map of registers relevant to A-D converter

2.6.2 Relevant registers

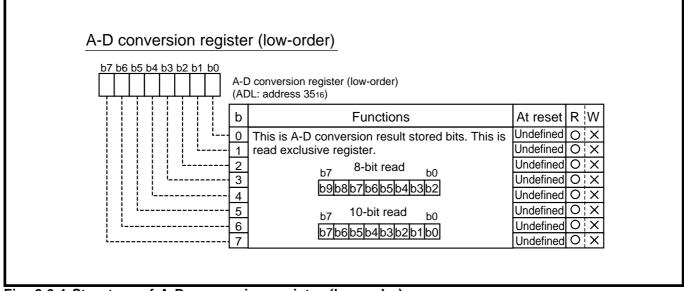




2.6 A-D converter









2.6 A-D converter

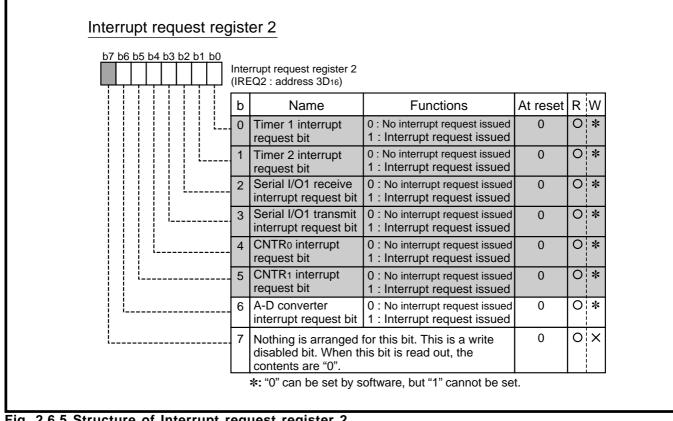
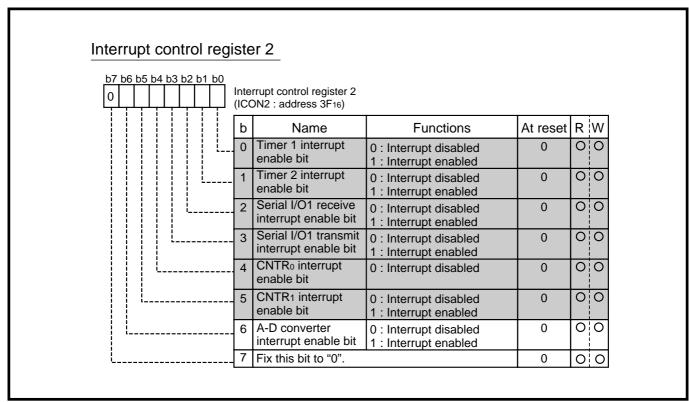


Fig. 2.6.5 Structure of Interrupt request register 2





2.6 A-D converter

2.6.3 A-D converter application examples

(1) Conversion of analog input voltage

Outline : The analog input voltage input from a sensor is converted to digital values.

Figure 2.6.7 shows a connection diagram, and Figure 2.6.8 shows the relevant registers setting.

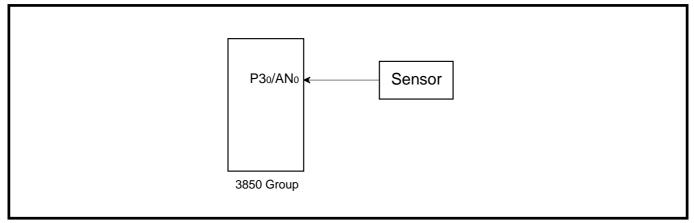


Fig. 2.6.7 Connection diagram

Specifications : •The analog input voltage input from a sensor is converted to digital values. •P3₀/AN₀ pin is used as an analog input pin.

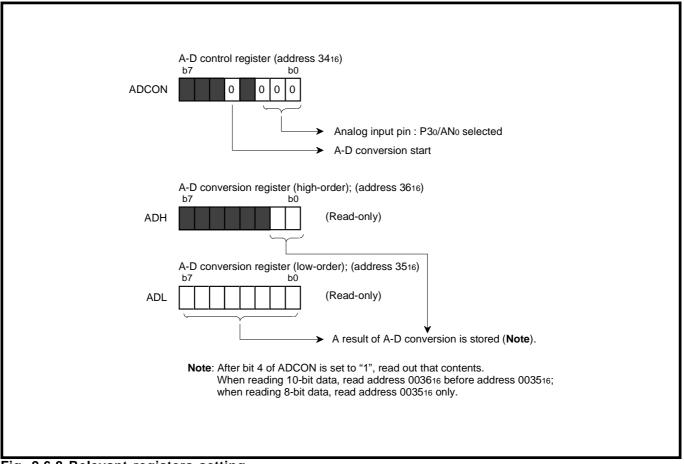


Fig. 2.6.8 Relevant registers setting

2.6 A-D converter

An analog input signal from a sensor is converted to the digital value according to the relevant registers setting shown by Figure 2.6.8. Figure 2.6.9 shows the control procedure for 8-bit read, and Figure 2.6.10 shows the control procedure for 10-bit read.

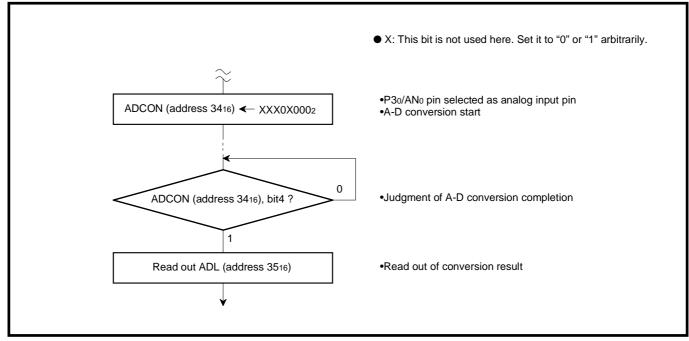


Fig. 2.6.9 Control procedure for 8-bit read

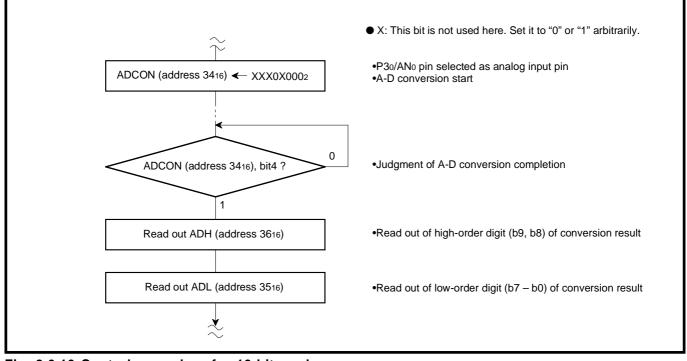


Fig. 2.6.10 Control procedure for 10-bit read

2.6 A-D converter

2.6.4 Notes on A-D converter

(1) Analog input pin

Make the signal source impedance for analog input low, or equip an analog input pin with an external capacitor of 0.01 μ F to 1 μ F. Further, be sure to verify the operation of application products on the user side.

Reason

An analog input pin includes the capacitor for analog voltage comparison. Accordingly, when signals from signal source with high impedance are input to an analog input pin, charge and discharge noise generates. This may cause the A-D conversion precision to be worse.

(2) A-D converter power source pin

The AVss pin is A-D converter power source pin. Regardless of using the A-D conversion function or not, connect it as following :

• AVss : Connect to the Vss line

Reason

If the AVss pin is opened, the microcomputer may have a failure because of noise or others.

(3) Clock frequency during A-D conversion

The comparator consists of a capacity coupling, and a charge of the capacity will be lost if the clock frequency is too low. Thus, make sure the following during an A-D conversion.

- f(XIN) is 500 kHz or more in middle-/high-speed mode.
- Do not execute the **STP** instruction.
- When the A-D converter is operated at low-speed mode, $f(X_{IN})$ do not have the lower limit of frequency, because of the A-D converter has a built-in self-oscillation circuit.

2.7 Watchdog timer

This paragraph explains the registers setting method and the notes relevant to the watchdog timer.

2.7.1 Memory map

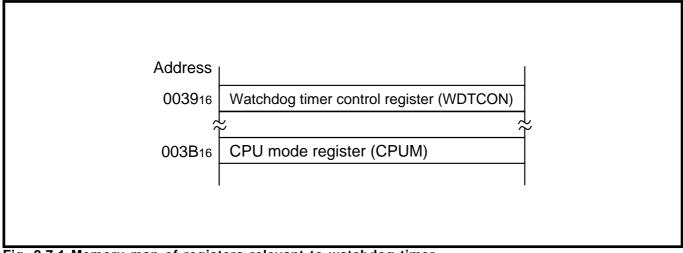
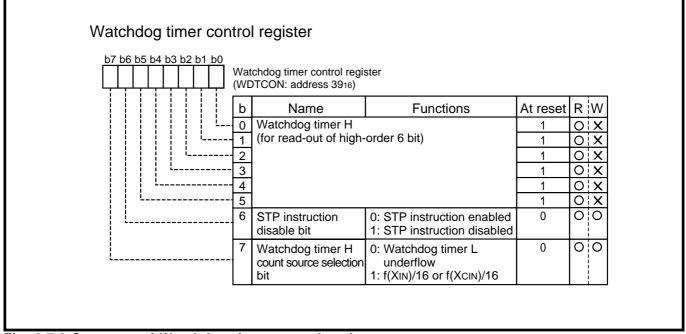


Fig. 2.7.1 Memory map of registers relevant to watchdog timer

2.7.2 Relevant registers





2.7 Watchdog timer

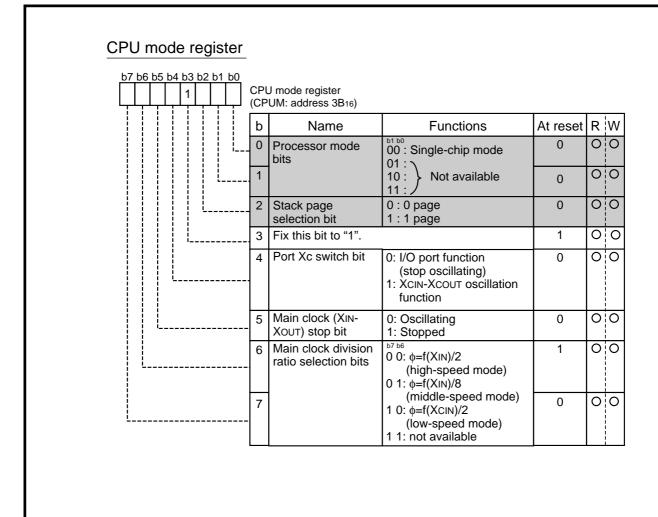


Fig. 2.7.3 Structure of CPU mode register

2.7.3 Watchdog timer application examples

(1) Detection of program runaway

Outline: If program runaway occurs, let the microcomputer reset, using the internal timer for detection of program runaway.

- **Specifications**: •An underflow of watchdog timer H is judged to be program runaway, and the microcomputer is returned to the reset status.
 - •Before the watchdog timer underflows, "0" is set into bits 6 and 7 of the watchdog timer control register at every cycle in a main routine.
 - •High-speed mode is used as a main clock division ratio.
 - •An underflow signal of the watchdog timer L is supplied as the count source of watchdog timer H.

Figure 2.7.4 shows a watchdog timer connection and division ratio setting; Figure 2.7.5 shows the relevant registers setting; Figure 2.7.6 shows the control procedure.

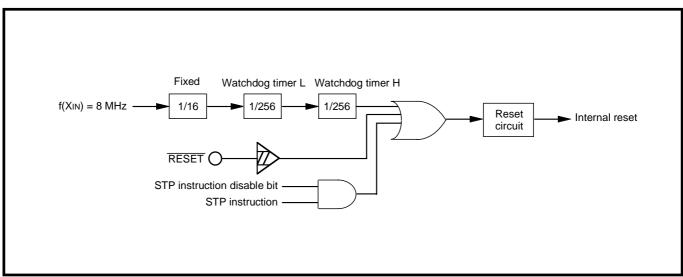


Fig. 2.7.4 Watchdog timer connection and division ratio setting

2.7 Watchdog timer

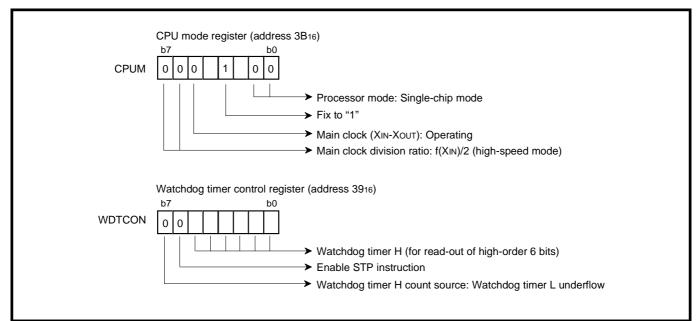


Fig. 2.7.5 Relevant registers setting

	RESET	
SEI CLT CLI	T ⊃ UM (address 3B16)	•All interrupts disabled •Processor mode: Single-chip mode •Main clock f(Xוر): Operating •High-speed mode selected as main clock division ratio •Interrupts enabled
	CON (address 3916), bit7, bit6 ← 002	 Watchdog timer L underflow selected as Watchdog timer H count source STP instruction enabled ("FF1e" is set to Watchdog timer H and Watchdog timer L, respectively.)

Fig. 2.7.6 Control procedure

2.7.4 Notes on watchdog timer

- •Make sure that the watchdog timer does not underflow while waiting Stop release, because the watchdog timer keeps counting during that term.
- •When the STP instruction disable bit has been set to "1", it is impossible to switch it to "0" by a program.

2.8 Reset

2.8 Reset

2.8.1 Connection example of reset IC

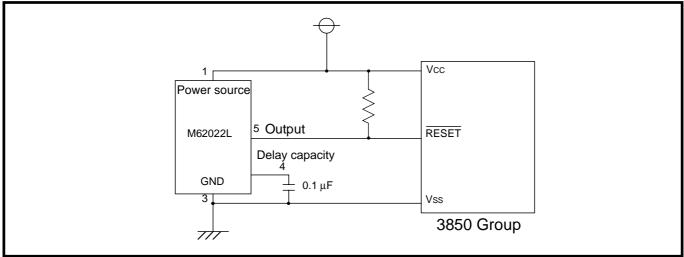


Fig. 2.8.1 Example of poweron reset circuit

Figure 2.8.2 shows the system example which switches to the RAM backup mode by detecting a drop of the system power source voltage with the INT interrupt.

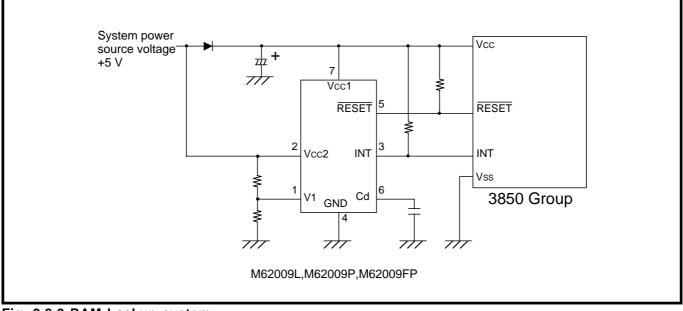


Fig. 2.8.2 RAM backup system

2.8 Reset

2.8.2 Notes on RESET pin

(1) Connecting capacitor

In case where the $\overline{\text{RESET}}$ signal rise time is long, connect a ceramic capacitor or others across the $\overline{\text{RESET}}$ pin and the Vss pin. Use a 1000 pF or more capacitor for high frequency use. When connecting the capacitor, note the following :

- Make the length of the wiring which is connected to a capacitor as short as possible.
- Be sure to verify the operation of application products on the user side.

Reason

If the several nanosecond or several ten nanosecond impulse noise enters the RESET pin, it may cause a microcomputer failure.

(2) Reset release after power on

When releasing the reset after power on, such as power-on reset, release reset after X_{IN} passes more than 20 cycles in the state where the power supply voltage is 2.7 V or more and the X_{IN} oscillation is stable.

• Reason

To release reset, the RESET pin must be held at an "L" level for 20 cycles or more of X_{IN} in the state where the power source voltage is between 2.7 V and 5.5 V, and X_{IN} oscillation is stable.

2.9 Clock generating circuit

This paragraph explains how to set the registers relevant to the clock generating circuit and describes an application example.

2.9.1 Relevant registers

b7 b6 b5 b4 b3 b2 b1 b0	CPI (CP	J mode register UM: address 3B16)			
	b	Name	Functions	At reset	RW
	0	Processor mode bits	00 : Single-chip mode	0	00
	1		10 : Not available	0	00
	2	Stack page selection bit	0 : 0 page 1 : 1 page	0	00
	3	Fix this bit to "1".	•	1	00
	4	Port Xc switch bit	0: I/O port function (stop oscillating) 1: XCIN-XCOUT oscillation function	0	00
	5	Main clock (XIN- XOUT) stop bit	0: Oscillating 1: Stopped	0	00
	6	Main clock division ratio selection bits	^{b7 b6} 0 0: φ=f(XIN)/2 (high-speed mode) 0 1: φ=f(XIN)/8	1	00
	7		(middle-speed mode) 1 0: φ=f(XcιN)/2 (low-speed mode) 1 1: not available	0	00

Fig. 2.9.1 Structure of CPU mode register

2.9 Clock generating circuit

2.9.2 Clock generating circuit application example

Status transition during power failure
 Outline: The clock counts up every second by using the timer interrupt during a power failure.

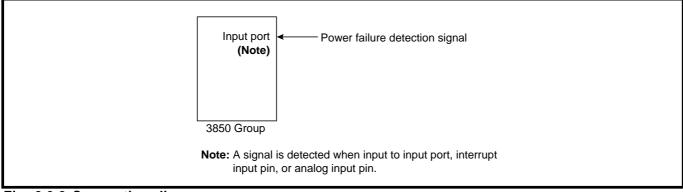


Fig. 2.9.2 Connection diagram

Specifications: •Reducing power dissipation as low as possible while maintaining clock function •Clock: f(X_{IN}) = 8 MHz, f(X_{CIN}) = 32.768 kHz •Port processing Input port: Fixed to "H" or "L" level externally. Output port: Fixed to output level that does not cause current flow to the external. (Example) Fix to "H" for an LED circuit that turns on at "L" output level. I/O port: Input port → Fixed to "H" or "L" level externally. Output port → Output of data that does not consume current VREF pin: Stop VREF current dissipation by terminating A-D conversion operation.

Figure 2.9.3 shows the status transition diagram during power failure and Figure 2.9.4 shows the setting of relevant registers.

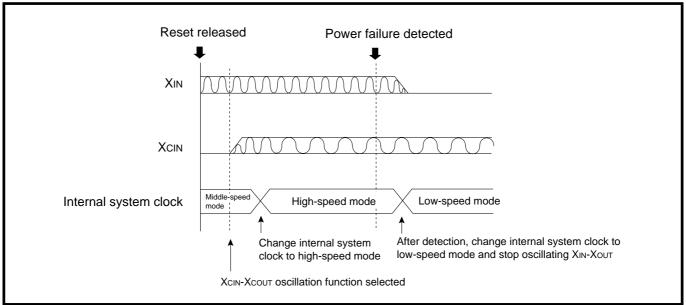
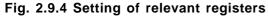


Fig. 2.9.3 Status transition diagram during power failure

2.9 Clock generating circuit

CPU mode register (address 3B16)
Main clock: High-speed mode (f(XIN)/2) (Note 1)
CPU mode register (address 3B16)
CPUM 0 0 1 1 0 0 (Note 2)
Port Xc: XciN–Xcout oscillation function
CPU mode register (address 3B16)
CPUM 10011100 Internal system clock: Low-speed mode (f(Xcin)/2)
CPU mode register (address 3B16)
Main clock f(XIN): Stopped
Notes 1: This setting is necessary only when selecting the high-speed mode. 2: When selecting the middle-speed mode, bit 6 is "1".



2.9 Clock generating circuit

Control procedure: To prepare for a power failure, set the relevant registers in the order shown below.

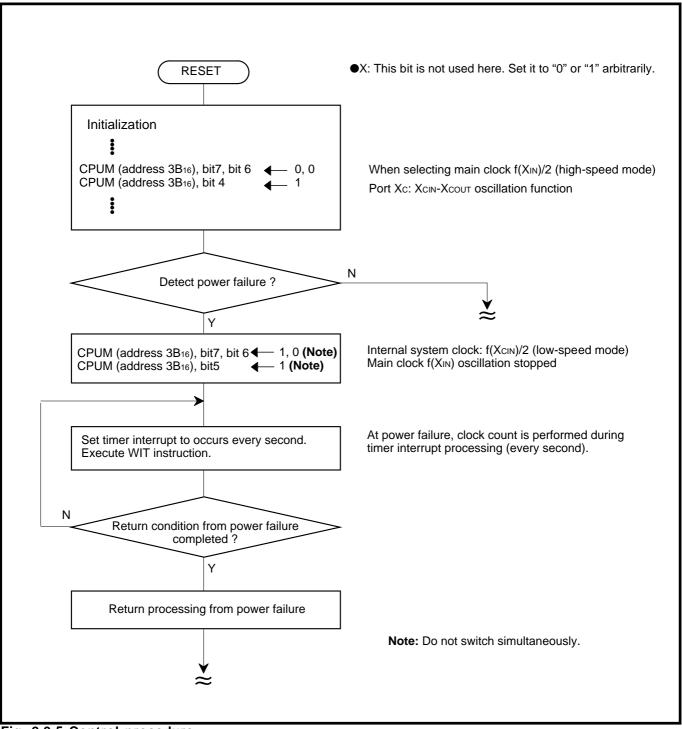


Fig. 2.9.5 Control procedure

2.10 Standby function

The 3850 group is provided with standby functions to stop the CPU by software and put the CPU into the low-power operation.

The following two types of standby functions are available.

•Stop mode using STP instruction

•Wait mode using WIT instruction

2.10.1 Relevant registers

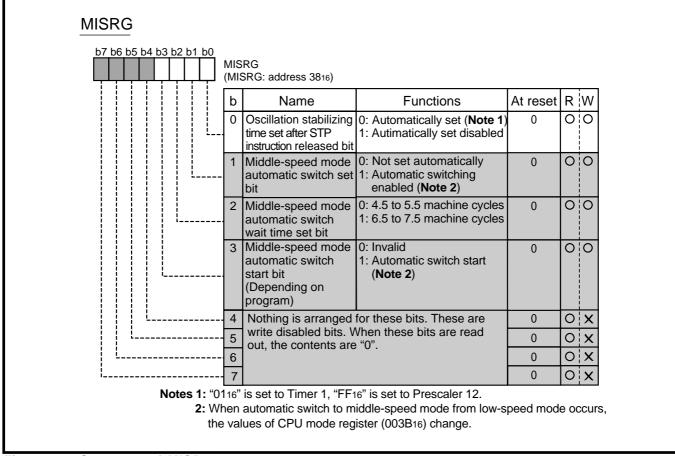


Fig. 2.10.1 Structure of MISRG

2.10 Standby function

2.10.2 Stop mode

The stop mode is set by executing the STP instruction. In the stop mode, the oscillation of both clocks $(X_{IN}-X_{OUT}, X_{CIN}-X_{COUT})$ stop and the internal clock ϕ stops at the "H" level. The CPU stops and peripheral units stop operating. As a result, power dissipation is reduced.

(1) State in stop mode

Table 2.10.1 shows the state in the stop mode.

Item	State in stop mode
Oscillation	Stopped.
CPU	Stopped.
Internal clock ϕ	Stopped at "H" level.
I/O ports P0-P4	Retains the state at the STP instruction execution.
Timer	Stopped. (Timers 1, 2, X, Y)
	However, Timers X and Y can be operated in the event counter
	mode.
PWM	Stopped.
Watchdog timer	Stopped.
Serial I/O1, Serial I/O2	Stopped.
	However, these can be operated only when an external clock
	is selected.
A-D converter	Stopped.

(2) Release of stop mode

The stop mode is released by a reset input or by the occurrence of an interrupt request. Note the differences in the restoration process according to reset input or interrupt request, as described below.

■Restoration by reset input

The stop mode is released by holding the $\overrightarrow{\text{RESET}}$ pin to the "L" input level during the stop mode. Oscillation is started when all ports are in the input state and the stop mode of the main clock (X_{IN-XouT}) is released.

Oscillation is unstable when restarted. For this reason, time for stabilizing of oscillation (oscillation stabilizing time) (Note) is required. The input of the $\overrightarrow{\text{RESET}}$ pin should be held at the "L" level until oscillation stabilizes.

When the RESET pin is held at the "L" level for 20 cycles or more of X_{IN} after the oscillation has stabilized, the microcomputer will go to the reset state. After the input level of the RESET pin is returned to "H", the reset state is released in approximately 10.5 to 18.5 cycles of the X_{IN} input. Figure 2.10.2 shows the oscillation stabilizing time at restoration by reset input.

At release of the stop mode by reset input, the internal RAM retains its contents previous to the reset. However, the previous contents of the CPU register and SFR are not retained.

For more details concerning reset, refer to "2.8 Reset".

Note: For the setting of oscillation stabilizing time, refer to MISRG (address 003816).

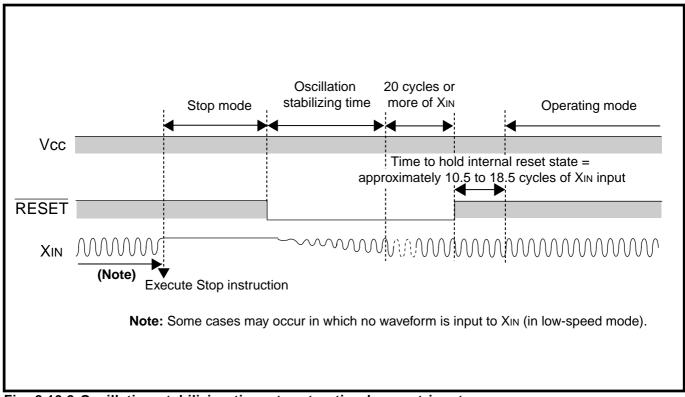


Fig. 2.10.2 Oscillation stabilizing time at restoration by reset input

2.10 Standby function

■Restoration by interrupt request

The occurrence of an interrupt request in the stop mode releases the stop mode. As a result, oscillation is resumed. The interrupts available for restoration are:

- •INT₀-INT₃
- •CNTR₀, CNTR₁
- •Serial I/O (1, 2) using an external clock
- •Timer X, Y using an external event count

However, when using any of these interrupt requests for restoration from the stop mode, <u>in order</u> to enable the selected interrupt, you must execute the STP instruction after setting the following conditions.

[Necessary register setting]

- ① Interrupt disable flag I = "0" (interrupt enabled)
- ② Timer 1 interrupt enable bit = "0" (interrupt disabled)
- ③ Interrupt request bit of interrupt source to be used for restoration = "0" (no interrupt request issued)
- ④ Interrupt enable bit of interrupt source to be used for restoration = "1" (interrupts enabled)

For more details concerning interrupts, refer to "2.2 Interrupts".

Oscillation is unstable when restarted. For this reason, time for stabilizing of oscillation (oscillation stabilizing time) is required. For restoration by an interrupt request, waiting time prior to supplying internal clock ϕ to the CPU is automatically generated^{*2} by Prescaler 12 and Timer 1^{*1}. This waiting time is reserved as the oscillation stabilizing time on the system clock side. The supply of internal clock ϕ to the CPU is started at the Timer 1 underflow.

Figure 2.10.3 shows an execution sequence example at restoration by the occurrence of an INT_0 interrupt request.

- *1: If the STP instruction is executed when the oscillation stabilizing time set after STP instruction released bit is "0", "FF₁₆" and "01₁₆" are automatically set in the Prescaler 12 counter/latch and Timer 1 counter/latch, respectively. When the oscillation stabilizing time set after STP instruction released bit is "1", nothing is automatically set to either Prescaler 12 or Timer 1. For this reason, any suitable value can be set to Prescaler 12 and Timer 1 for the oscillation stabilizing time.
- *2: Immediately after the oscillation is started, the count source is supplied to the prescaler 12 so that a count operation is started.

2.10 Standby function

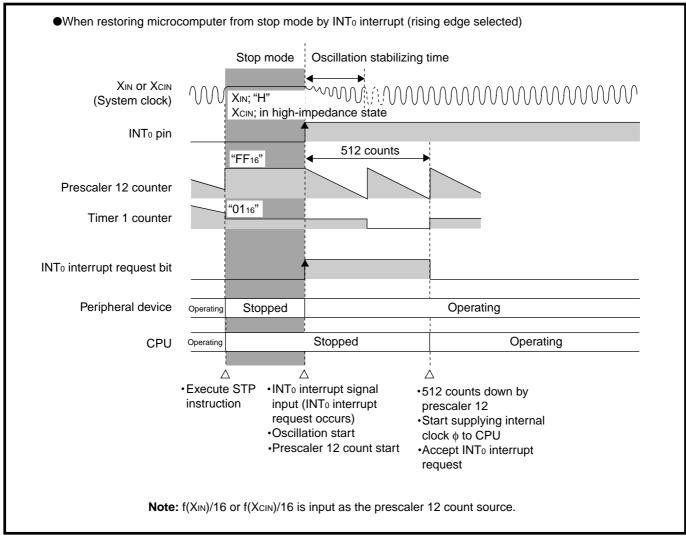


Fig. 2.10.3 Execution sequence example at restoration by occurrence of INT₀ interrupt request

(3) Notes on using stop mode

■Register setting

Since values of the prescaler 12 and Timer 1 are automatically reloaded when returning from the stop mode, set them again, respectively. (When the oscillation stabilizing time set after STP instruction released bit is "0")

■Clock restoration

After restoration from the stop mode to the normal mode by an interrupt request, the contents of the CPU mode register previous to the STP instruction execution are retained. Accordingly, if both main clock and sub clock were oscillating before execution of the STP instruction, the oscillation of both clocks is resumed at restoration.

In the above case, when the main clock side is set as a system clock, the oscillation stabilizing time for approximately 8,000 cycles of the X_{IN} input is reserved at restoration from the stop mode. At this time, note that the oscillation on the sub clock side may not be stabilized even after the lapse of the oscillation stabilizing time of the main clock side.

2.10 Standby function

2.10.3 Wait mode

The wait mode is set by execution of the WIT instruction. In the wait mode, oscillation continues, but the internal clock ϕ stops at the "H" level.

The CPU stops, but most of the peripheral units continue operating.

(1) State in wait mode

The continuation of oscillation permits clock supply to the peripheral units. Table 2.10.2 shows the state in the wait mode.

Table 2.10.2 State in wait mode

ltem	State in wait mode
Oscillation	Operating.
CPU	Stopped.
Internal clock ø	Stopped at "H" level.
I/O ports P0-P4	Retains the state at the WIT instruction execution.
Timer	Operating.
PWM	Operating.
Watchdog timer	Operating.
Serial I/O1, Serial I/O2	Operating.
A-D converter	Operating.

(2) Release of wait mode

The wait mode is released by reset input or by the occurrence of an interrupt request. Note the differences in the restoration process according to reset input or interrupt request, as described below.

In the wait mode, oscillation is continued, so an instruction can be executed immediately after the wait mode is released.

■Restoration by reset input

The wait mode is released by holding the input level of the $\overrightarrow{\text{RESET}}$ pin at "L" in the wait mode. Upon release of the wait mode, all ports are in the input state, and supply of the internal clock ϕ to the CPU is started. To reset the microcomputer, the $\overrightarrow{\text{RESET}}$ pin should be held at an "L" level for 20 cycles or more of X_{IN}. The reset state is released in approximately 10.5 cycles to 18.5 cycles of the X_{IN} input after the input of the $\overrightarrow{\text{RESET}}$ pin is returned to the "H" level.

At release of wait mode, the internal RAM retains its contents previous to the reset. However, the previous contents of the CPU register and SFR are not retained.

Figure 2.10.4 shows the reset input time.

For more details concerning reset, refer to "2.8 Reset".

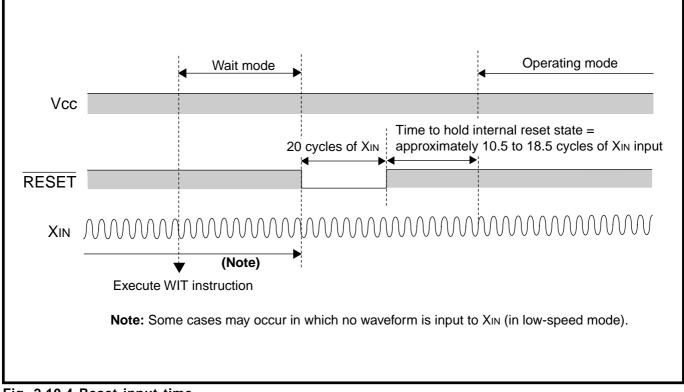


Fig. 2.10.4 Reset input time

2.10 Standby function

■Restoration by interrupt request

In the wait mode, the occurrence of an interrupt request releases the wait mode and supply of the internal clock ϕ to the CPU is started. At the same time, the interrupt request used for restoration is accepted, so the interrupt processing routine is executed.

However, when using an interrupt request for restoration from the wait mode, <u>in order to enable</u> the selected interrupt, you must execute the WIT instruction after setting the following conditions.

[Necessary register setting]

- ① Interrupt disable flag I = "0" (interrupt enabled)
- ② Interrupt request bit of interrupt source to be used for restoration = "0" (no interrupt request issued)
- ③ Interrupt enable bit of interrupt source to be used for restoration = "1" (interrupts enabled)

For more details concerning interrupts, refer to "2.2 Interrupts".

(3) Notes on wait mode

■Clock restoration

If the wait mode is released by a reset when X_{CIN} is set as the system clock and X_{IN} oscillation is stopped during execution of the WIT instruction, X_{CIN} oscillation stops, X_{IN} oscillations starts, and X_{IN} is set as the system clock.

In the above case, the $\overline{\text{RESET}}$ pin should be held at "L" until the oscillation is stabilized.

2.11 Flash memory mode

This paragraph explains the registers setting method and the notes relevant to the flash memory version.

2.11.1 Overview

The functions of the flash memory version are similar to those of the mask ROM version except that the flash memory is built-in and some of the SFR area differ from that of the mask ROM version (refer to "2.11.2 Memory map").

In the flash memory version, the built-in flash memory can be programmed or erased by using the following three modes.

- CPU rewrite mode
- Parallel I/O mode
- Standard serial I/O mode

2.11.2 Memory map

M38507F8FP/SP have 32 Kbytes of built-in flash memory. Figure 2.11.1 shows the memory map of the flash memory version.

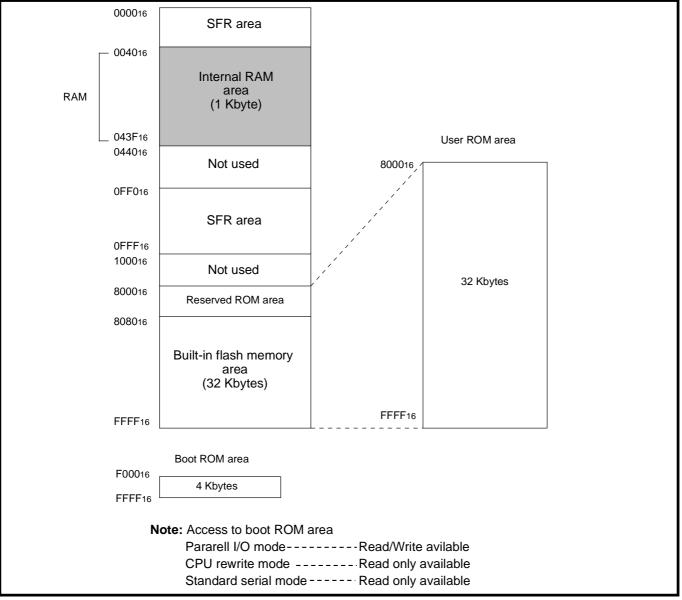
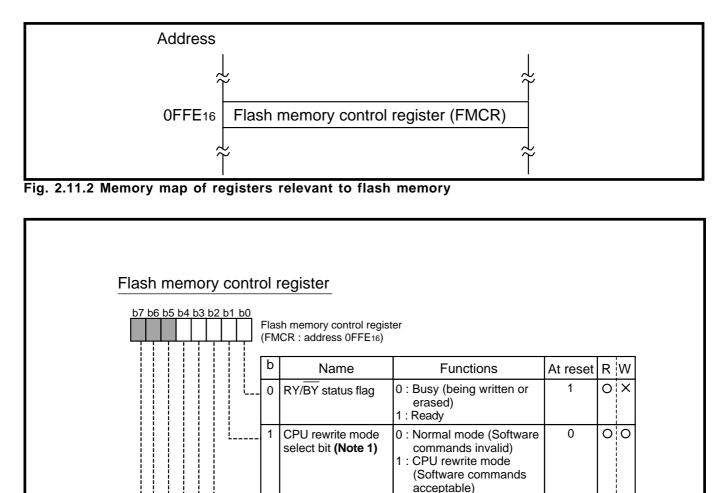
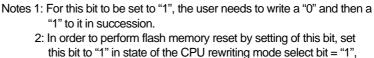


Fig. 2.11.1 Memory map of flash memory version for 3850 Group

2.11 Flash memory mode

2.11.3 Relevant registers





0: Normal mode

1: Reset

Nothing is arranged for these bits. When write,

set "0". When these bits are read out, the

1: CPU rewrite mode

0: Normal operation

0: User ROM area

1: Boot ROM area

0 X

00

00

XX

XX

XX

0

0

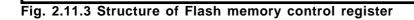
0

Undefined

Undefined

Undefined

and then set to "0" to release the reset state.



2

3

4

5

6

7

CPU rewrite mode

Flash memory reset

entry flag

bit (Note 2)

User area/Boot

area selection bit

contents are undefined.

2.11.4 Parallel I/O mode

In the parallel I/O mode, program/erase to the built-in flash memory can be performed by a EPROM programmer (EFP-I).

The memory area of program/erase is from 0F000₁₆ to 0FFFF₁₆ (boot ROM area) or from 08000₁₆ to 0FFFF₁₆ (user ROM area). Be especially careful when erasing; if the memory area is not set correctly, the products will be damaged eternally.

Table 2.11.1 shows the setting of programmers when programming in the parallel I/O mode. •EFP-I provided by Suisei Electronics System Co., Ltd. (http://www.suisei.co.jp/index_e.htm) (product available in Asia and Oceania only)

Table 2.11.1 Setting of progra	mmers when parallel programming
--------------------------------	---------------------------------

Products	Parallel unit	Boot ROM area	User ROM area	
M38507F8FP	EF3850F-42E			
M38507F8SP	EF3850F-42S	0F000 ₁₆ to 0FFFF ₁₆	0800016 to 0FFFF16	

2.11.5 Standard serial I/O mode

Table 2.11.2 shows a pin connection example (4 wires) between the programmer (EFP-I; Serial unit EF1SRP-01U is required additionally) and the microcomputer when programming in the serial I/O mode. •EFP-I provided by Suisei Electronics System Co., Ltd. (http://www.suisei.co.jp/index_e.htm) (product available in Asia and Oceania only)

Table 2.11.2 Connection example to programmer when serial programming (4 wires)

	EFP-I (E	F1SRP-01U)	3850 Group flash memory version		
Function	Signal name EF1RP-01U side		Pin name	Pin number	
		connector Line number			
Transfer clock input	T_SCLK1	9	P26/Sclk1	10	
Serial data input	T_RXD	11	P2₅/TxD	11	
Serial data output	T_TXD	10	P2₄/RxD	12	
Transmit/Receive enable output	T_BUSY	12	P27/CNTR0/SRDY1	9	
5 V input	T_VPP	3	CNVss	15	
Reset input	T_RESET	14	RESET (Note 1)	18	
Target board power source monitor input	T_VDD (Note 2)	4	Vcc (Note 2)	1	
GND	GND (Note 3)	1, 2, 15, 16	Vss, AVss (Note 3)	21, 3	

Notes 1: Since reset release after write verification is not performed, when operating MCU after writing, separate a target connection cable.

2: Supply Vcc of EFP-I side from user side so that the power supply voltage of the output buffer used by the EFP-I side becomes the same as user side power supply voltage (Vcc).

3: Four pins (No. 1, 2, 15, and 16) of the EF1SRP-01U side connector are prepared for GND signal. When connecting with a target board, although connection of only one pin does not have a problem, we recommend connecting with two or more pins.

2.11 Flash memory mode

2.11.6 CPU rewrite mode

In the CPU rewrite mode, issuing software commands through the Central Processing Unit (CPU) can rewrite the built-in flash memory. Accordingly, the contents of the built-in flash memory can be rewritten with the microcomputer itself mounted on board, without using the programmer.

Store the rewrite control program to the built-in flash memory in advance. The built-in flash memory cannot be read in the CPU rewrite mode. Accordingly, after transferring the rewrite control program to the internal RAM, execute it on the RAM.

The following commands can be used in the CPU rewrite mode: read array, read status register, clear status register, program, erase all block, and block erase. For details concerning each command, refer to "CHAPTER 1 Flash memory mode (CPU rewrite mode)".

(1) CPU rewrite mode beginning/release procedures

Operation procedure in the CPU rewrite mode for the built-in flash memory is described below. As for the control example, refer to "2.11.7 (2) Control example in the CPU rewrite mode".

[Beginning procedure]

- ① Apply 5 V±10 % to the CNVss/VPP pin (at selecting boot ROM area).
- 2 Release reset.
- ③ Set bits 6 and 7 (main clock division ratio selection bits) of the CPU mode register.
- ④ After CPU rewrite mode control program is transferred to internal RAM, jump to this control program on RAM. (The following operations are controlled by this control program).
- 5 Apply 5 V±10 % to the CNVss/VPP pin (in single-chip mode).
- 6 Set "1" to the CPU rewrite mode select bit (bit 1 of address 0FFE₁₆).
- For this bit to be set to "1", the user needs to write "0" and then "1" to it in succession.
- ⑦ Read the CPU rewrite mode entry flag (bit 2 of address 0FFE₁₆) to confirm that the CPU rewrite mode is set to "1".
- [®] Flash memory operations are executed by using software commands.

Note: The following procedures are also necessary.

- Control for data which is input from the external (serial I/O etc.) and to be programmed to the flash memory.
- Initial setting for ports, etc.
- Writing to the watchdog timer

[Release procedure]

① Execute the read command or set the flash memory reset bit (bit 3 of address 0FFE16).

② Set the CPU rewrite mode select bit (bit 0 of address 0FFE16) to "0".

2.11 Flash memory mode

Also, execute the following processing before the CPU reprogramming mode is selected so that interrupts will not occur during the CPU reprogramming mode.

• Set the interrupt disable flag (I) to "1"

When the watchdog timer has already started, write to the watchdog timer control register (address 1E₁₆) periodically during the CPU reprogramming mode in order not to generate the reset by the underflow of the watchdog timer H.

During the program or erase execution, watchdog timer is automatically cleared. Accordingly, the inernal reset by underflow does not occur.

When the interrupt request or reset occurs in the CPU reprogramming mode, the microcomputer enters the following state;

• Interrupt occurs

This may cause a program runaway because the read from the flash memory which has the interrupt vector area cannot be performed.

• Underflow of watchdog timer H, reset

This may cause a microcomputer reset; the built-in flash memory control circuit and the flash memory control register are reset. When reset state is released with CNVss = "H", CPU starts in the boot mode.

Also, when the above interrupt and reset occur during program/erase, error data may still exist after reset release because the reprogramming of the flash memory is not completed, so that reprogramming of the flash memory in the parallel I/O mode or serial I/O mode is required.

2.11 Flash memory mode

2.11.7 Flash memory mode application examples

The control pin processing example on the system board in the serial I/O mode and the control example in the CPU rewrite mode are described below.

(1) Control pin connection example on the system board in serial I/O mode

As shown in Figure 2.11.4, in the serial I/O mode, the built-in flash memory can be rewritten with the microcomputer mounted on board. Connection examples of control pins (P2₄/RxD, P2₅/TxD, P2₅/ScLk1, P27/SRDY1, P41, CNVss, and RESET pin) in the serial I/O mode are described below.

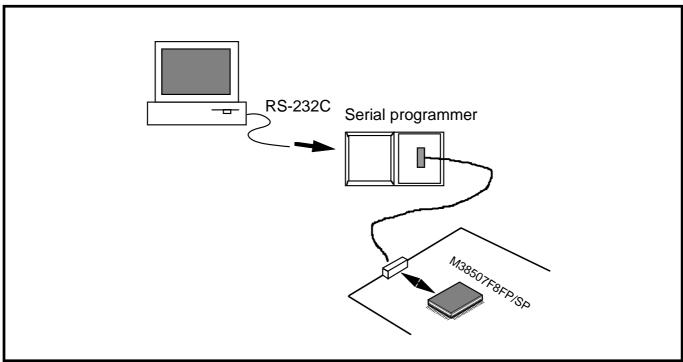
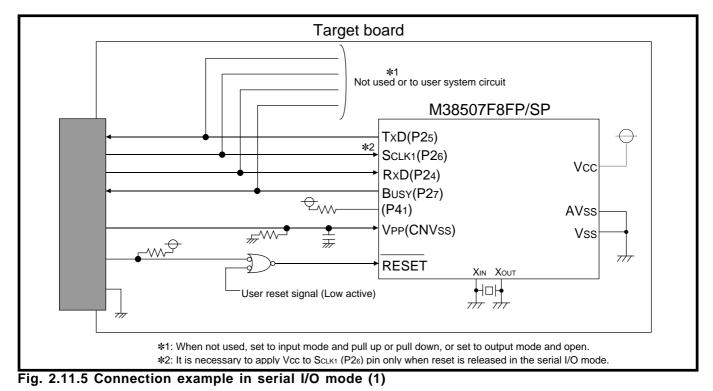


Fig. 2.11.4 Rewrite example of built-in flash memory in serial I/O mode

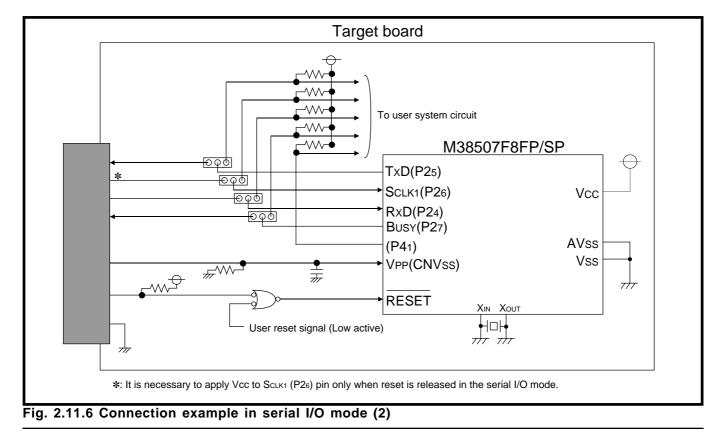
① When control signals are not affected to user system circuit

When the control signals in the serial I/O mode are not used or not affected to the user system circuit, they can be connected as shown in Figure 2.11.5.



② When control signals are affected to user system circuit-1

Figure 2.11.6 shows an example that the jumper switch cut-off the control signals not to supply to the user system circuit in the serial I/O mode.



2.11 Flash memory mode

③ When control signals are affected to user system circuit-2

Figure 2.11.7 shows an example that the analog switch (74HC4066) cut-off the control signals not to supply to the user system circuit in the serial I/O mode.

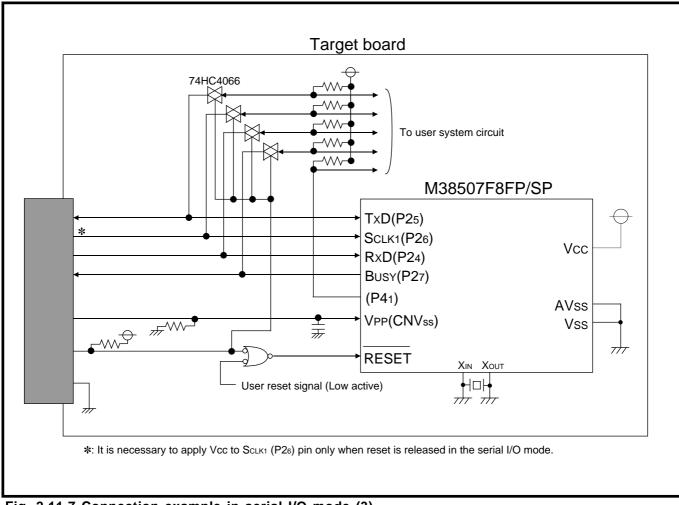
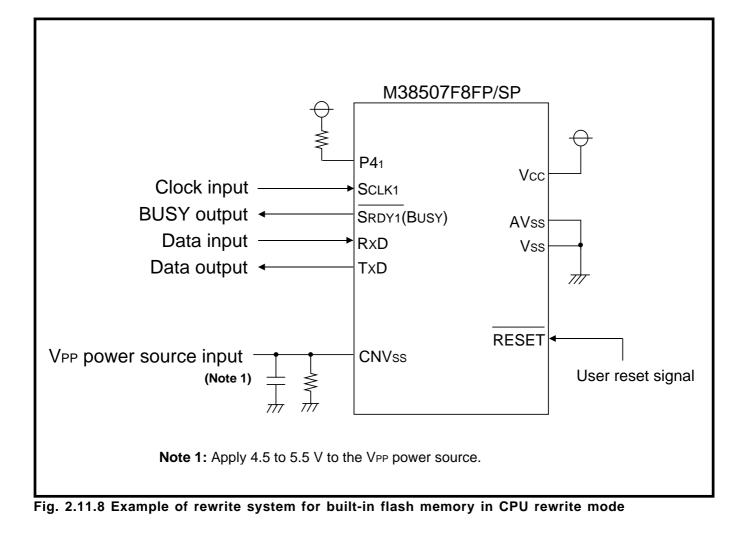


Fig. 2.11.7 Connection example in serial I/O mode (3)

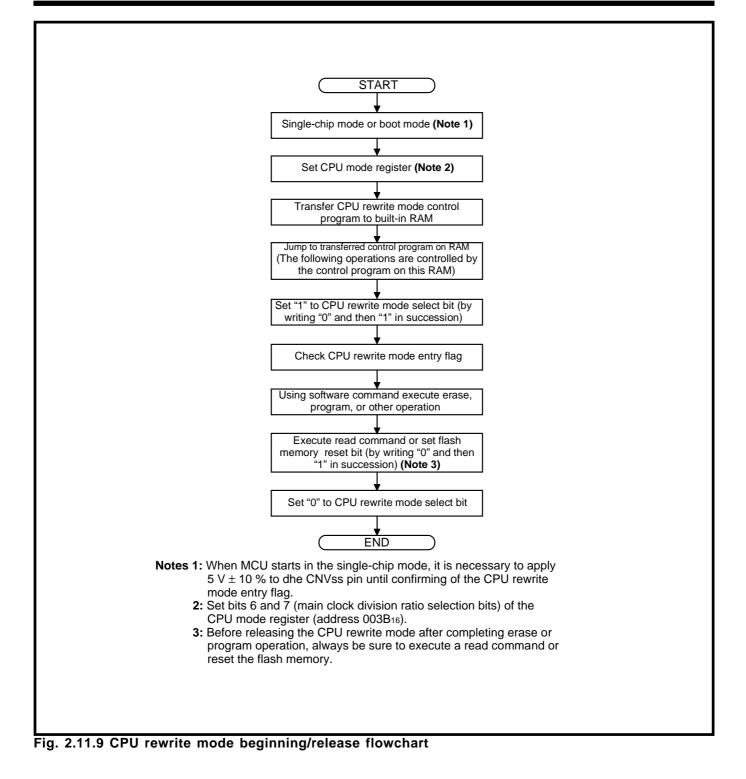
(2) Control example in CPU rewrite mode

In this example, data is received by using serial I/O, and the data is programmed to the built-in flash memory in the CPU rewrite mode.

Figure 2.11.8 shows an example of the reprogramming system for the built-in flash memory in the CPU rewrite mode. Figure 2.11.9 shows the CPU rewrite mode beginning/release flowchart.



2.11 Flash memory mode



2.11.8 Notes on CPU rewrite mode

(1) Operation speed

During CPU rewrite mode, set the internal clock ϕ 4 MHz or less using the main clock division ratio selection bits (bits 6 and 7 of address 003B₁₆).

(2) Instructions inhibited against use The instructions which refer to the internal data of the flash memory cannot be used during the CPU rewrite mode.

(3) Interrupts inhibited against use

The interrupts cannot be used during the CPU rewrite mode because they refer to the internal data of the flash memory.

(4) Watchdog timer

In case of the watchdog timer has been running already, the internal reset generated by watchdog timer underflow does not happen, because of watchdog timer is always clearing during program or erase operation.

(5) Reset

Reset is always valid. In case of $CNV_{SS} = "H"$ when reset is released, boot mode is active. So the program starts from the address contained in address FFFC₁₆ and FFFD₁₆ in boot ROM area.

2.11 Flash memory mode

MEMORANDUM



- 3.1 Electrical characteristics
- 3.2 Standard characteristics
- 3.3 Notes on use
- 3.4 Countermeasures against noise
- 3.5 List of registers
- 3.6 Package outline
- 3.7 Machine instructions
- 3.8 List of instruction code
- 3.9 SFR memory map
- 3.10 Pin configurations

3.1 Electrical characteristics

3.1.1 Absolute maximum ratings

Table 3.1.1 Absolute maximum ratings

Symbol	Parameter	Conditions	Ratings	Unit
Vcc	Power source voltage		-0.3 to 6.5	V
VI	Input voltage P00–P07, P10–P17, P20, P21, P24–P27, P30–P34, P40–P44, VREF		-0.3 to Vcc +0.3	V
VI	Input voltage P22, P23		-0.3 to 5.8	V
VI	Input voltage RESET, XIN	All voltages are based on Vss. Output transistors are cut off.	-0.3 to Vcc +0.3	V
VI	Input voltage CNVss		-0.3 to Vcc +0.3	V
Vo	Output voltage P00–P07, P10–P17, P20, P21, P24–P27, P30–P34, P40–P44, Xout		-0.3 to Vcc +0.3	V
Vo	Output voltage P22, P23		-0.3 to 5.8	V
Pd	Power dissipation	Ta = 25 °C	1000 (Note)	mW
Topr	Operating temperature		-20 to 85	°C
Tstg	Storage temperature]	-40 to 125	°C

Note : The rating becomes 300mW at the 42P2R-A/E package.

3.1.2 Recommended operating conditions

Table 3.1.2 Recommended operating conditions (1)

(Vcc = 2.7 to 5.5 V, $T_a = -20$ to 85 °C, unless otherwise noted)

Symbol		Dorom	otor		Limits		
Symbol		Parame	eter	Min.	Тур.	Max.	Unit
Vcc	Power source voltage	8 MHz (high-spee	ed mode)	4.0	5.0	5.5	V
	r offor obditor tollage	8 MHz (middle-sp	peed mode), 4 MHz (high-speed mode)	2.7	5.0	5.5	v
Vss	Power source voltage				0		V
Vref	A-D convert reference voltage		2.0		Vcc	V	
AVss	Analog power source vol	Itage			0		V
Via	Analog input voltage	AN0-AN	4	AVss		Vcc	V
Vih	"H" input voltage	P00-P07	7, P10–P17, P20–P27, P30–P34, P40–P44	0.8Vcc		Vcc	V
Viн	"H" input voltage	RESET,	XIN, CNVss	0.8Vcc		Vcc	V
VIL	"L" input voltage	P00-P07	7, P10–P17, P20–P27, P30–P34, P40–P44	0		0.2Vcc	V
VIL	"L" input voltage	RESET,	CNVss	0		0.2Vcc	V
VIL	"L" input voltage	Xin		0		0.16Vcc	V
Σ IOH(peak)	"H" total peak output cur	rent (Note) P00-	–P07, P10–P17, P30–P34			-80	mA
Σ IOH(peak)	"H" total peak output cur	rent (Note) P20,	, P21, P24–P27, P40–P44			-80	mA
Σ IOL(peak)	"L" total peak output curr	ent (Note) P00-	–P07, P30–P34			80	mA
Σ IOL(peak)	"L" total peak output curr	ent (Note) P10-	–P17			120	mA
Σ IOL(peak)	"L" total peak output curr	rent (Note) P20-	-P27,P40-P44			80	mA
Σ IOH(avg)	"H" total average output	current (Note) P00-	–P07, P10–P17, P30–P34			-40	mA
Σ IOH(avg)	"H" total average output	current (Note) P20,	, P21, P24–P27, P40–P44			-40	mA
Σ IOL(avg)	"L" total average output cu	rrent (Note) P00-	–P07, P30–P34			40	mA
Σ IOL(avg)	"L" total average output cu	rrent (Note) P10-	–P17			60	mA
Σ IOL(avg)	"L" total average output	current (Note) P20-	-P27,P40-P44			40	mA

Note : The total output current is the sum of all the currents flowing through all the applicable ports. The total average current is an average value measured over 100 ms. The total peak current is the peak value of all the currents.

Table 3.1.3 Recommended operating conditions (2)

(Vcc = 2.7 to 5.5 V, Ta = -20 to 85 °C, unless otherwise noted)

Symbol		Parameter		Limits	Limits	
Symbol				Тур.	Max.	Unit
IOH(peak)	"H" peak output current (Note 1)	P00-P07, P10-P17, P20, P21, P24-P27, P30-P34, P40-P44			-10	mA
IOL(peak)	"L" peak output current (Note 1)	P00-P07, P20-P27, P30-P34, P40-P44			10	mA
IOL(peak)	"L" peak output current (Note 1)	P10-P17			20	mA
IOH(avg)	"H" average output current (Note 2)	P00-P07, P10-P17, P20, P21, P24-P27, P30-P34, P40-P44			-5	mA
IOL(avg)	"L" average output current (Note 2)	P00–P07, P20–P27, P30–P34, P40–P44			5	mA
IOL(avg)	"L" average output current (Note 2)	P10–P17			15	mA
f(XIN)	Internal clock oscillation frequer	ncy (Vcc = 4.0 to 5.5V) (Note 3)			8	MHz
f(XIN)	Internal clock oscillation frequer	ncy (Vcc = 2.7 to 5.5V) (Note 3)			4	MHz

Notes 1: The peak output current is the peak current flowing in each port.

2: The average output current IOL(avg), IOH(avg) are average value measured over 100 ms.

3: When the oscillation frequency has a duty cycle of 50%.

3.1.3 Electrical characteristics

Table 3.1.4 Electrical characteristics (1)

(Vcc = 2.7 to 5.5 V, Vss = 0 V, Ta = -20 to 85 °C, unless otherwise noted)

Symbol	Parameter	Test conditions				
			Min.	Тур.	Max.	Unit
Voн	"H" output voltage	Iон = -10 mA	Vcc-2.0			V
	P00–P07, P10–P17, P20, P21,	Vcc = 4.0–5.5 V				
	P24–P27, P30–P34, P40–P44	IOH = -1.0 mA	Vcc-1.0			V
	(Note)	Vcc = 2.7–5.5 V				
Vol	"L" output voltage	IOL = 10 mA			2.0	V
	P00–P07, P20–P27, P30–P34,	VCC = 4.0–5.5 V				
	P40-P44	IOL = 1.0 mA			1.0	V
		Vcc = 2.7–5.5 V				
Vol	"L" output voltage	IOL = 20 mA			2.0	V
	P10–P17	Vcc = 4.0-5.5 V				
		IOL = 10 mA			1.0	V
		Vcc = 2.7–5.5 V				
VT+-VT-	Hysteresis			0.4		V
	CNTR0, CNTR1, INT0-INT3					
VT+-VT-	Hysteresis			0.5		V
	RxD, SCLK1, SCLK2, SIN2					
VT+-VT-	Hysteresis RESET			0.5		V
Ін	"H" input current	VI = VCC			5.0	μΑ
	P00–P07, P10–P17, P20, P21,					
	P24–P27, P30–P34, P40–P44					
Іін	"H" input current RESET, CNVss	VI = VCC			5.0	μΑ
Іін	"H" input current XIN	VI = VCC		4		μA
lı∟	"L" input current	VI = VSS			-5.0	μΑ
	P00–P07, P10–P17, P20–P27					
	P30–P34, P40–P44					
lı∟	"L" input current RESET, CNVss	VI = VSS			-5.0	μΑ
lı∟	"L" input current XIN	VI = VSS		-4		μΑ
VRAM	RAM hold voltage	When clock stopped	2.0		5.5	V

Note: P25 is measured when the P25/TxD P-channel output disable bit of the UART control register (bit 4 of address 001B16) is "0".

Table 3.1.5 Electrical characteristics (2)

(Vcc = 2.7 to 5.5 V, Vss = 0 V, Ta = -20 to 85 °C, unless otherwise noted)

Symbol Parameter	Deremeter	Test conditions					
	Parameter			Min.	Тур.	Max.	Unit
		High-speed mode f(XIN) = 8 MHz f(XCIN) = 32.768 kHz Output transistors "off"			6.8	13	mA
		High-speed mode f(XIN) = 8 MHz (in WIT state) f(XCIN) = 32.768 kHz Output transistors "off"			1.6		mA
		Low-speed mode f(XIN) = stopped	Except M38507F8FP/SP		60	200	μA
		f(Xcin) = 32.768 kHz Output transistors "off"	M38507F8FP/SP		250		μA
		Low-speed mode f(XIN) = stopped f(XCIN) = 32.768 kHz (in WIT state) Output transistors "off"	Except M38507F8FP/SP		20	40	μA
Icc			M38507F8FP/SP		70		μA
	Power source current		Except M38507F8FP/SP		20	55	μA
			M38507F8FP/SP		150		μA
		f(XIN) = stopped	Except M38507F8FP/SP		5.0	10.0	μA
		f(XCIN) = 32.768 kHz (in WIT state Output transistors "off"	M38507F8FP/SP		20		μA
		Middle-speed mode f(XIN) = 8 MHz f(XCIN) = stopped Output transistors "off"			4.0	7.0	mA μA μΑ μΑ μΑ μΑ μΑ μΑ μΑ μΑ μΑ μΑ
		Middle-speed mode f(XIN) = 8 MHz (in WIT state) f(XCIN) = stopped Output transistors "off"		1.5		mA	
		Increment when A-D conversion executed f(XIN) = 8 MHz	is		800		μA
		All oscillation stopped (in STP state)	ā = 25 °C		0.1	1.0	μA
		Output transistors "off"	ā = 85 °C			10	μA

3.1.4 A-D converter characteristics

Table 3.1.6 A-D converter characteristics

<u>(Vcc = 2.7 to 5.5 V, Vss = AVss = 0 V, Ta = -20 to 85 °C, f(XIN) = 8 MHz, unless otherwise noted</u>)

Currente e l	Parameter		Test senditions		Linit		
Symbol			Test conditions	Min.	Тур.	Max.	Unit
_	Resolution					10	bit
-	Absolute accuracy (excluding quantization error)					±4	LSB
tCONV	Conversion time		High-speed mode, Middle-speed mode			61	2tc(XIN)
			Low-speed mode		40		μs
RLADDER	Ladder resistor				35		kΩ
IVREF	Reference power source input current	VREF "on"	VREF = 5.0 V	50	150	200	μA
		VREF "off"				5.0	
li(AD)	A-D port input current				0.5	5.0	μA

3.1.5 Timing requirements and switching characteristics

Table 3.1.7 Timing requirements (1)

(Vcc = 4.0 to 5.5 V, Vss = 0 V, Ta = –20 to 85 °C, unless otherwise noted)

Sumbol	Desemator		Limits		Unit
Symbol	Parameter	Min.	Тур.	Max.	
tw(RESET)	Reset input "L" pulse width	20			XIN cycle
tC(XIN)	External clock input cycle time	125			ns
twh(Xin)	External clock input "H" pulse width	50			ns
twl(Xin)	External clock input "L" pulse width	50			ns
tc(CNTR)	CNTR0, CNTR1 input cycle time	200			ns
twн(CNTR)	CNTR0, CNTR1 input "H" pulse width	80			ns
twL(CNTR)	CNTR0, CNTR1 input "L" pulse width	80			ns
twн(INT)	INTo to INT3 input "H" pulse width	80			ns
tw∟(INT)	INTo to INT3 input "L" pulse width	80			ns
tC(SCLK1)	Serial I/O1 clock input cycle time (Note)	800			ns
tWH(SCLK1)	Serial I/O1 clock input "H" pulse width (Note)	370			ns
tWL(SCLK1)	Serial I/O1 clock input "L" pulse width (Note)	370			ns
tsu(RxD-SCLK1)	Serial I/O1 input setup time	220			ns
th(SCLK1-RxD)	Serial I/O1 input hold time	100			ns
tC(SCLK2)	Serial I/O2 clock input cycle time	1000			ns
tWH(SCLK2)	Serial I/O2 clock input "H" pulse width	400			ns
tWL(SCLK2)	Serial I/O2 clock input "L" pulse width	400			ns
tsu(SIN2-SCLK2)	Serial I/O2 clock input setup time	200			ns
th(SCLK2-SIN2)	Serial I/O2 clock input hold time	200			ns

Note : When f(XIN) = 8 MHz and bit 6 of address 001A16 is "1" (clock synchronous).

Divide this value by four when f(XIN) = 8 MHz and bit 6 of address 001A16 is "0" (UART).

Table 3.1.8 Timing requirements (2)

(Vcc = 2.7 to 5.5 V, Vss = 0 V, Ta = -20 to 85 °C, unless otherwise noted)

Sumbol	Demension		Limits		Unit
Symbol	Parameter	Min.	Тур.	Max.	
tw(RESET)	Reset input "L" pulse width	20			XIN cycle
tC(XIN)	External clock input cycle time	250			ns
twh(Xin)	External clock input "H" pulse width	100			ns
twl(Xin)	External clock input "L" pulse width	100			ns
tc(CNTR)	CNTR0, CNTR1 input cycle time	500			ns
twH(CNTR)	CNTR0, CNTR1 input "H" pulse width	230			ns
twL(CNTR)	CNTR0, CNTR1 input "L" pulse width	230			ns
twн(INT)	INTo to INT3 input "H" pulse width	230			ns
tw∟(INT)	INTo to INT3 input "L" pulse width	230			ns
tC(SCLK1)	Serial I/O1 clock input cycle time (Note)	2000			ns
tWH(SCLK1)	Serial I/O1 clock input "H" pulse width (Note)	950			ns
tWL(SCLK1)	Serial I/O1 clock input "L" pulse width (Note)	950			ns
tsu(RxD-SCLK1)	Serial I/O1 input setup time	400			ns
th(SCLK1-RxD)	Serial I/O1 input hold time	200			ns
tC(SCLK2)	Serial I/O2 clock input cycle time	2000			ns
tWH(SCLK2)	Serial I/O2 clock input "H" pulse width	950			ns
twL(SCLK2)	Serial I/O2 clock input "L" pulse width	950			ns
tsu(SIN2-SCLK2)	Serial I/O2 clock input setup time	400			ns
th(SCLK2-SIN2)	Serial I/O2 clock input hold time	300			ns

Note : When f(XIN) = 4 MHz and bit 6 of address 001A16 is "1" (clock synchronous).

Divide this value by four when f(XIN) = 4 MHz and bit 6 of address 001A16 is "0" (UART).

Table 3.1.9 Switching characteristics (1)

(Vcc = 4.0 to 5.5 V, Vss = 0 V, Ta = -20 to 85 °C, unless otherwise noted)

Symbol	Parameter	Test conditions	Limits	Linit		
	Parameter	Test conditions	Min.	Тур.	Max.	Unit
tWH (SCLK1)	Serial I/O1 clock output "H" pulse width		tC(SCLK1)/2-30			ns
tWL (SCLK1)	Serial I/O1 clock output "L" pulse width	Fig.3.1.1	tC(SCLK1)/2-30			ns
td (SCLK1-TXD)	Serial I/O1 output delay time (Note 1)				140	ns
tv (SCLK1-TXD)	Serial I/O1 output valid time (Note 1)		-30			ns
tr (SCLK1)	Serial I/O1 clock output rising time				30	ns
tf (SCLK1)	Serial I/O1 clock output falling time				30	ns
tWH (SCLK2)	Serial I/O2 clock output "H" pulse width	-	tC(SCLK2)/2-160			ns
tWL (SCLK2)	Serial I/O2 clock output "L" pulse width	1	tC(SCLK2)/2-160			ns
td (SCLK2-SOUT2)	Serial I/O2 output delay time (Note 2)]			200	ns
tv (SCLK2-SOUT2)	Serial I/O2 output valid time (Note 2)]	0			ns
tf (SCLK2)	Serial I/O2 clock output falling time				30	ns
tr (CMOS)	CMOS output rising time (Note 3)	1		10	30	ns
tf (CMOS)	CMOS output falling time (Note 3)]		10	30	ns

Notes 1: When the P25/TxD P-channel output disable bit of the UART control register (bit 4 of address 001B16) is "0".

2: When the P01/SOUT2 and P02/SCLK2 P-channel output disable bit of the Serial I/O2 control register 1 (bit 7 of address 001516) is "0". 3: The XOUT pin is excluded.

Table 3.1.10 Switching characteristics (2)

(Vcc = 2.7 to 5.5 V, Vss = 0 V, Ta = –20 to 85 °C, unless otherwise noted)

Symbol	Denemeter	Test senditions	Limits		Unit	
	Parameter	Test conditions	Min.	Тур.	Max.	Unit
tWH (SCLK1)	Serial I/O1 clock output "H" pulse width		tC(SCLK1)/2-50			ns
tWL (SCLK1)	Serial I/O1 clock output "L" pulse width	Fig.3.1.1	tC(SCLK1)/2-50			ns
td (SCLK1-TXD)	Serial I/O1 output delay time (Note 1)				350	ns
tv (SCLK1-TXD)	Serial I/O1 output valid time (Note 1)		-30			ns
tr (SCLK1)	Serial I/O1 clock output rising time				50	ns
tf (SCLK1)	Serial I/O1 clock output falling time				50	ns
tWH (SCLK2)	Serial I/O2 clock output "H" pulse width		tC(SCLK2)/2-240			ns
tWL (SCLK2)	Serial I/O2 clock output "L" pulse width		tC(SCLK2)/2-240			ns
td (SCLK2-SOUT2)	Serial I/O2 output delay time (Note 2)				400	ns
tv (SCLK2-SOUT2)	Serial I/O2 output valid time (Note 2)	7	0			ns
tf (SCLK2)	Serial I/O2 clock output falling time				50	ns
tr (CMOS)	CMOS output rising time (Note 3)			20	50	ns
tf (CMOS)	CMOS output falling time (Note 3)			20	50	ns

Notes 1: When the P25/TxD P-channel output disable bit of the UART control register (bit 4 of address 001B16) is "0".

2: When the P01/SOUT2 and P02/SCLK2 P-channel output disable bit of the Serial I/O2 control register 1 (bit 7 of address 001516) is "0". 3: The XOUT pin is excluded.

3.1 Electrical characteristics

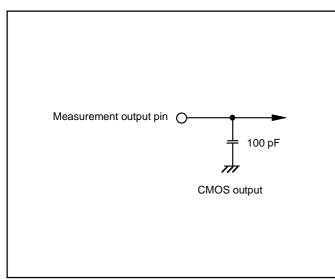


Fig. 3.1.1 Circuit for measuring output switching characteristics

3.1 Electrical characteristics

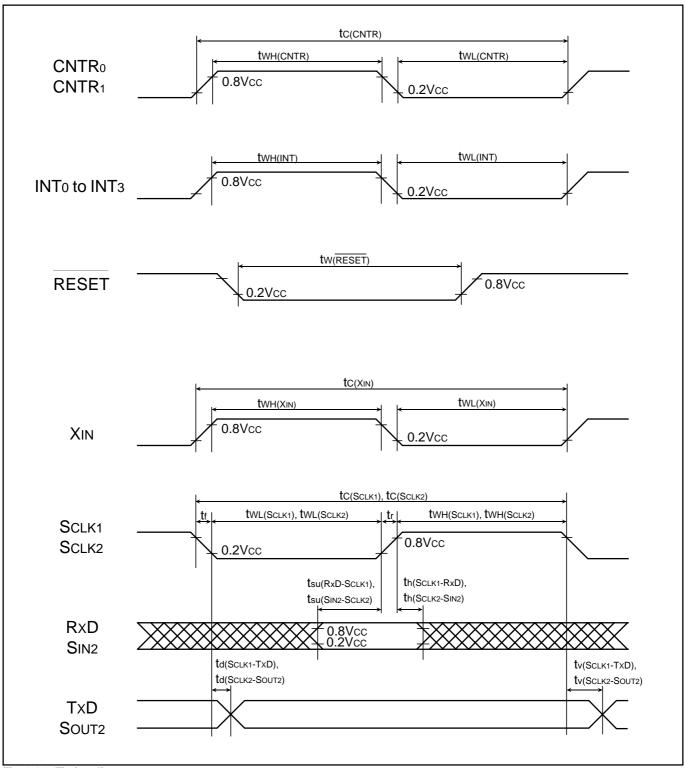


Fig. 3.1.2 Timing diagram

3.2 Standard characteristics

Standard characteristics described below are just examples of the 3850 Group (spec. H)'s characteristics and are not guaranteed. For rated values, refer to **"3.1 Electrical characteristics"**.

3.2.1 Flash memory version power source current standard characteristics

Figure 3.2.1, Figure 3.2.2, Figure 3.2.3, Figure 3.2.4, and Figure 3.2.5 show flash memory version (M38507F8) power source current standard characteristics.

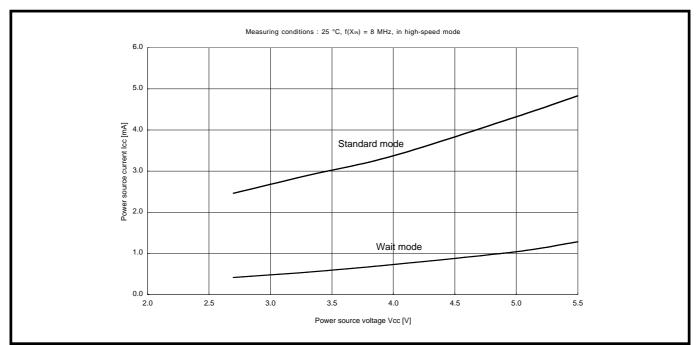


Fig. 3.2.1 Flash memory version power source current standard characteristics (in high-speed mode, $f(X_{IN}) = 8 \text{ MHz}$)

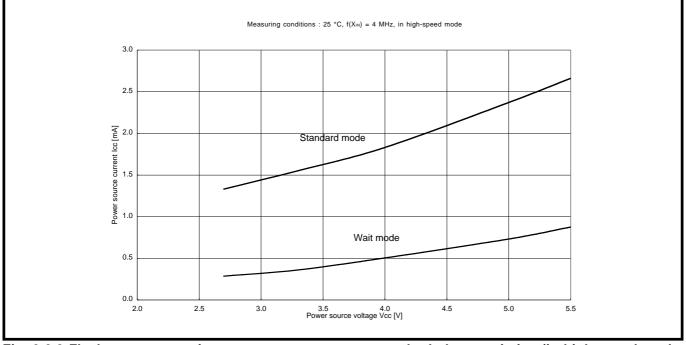


Fig. 3.2.2 Flash memory version power source current standard characteristics (in high-speed mode, $f(X_{IN}) = 4 \text{ MHz}$)

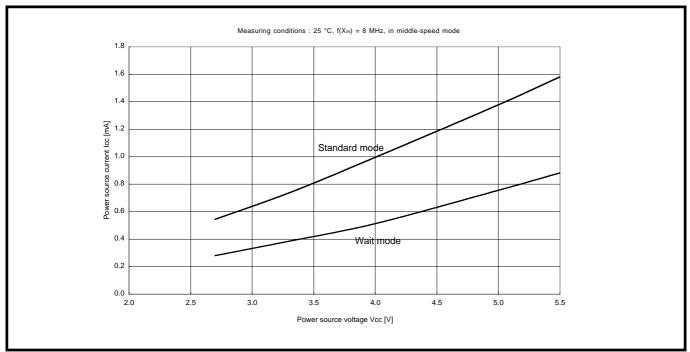


Fig. 3.2.3 Flash memory version power source current standard characteristics (in middle-speed mode, f(X_{IN}) = 8 MHz)

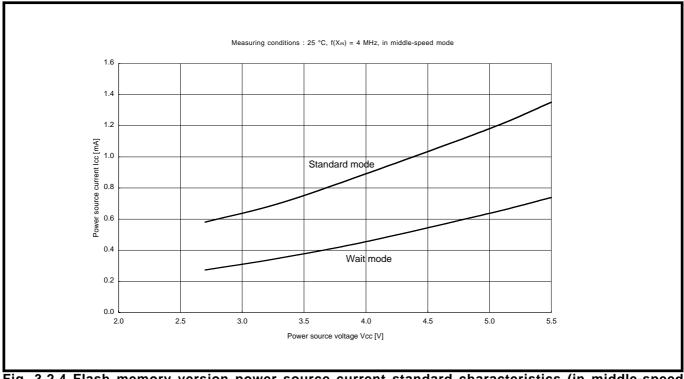


Fig. 3.2.4 Flash memory version power source current standard characteristics (in middle-speed mode, f(X_{IN}) = 4 MHz)

3.2 Standard characteristics

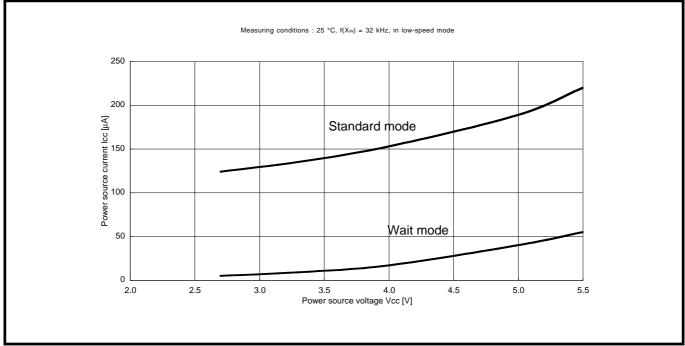


Fig. 3.2.5 Flash memory version power source current standard characteristics (in low-speed mode)

3.2 Standard characteristics

3.2.2 Mask ROM version power source current standard characteristics

Figure 3.2.6, Figure 3.2.7, Figure 3.2.8, Figure 3.2.9 and Figure 3.2.10 show mask ROM version (M38503M2H, M38503M4H, M38504M6, M38507M8) power source current standard characteristics.

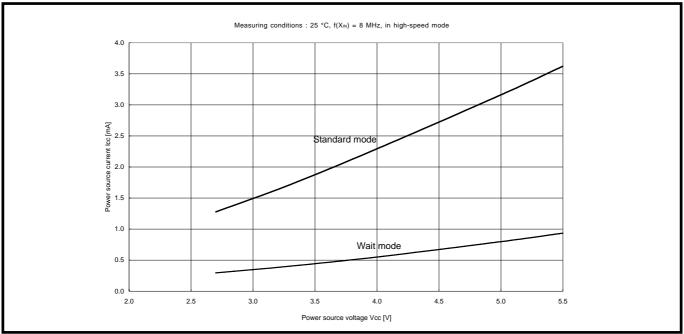


Fig. 3.2.6 Mask ROM version power source current standard characteristics (in high-speed mode, $f(X_{IN}) = 8 \text{ MHz}$)

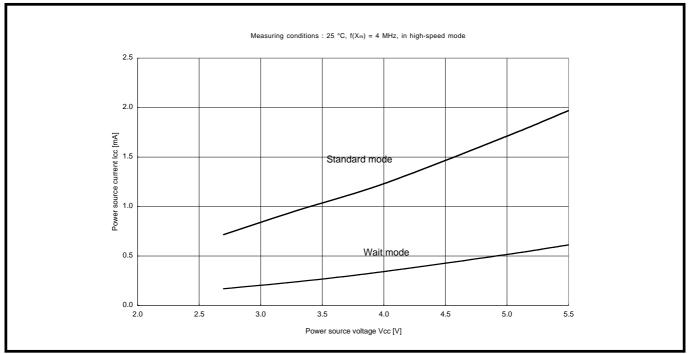


Fig. 3.2.7 Mask ROM version power source current standard characteristics (in high-speed mode, $f(X_{IN}) = 4 \text{ MHz}$)

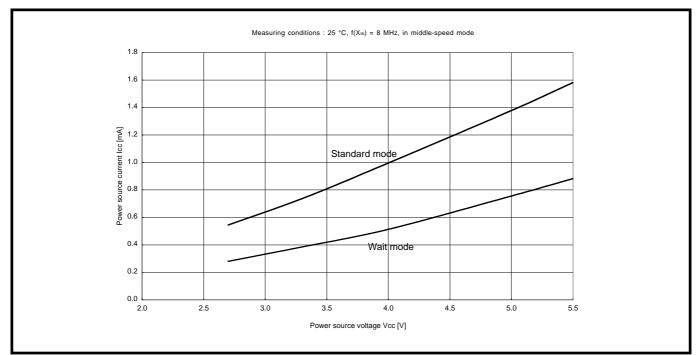
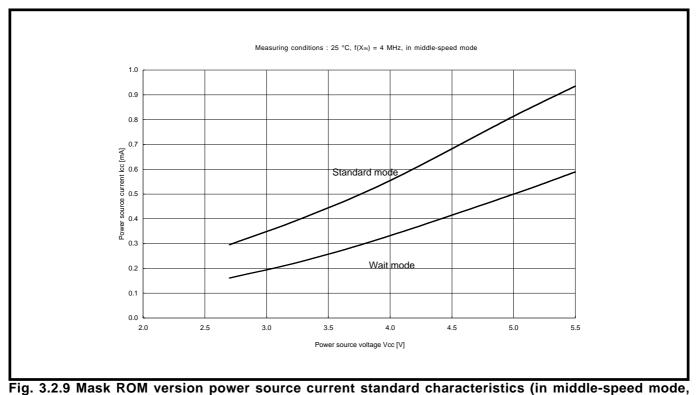
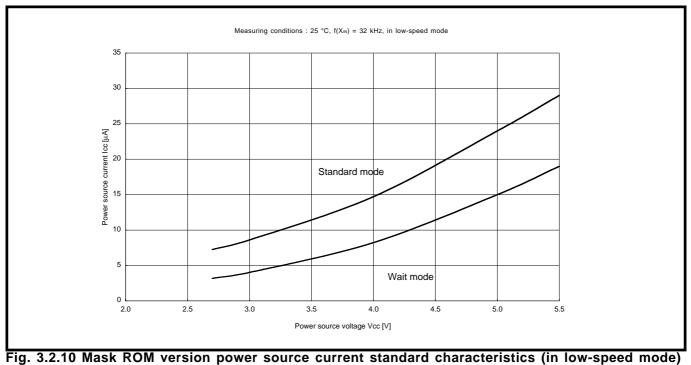


Fig. 3.2.8 Mask ROM version power source current standard characteristics (in middle-speed mode, $f(X_{IN}) = 8 \text{ MHz}$)



 $f(X_{IN}) = 4 MHz$



ig. 5.2.10 mask from version power source current standard characteristics (in low-speed mou

3.2 Standard characteristics

3.2.3 PROM version power source current standard characteristics

Figure 3.2.11, Figure 3.2.12, Figure 3.2.13, Figure 3.2.14, and Figure 3.2.15 show flash memory version (M38504E6) power source current standard characteristics.

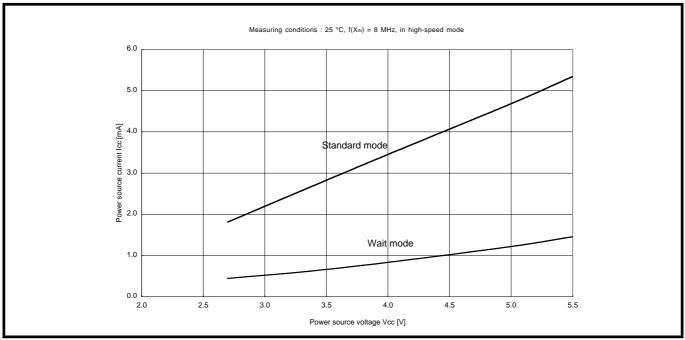


Fig. 3.2.11 PROM version power source current standard characteristics (in high-speed mode, f(X_{IN}) = 8 MHz)

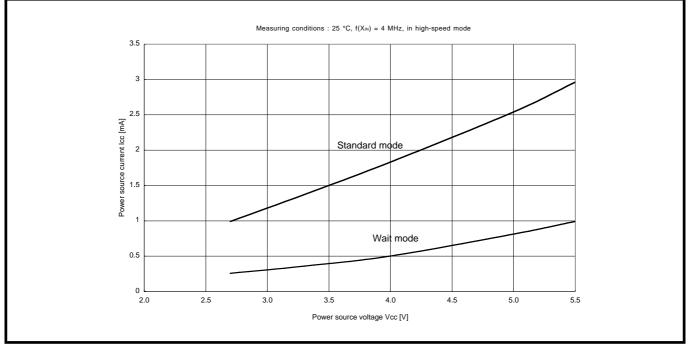


Fig. 3.2.12 PROM version power source current standard characteristics (in high-speed mode, f(X_{IN}) = 4 MHz)

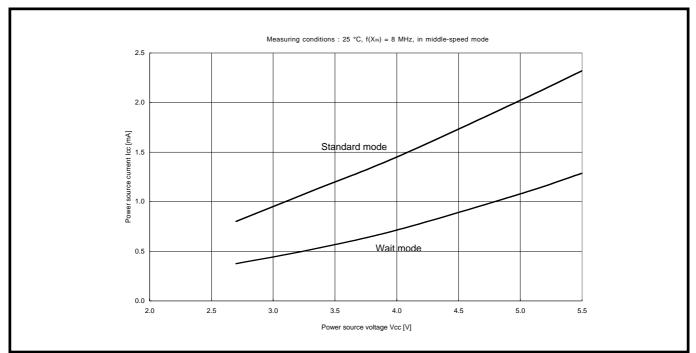
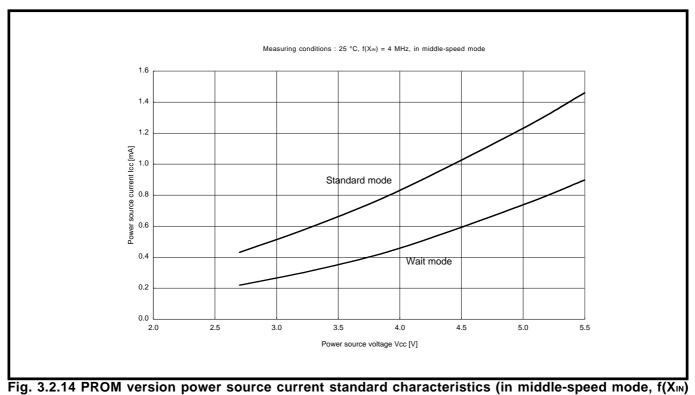


Fig. 3.2.13 PROM version power source current standard characteristics (in middle-speed mode, f(X_{IN}) = 8 MHz)



= 4 MHz)

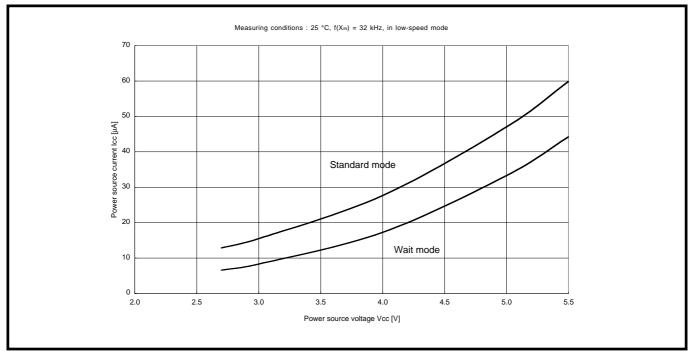


Fig. 3.2.15 PROM version power source current standard characteristics (in low-speed mode)

3.2 Standard characteristics

3.2.4 Flash memory version port standard characteristics

Figure 3.2.16, Figure 3.2.17, Figure 3.2.18 and Figure 3.2.19 show flash memory version (M38507F8) port standard characteristics.

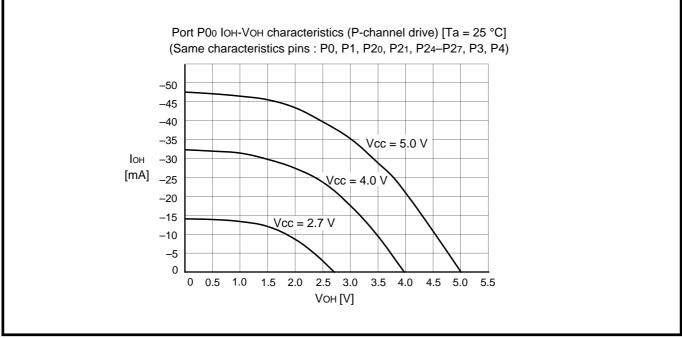


Fig. 3.2.16 CMOS output port P-channel side characteristics (Ta = 25 °C)

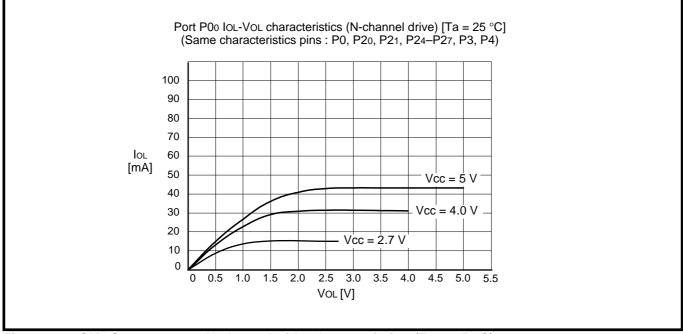


Fig. 3.2.17 CMOS output port N-channel side characteristics (Ta = 25 °C)

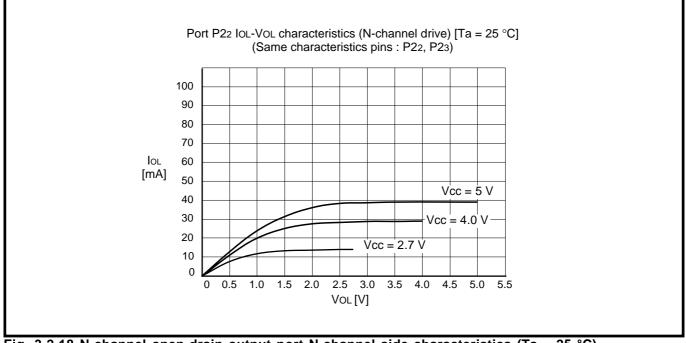


Fig. 3.2.18 N-channel open-drain output port N-channel side characteristics (Ta = 25 °C)

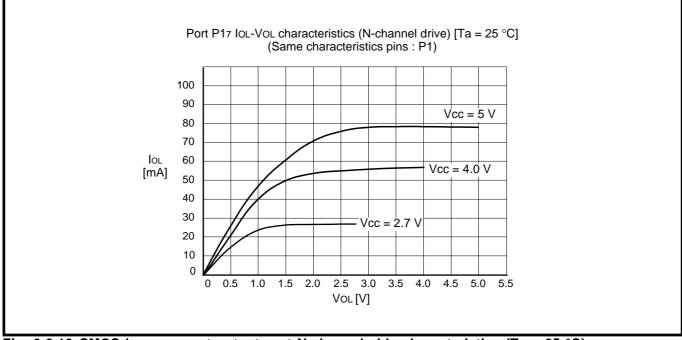


Fig. 3.2.19 CMOS large current output port N-channel side characteristics (Ta = 25 °C)

3.2 Standard characteristics

3.2.5 Mask ROM version port standard characteristics

Figure 3.2.20, Figure 3.2.21, Figure 3.2.22 and Figure 3.2.23 show mask ROM version (M38503M2H, M38503M4H, M38504M6, M38507M8) port standard characteristics.

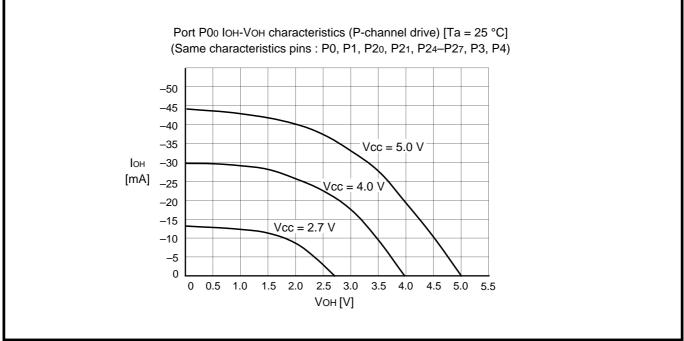


Fig. 3.2.20 CMOS output port P-channel side characteristics (Ta = 25 °C)

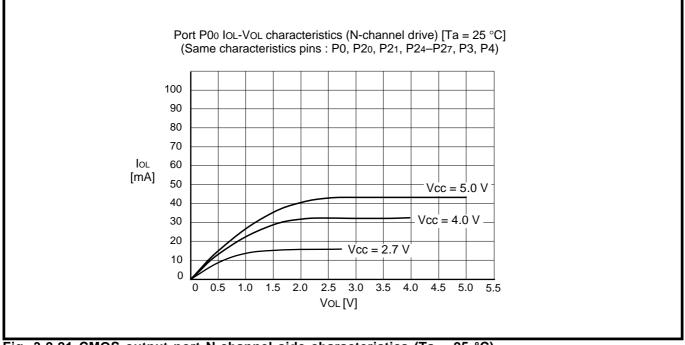
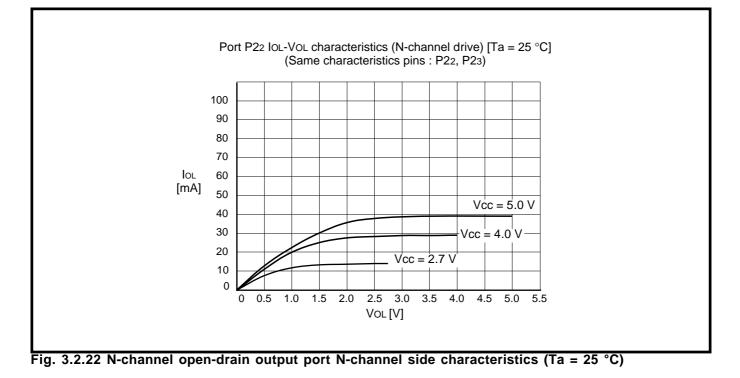


Fig. 3.2.21 CMOS output port N-channel side characteristics (Ta = 25 °C)



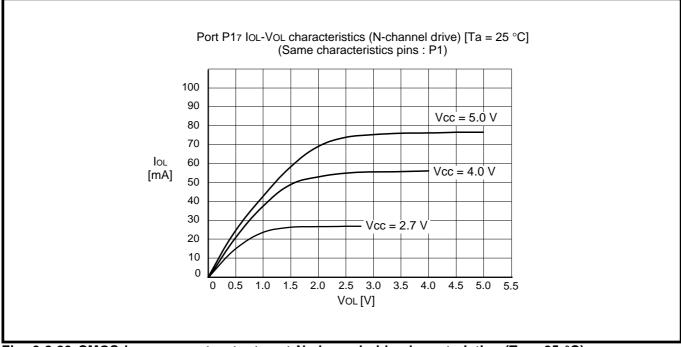


Fig. 3.2.23 CMOS large current output port N-channel side characteristics (Ta = 25 °C)

3.2 Standard characteristics

3.2.6 PROM version port standard characteristics

Figure 3.2.24, Figure 3.2.25, Figure 3.2.26 and Figure 3.2.27 show PROM version (M38504E6) port standard characteristics.

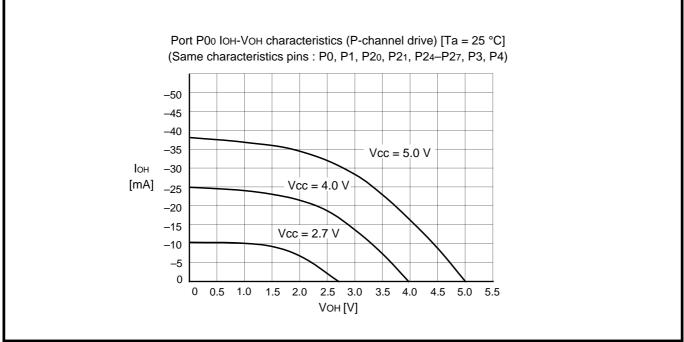


Fig. 3.2.24 CMOS output port P-channel side characteristics (Ta = 25 °C)

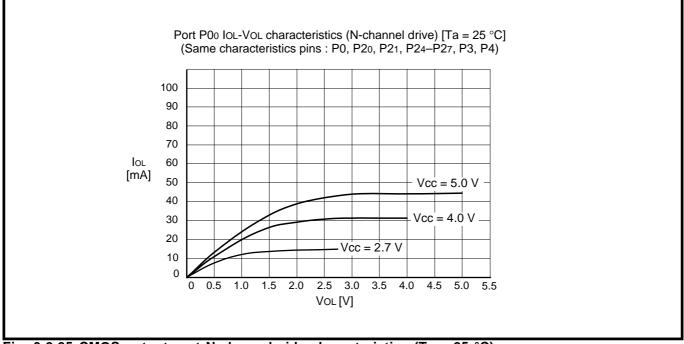


Fig. 3.2.25 CMOS output port N-channel side characteristics (Ta = 25 °C)

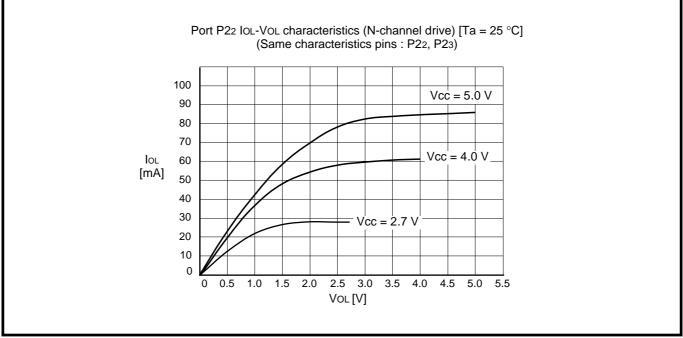


Fig. 3.2.26 N-channel open-drain output port N-channel side characteristics (Ta = 25 °C)

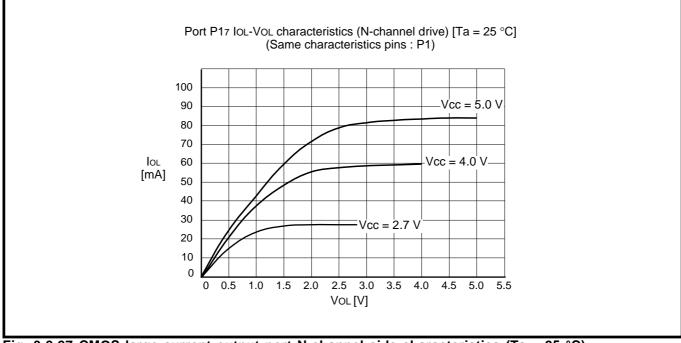


Fig. 3.2.27 CMOS large current output port N-channel side characteristics (Ta = 25 °C)

3.2 Standard characteristics

3.2.7 A-D conversion standard characteristics

(1) Definition of A-D conversion accuracy

The A-D conversion accuracy is defined below.

●Relative accuracy

① Zero transition voltage (VoT)

This means an analog input voltage when the actual A-D conversion output data changes from "0" to "1".

2 Full-scale transition voltage (VFST)

This means an analog input voltage when the actual A-D conversion output data changes from "1023" to "1022".

③ Linearity error

This means a deviation from the lone between V_{0T} and V_{FST} of a converted value between V_{0T} and V_{FST} .

④ Differential non-linearity error

This means a deviation from the input potential difference required to change a converted value between V_{0T} and V_{FST} by 1 LSB of the 1 LSB at the relative accuracy.

Absolute accuracy

This means a deviation from the ideal characteristics between 0 to V_{REF} of actual A-D conversion characteristics.

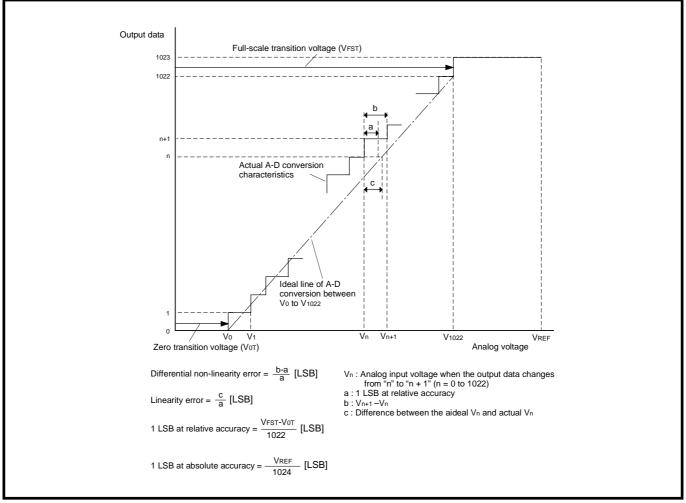


Fig. 3.2.28 Definition of A-D conversion accuracy

(2) A-D conversion standard characteristics

Figure 3.2.29, Figure 3.2.30, and Figure 3.2.31 show the A-D conversion standard characteristics of flash memory version, mask ROM version, and PROM version, respectively.

The thick lines of the graph indicate the absolute precision errors, These are expressed as the deviation from the ideal value when the output code changes. For example, the change in output code from 256 to 257 should occur at 1280 mV, but the measured value is 2.5 mV. Accordingly, the measured point of change is 1280 + 2.5 = 1282.5 mV.

The thin lines of the graph indicate the input voltage width for which the output code is constant. For example, the measured input voltage width for which the output code is 256 is 5.0 mV, so that the differential non-linear error is 5.0 - 5.0 = 0 mV (0 LSB).

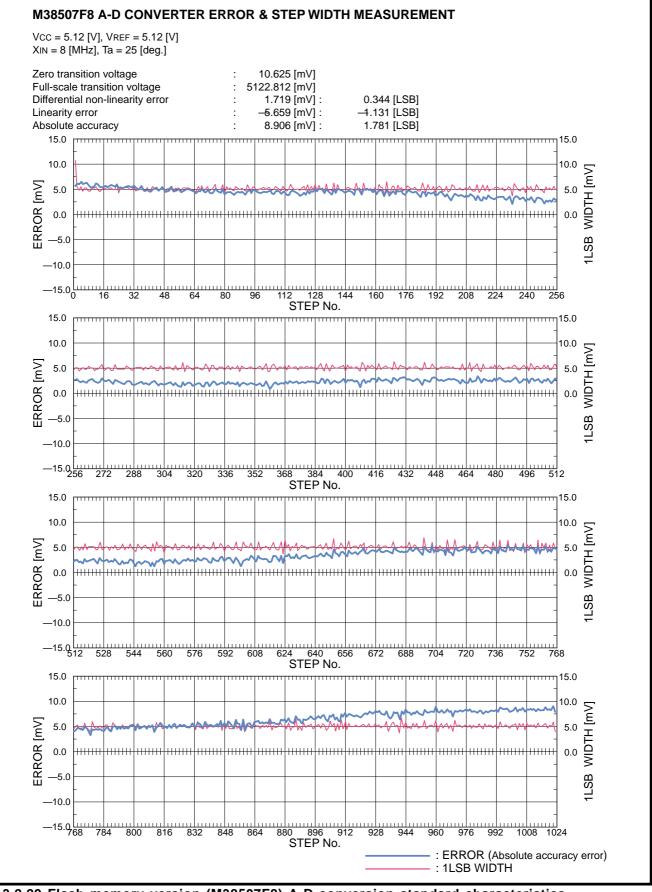


Fig. 3.2.29 Flash memory version (M38507F8) A-D conversion standard characteristics

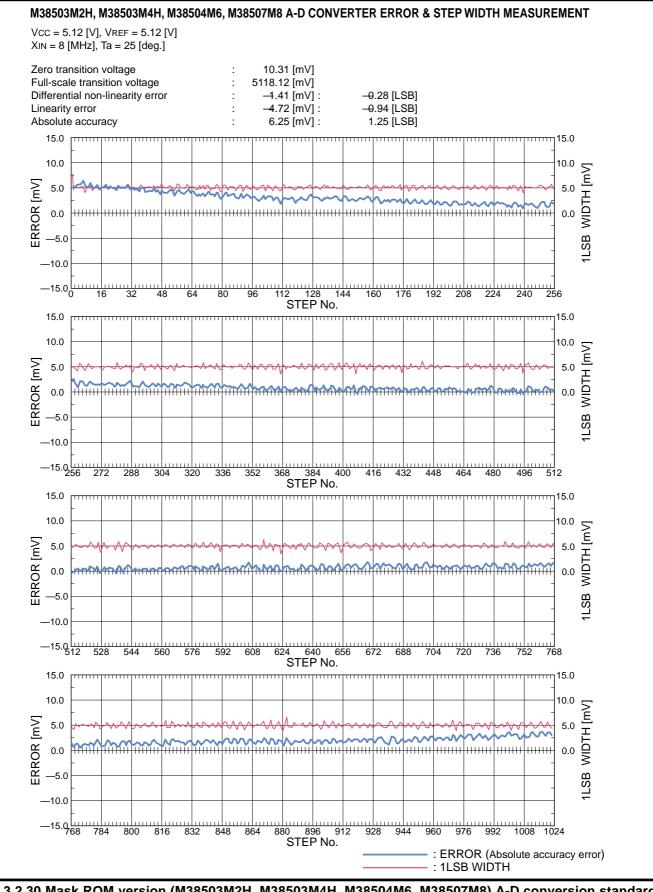


Fig. 3.2.30 Mask ROM version (M38503M2H, M38503M4H, M38504M6, M38507M8) A-D conversion standard characteristics

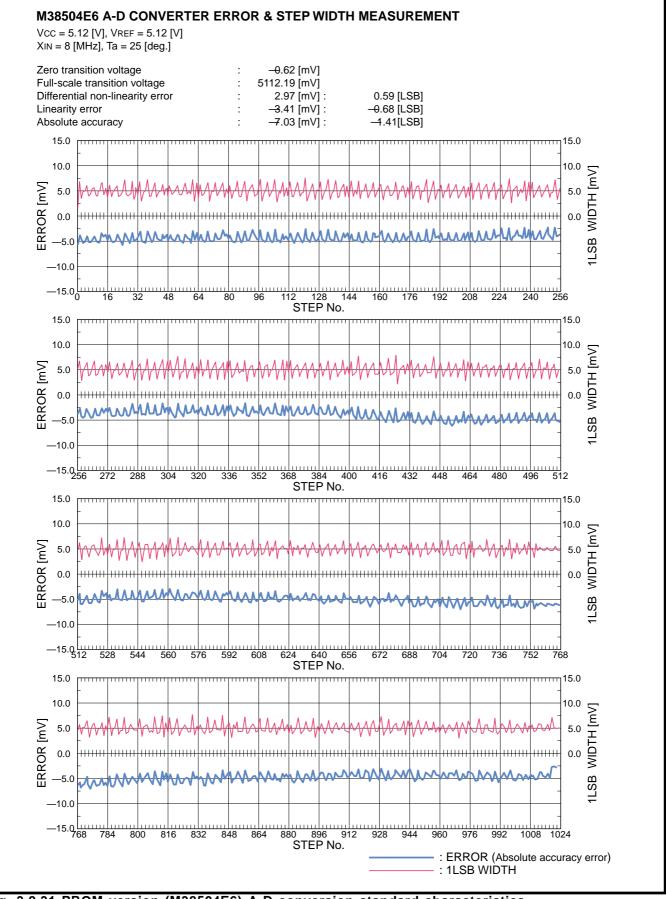


Fig. 3.2.31 PROM version (M38504E6) A-D conversion standard characteristics

3.3 Notes on use

3.3.1 Notes on input and output ports

(1) Notes in standby state

In standby state^{*1}, do not make input levels of an I/O port "undefined", especially for I/O ports of the N-channel open-drain. When setting the N-channel open-drain port as an output, do not make input levels of an I/O port "undefined", too.

Pull-up (connect the port to Vcc) or pull-down (connect the port to Vss) these ports through a resistor.

When determining a resistance value, note the following points:

- External circuit
- Variation of output levels during the ordinary operation

Reason

When setting as an input port with its direction register, the transistor becomes the OFF state, which causes the ports to be the high-impedance state.

Accordingly, the potential which is input to the input buffer in a microcomputer is unstable in the state that input levels of an I/O port are "undefined". This may cause power source current.

In I/O ports of N-channel open-drain, when the contents of the port latch are "1", even if it is set as an output port with its direction register, it becomes the same phenomenon as the case of an input port.

*1 standby state: stop mode by executing **STP** instruction wait mode by executing **WIT** instruction

(2) Modifying output data with bit managing instruction

When the port latch of an I/O port is modified with the bit managing instruction^{*2}, the value of the unspecified bit may be changed.

Reason

The bit managing instructions are read-modify-write form instructions for reading and writing data by a byte unit. Accordingly, when these instructions are executed on a bit of the port latch of an I/O port, the following is executed to all bits of the port latch.

•As for bit which is set for input port:

The pin state is read in the CPU, and is written to this bit after bit managing.

•As for bit which is set for output port:

The bit value is read in the CPU, and is written to this bit after bit managing.

Note the following:

- •Even when a port which is set as an output port is changed for an input port, its port latch holds the output data.
- •As for a bit of which is set for an input port, its value may be changed even when not specified with a bit managing instruction in case where the pin state differs from its port latch contents.

*2 Bit managing instructions: **SEB** and **CLB** instructions

3.3 Notes on use

3.3.2 Termination of unused pins

(1) Terminate unused pins

- ① Output ports : Open
- 2 Input ports :

Connect each pin to V_{cc} or V_{ss} through each resistor of 1 k Ω to 10 k Ω . As for pins whose potential affects to operation modes such as pins CNV_{ss}, INT or others, select the V_{cc} pin or the V_{ss} pin according to their operation mode.

③ I/O ports :

- Set the I/O ports for the input mode and connect them to Vcc or Vss through each resistor of 1 k Ω to 10 k $\Omega.$
 - Set the I/O ports for the output mode and open them at "L" or "H".
- When opening them in the output mode, the input mode of the initial status remains until the mode of the ports is switched over to the output mode by the program after reset. Thus, the potential at these pins is undefined and the power source current may increase in the input mode. With regard to an effects on the system, thoroughly perform system evaluation on the user side.
- Since the direction register setup may be changed because of a program runaway or noise, set direction registers by program periodically to increase the reliability of program.

④ The AVss pin when not using the A-D converter :

• When not using the A-D converter, handle a power source pin for the A-D converter, AVss pin as follows:

AVss: Connect to the Vss pin.

(2) Termination remarks

① Input ports and I/O ports :

Do not open in the input mode.

Reason

- The power source current may increase depending on the first-stage circuit.
- An effect due to noise may be easily produced as compared with proper termination 2 and 3 shown on the above.

2 I/O ports :

When setting for the input mode, do not connect to VCC or VSS directly.

Reason

If the direction register setup changes for the output mode because of a program runaway or noise, a short circuit may occur between a port and Vcc (or Vss).

③ I/O ports :

When setting for the input mode, do not connect multiple ports in a lump to Vcc or Vss through a resistor.

Reason

If the direction register setup changes for the output mode because of a program runaway or noise, a short circuit may occur between ports.

• At the termination of unused pins, perform wiring at the shortest possible distance (20 mm or less) from microcomputer pins.

3.3.3 Notes on interrupts

(1) Change of relevant register settings

When the setting of the following registers or bits is changed, the interrupt request bit may be set to "1". When not requiring the interrupt occurrence synchronized with these setting, take the following sequence.

•Interrupt edge selection register (address 3A₁₆)

•Timer XY mode register (address 2316)

Set the above listed registers or bits as the following sequence.

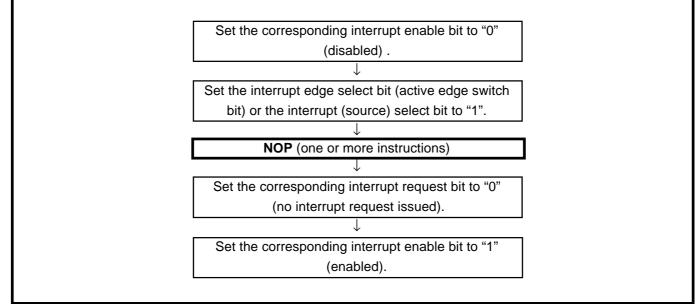


Fig. 3.3.1 Sequence of changing relevant register

Reason

When setting the followings, the interrupt request bit may be set to "1".

•When setting external interrupt active edge

Concerned register: Interrupt edge selection register (address 3A₁₆)

Timer XY mode register (address 23₁₆)

•When switching interrupt sources of an interrupt vector address where two or more interrupt sources are allocated.

Concerned register: Interrupt edge selection register (address 3A₁₆)

3.3 Notes on use

(2) Check of interrupt request bit

• When executing the **BBC** or **BBS** instruction to an interrupt request bit of an interrupt request register immediately after this bit is set to "0" by using a data transfer instruction, execute one or more instructions before executing the **BBC** or **BBS** instruction.

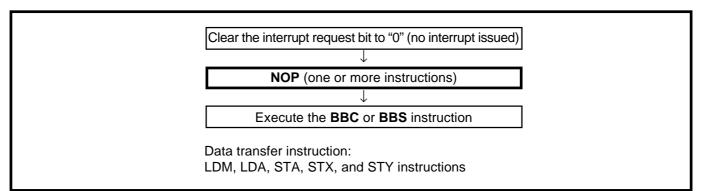


Fig. 3.3.2 Sequence of check of interrupt request bit

Reason

If the BBC or BBS instruction is executed immediately after an interrupt request bit of an interrupt request register is cleared to "0", the value of the interrupt request bit before being cleared to "0" is read.

3.3.4 Notes on timer

- If a value n (between 0 and 255) is written to a timer latch, the frequency division ratio is 1/(n+1).
- When switching the count source by the timer 12, X and Y count source selection bits, the value of timer count is altered in unconsiderable amount owing to generating of thin pulses in the count input signals.

Therefore, select the timer count source before set the value to the prescaler and the timer.

3.3.5 Notes on serial I/O

(1) Notes when selecting clock synchronous serial I/O (Serial I/O1)

① Stop of transmission operation

Clear the serial I/O1 enable bit and the transmit enable bit to "0" (Serial I/O1 and transmit disabled).

Reason

Since transmission is not stopped and the transmission circuit is not initialized even if only the serial I/O1 enable bit is cleared to "0" (Serial I/O1 disabled), the internal transmission is running (in this case, since pins TxD, RxD, SCLK1, and SRDY1 function as I/O ports, the transmission data is not output). When data is written to the transmit buffer register in this state, data starts to be shifted to the transmit shift register. When the serial I/O1 enable bit is set to "1" at this time, the data during internally shifting is output to the TxD pin and an operation failure occurs.

② Stop of receive operation

Clear the receive enable bit to "0" (receive disabled), or clear the serial I/O1 enable bit to "0" (Serial I/O1 disabled).

③ Stop of transmit/receive operation

Clear the transmit enable bit and receive enable bit to "0" simultaneously (transmit and receive disabled).

(when data is transmitted and received in the clock synchronous serial I/O mode, any one of data transmission and reception cannot be stopped.)

Reason

In the clock synchronous serial I/O mode, the same clock is used for transmission and reception. If any one of transmission and reception is disabled, a bit error occurs because transmission and reception cannot be synchronized.

In this mode, the clock circuit of the transmission circuit also operates for data reception. Accordingly, the transmission circuit does not stop by clearing only the transmit enable bit to "0" (transmit disabled). Also, the transmission circuit is not initialized by clearing the serial I/O1 enable bit to "0" (Serial I/O1 disabled) (refer to (1) ①).

④ S_{RDY1} output of reception side (Serial I/O1)

When signals are output from the S_{RDY1} pin on the reception side by using an external clock in the clock synchronous serial I/O mode, set all of the receive enable bit, the $\overline{S_{RDY1}}$ output enable bit, and the transmit enable bit to "1" (transmit enabled).

(2) Notes when selecting clock asynchronous serial I/O (Serial I/O1)

① Stop of transmission operation

Clear the transmit enable bit to "0" (transmit disabled).

Reason

Since transmission is not stopped and the transmission circuit is not initialized even if only the serial I/O1 enable bit is cleared to "0" (Serial I/O1 disabled), the internal transmission is running (in this case, since pins TxD, RxD, SCLK1, and SRDY1 function as I/O ports, the transmission data is not output). When data is written to the transmit buffer register in this state, data starts to be shifted to the transmit shift register. When the serial I/O1 enable bit is set to "1" at this time, the data during internally shifting is output to the TxD pin and an operation failure occurs.

② Stop of receive operation

Clear the receive enable bit to "0" (receive disabled).

③ Stop of transmit/receive operation

Only transmission operation is stopped.

Clear the transmit enable bit to "0" (transmit disabled).

Reason

Since transmission is not stopped and the transmission circuit is not initialized even if only the serial I/O1 enable bit is cleared to "0" (Serial I/O1 disabled), the internal transmission is running (in this case, since pins TxD, RxD, SCLK1, and SRDY1 function as I/O ports, the transmission data is not output). When data is written to the transmit buffer register in this state, data starts to be shifted to the transmit shift register. When the serial I/O1 enable bit is set to "1" at this time, the data during internally shifting is output to the TxD pin and an operation failure occurs.

Only receive operation is stopped.

Clear the receive enable bit to "0" (receive disabled).

3.3 Notes on use

(3) Setting serial I/O1 control register again (Serial I/O1)

Set the serial I/O1 control register again after the transmission and the reception circuits are reset by clearing both the transmit enable bit and the receive enable bit to "0".

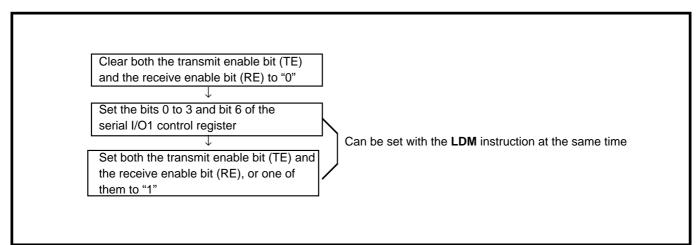


Fig. 3.3.3 Sequence of setting serial I/O1 control register again

(4) Data transmission control with referring to transmit shift register completion flag (Serial I/O1) The transmit shift register completion flag changes from "1" to "0" with a delay of 0.5 to 1.5 shift clocks. When data transmission is controlled with referring to the flag after writing the data to the transmit buffer register, note the delay.

(5) Transmit interrupt request when transmit enable bit is set (Serial I/O1)

When the transmit interrupt is used, set the transmit interrupt enable bit to transmit enabled as shown in the following sequence.

- Set the interrupt enable bit to "0" (disabled) with CLB instruction.
- ② Prepare serial I/O for transmission/reception.
- ③ Set the interrupt request bit to "0" with CLB instruction after 1 or more instruction has been executed.
- ④ Set the interrupt enable bit to "1" (enabled).

Reason

When the transmission enable bit is set to "1", the transmit buffer empty flag and transmit shift register completion flag are set to "1". The interrupt request is generated and the transmission interrupt request bit is set regardless of which of the two timings listed below is selected as the timing for the transmission interrupt to be generated.

- Transmit buffer empty flag is set to "1"
- Transmit shift register completion flag is set to "1"
- (6) Transmission control when external clock is selected (Serial I/O1 clock synchronous mode) When an external clock is used as the synchronous clock for data transmission, set the transmit

enable bit to "1" at "H" of the SCLK1 input level. Also, write the transmit data to the transmit buffer register (serial I/O shift register) at "H" of the SCLK1 input level.

(7) Transmit data writing (Serial I/O2)

In the clock synchronous serial I/O, when selecting an external clock as synchronous clock, write the transmit data to the serial I/O2 register (serial I/O shift register) at "H" of the transfer clock input level.

3.3.6 Notes on PWM

The PWM starts after the PWM enable bit is set to enable and "L" level is output from the PWM pin. The length of this "L" level output is as follows:

 $\frac{n+1}{2 \cdot f(X_{IN})}$ sec. (Count source selection bit = "0", where n is the value set in the prescaler)

 $\frac{n+1}{f(X_{IN})}$ sec. (Count source selection bit = "1", where n is the value set in the prescaler)

3.3.7 Notes on A-D converter

(1) Analog input pin

Make the signal source impedance for analog input low, or equip an analog input pin with an external capacitor of 0.01 μ F to 1 μ F. Further, be sure to verify the operation of application products on the user side.

Reason

An analog input pin includes the capacitor for analog voltage comparison. Accordingly, when signals from signal source with high impedance are input to an analog input pin, charge and discharge noise generates. This may cause the A-D conversion precision to be worse.

(2) A-D converter power source pin

The AVss pin is A-D converter power source pin. Regardless of using the A-D conversion function or not, connect it as following :

• AVss : Connect to the Vss line

Reason

If the AVss pin is opened, the microcomputer may have a failure because of noise or others.

(3) Clock frequency during A-D conversion

The comparator consists of a capacity coupling, and a charge of the capacity will be lost if the clock frequency is too low. Thus, make sure the following during an A-D conversion.

- f(XIN) is 500 kHz or more in middle-/high-speed mode.
- Do not execute the **STP** instruction.
- When the A-D converter is operated at low-speed mode, $f(X_{IN})$ do not have the lower limit of frequency, because of the A-D converter has a built-in self-oscillation circuit.

3.3.8 Notes on watchdog timer

•Make sure that the watchdog timer does not underflow while waiting Stop release, because the watchdog timer keeps counting during that term.

•When the STP instruction disable bit has been set to "1", it is impossible to switch it to "0" by a program.

3.3 Notes on use

3.3.9 Notes on RESET pin

(1) Connecting capacitor

In case where the RESET signal rise time is long, connect a ceramic capacitor or others across the RESET pin and the Vss pin. Use a 1000 pF or more capacitor for high frequency use. When connecting the capacitor, note the following :

• Make the length of the wiring which is connected to a capacitor as short as possible.

• Be sure to verify the operation of application products on the user side.

Reason

If the several nanosecond or several ten nanosecond impulse noise enters the RESET pin, it may cause a microcomputer failure.

(2) Reset release after power on

When releasing the reset after power on, such as power-on reset, release reset after X_{IN} passes more than 20 cycles in the state where the power supply voltage is 2.7 V or more and the X_{IN} oscillation is stable.

Reason

To release reset, the RESET pin must be held at an "L" level for 20 cycles or more of X_{IN} in the state where the power source voltage is between 2.7 V and 5.5 V, and X_{IN} oscillation is stable.

3.3.10 Notes on using stop mode

■Register setting

Since values of the prescaler 12 and Timer 1 are automatically reloaded when returning from the stop mode, set them again, respectively. (When the oscillation stabilizing time set after STP instruction released bit is "0")

When using the oscillation stabilizing time set after STP instruction released bit set to "1", evaluate time to stabilize oscillation of the used oscillator and set the value to the timer 1 and prescaler 12.

■Clock restoration

After restoration from the stop mode to the normal mode by an interrupt request, the contents of the CPU mode register previous to the STP instruction execution are retained. Accordingly, if both main clock and sub clock were oscillating before execution of the STP instruction, the oscillation of both clocks is resumed at restoration.

In the above case, when the main clock side is set as a system clock, the oscillation stabilizing time for approximately 8,000 cycles of the X_{IN} input is reserved at restoration from the stop mode. At this time, note that the oscillation on the sub clock side may not be stabilized even after the lapse of the oscillation stabilizing time of the main clock side.

3.3.11 Notes on wait mode

■Clock restoration

If the wait mode is released by a reset when X_{CIN} is set as the system clock and X_{IN} oscillation is stopped during execution of the WIT instruction, X_{CIN} oscillation stops, X_{IN} oscillations starts, and X_{IN} is set as the system clock.

In the above case, the RESET pin should be held at "L" until the oscillation is stabilized.

3.3.12 Notes on CPU rewrite mode of flash memory version

(1) Operation speed

During CPU rewrite mode, set the internal clock frequency 4MHz or less by using the main clock division ratio selection bits (bits 6, 7 at address $003B_{16}$).

(2) Instructions inhibited against use The instructions which refer to the internal data of the flash memory cannot be used during CPU rewrite mode.

(3) Interrupts inhibited against use

The interrupts cannot be used during CPU rewrite mode because they refer to the internal data of the flash memory.

(4) Watchdog timer

In case of the watchdog timer has been running already, the internal reset generated by watchdog timer underflow does not happen, because of watchdog timer is always clearing during program or erase operation.

(5) Reset

Reset is always valid. In case of $CNV_{SS} = "H"$ when reset is released, boot mode is active. So the program starts from the address contained in addresses FFFC₁₆ and FFFD₁₆ in boot ROM area.

3.3.13 Notes on restarting oscillation

■Restarting oscillation

Usually, when the MCU stops the clock oscillation by STP instruction and the STP instruction has been released by an external interrupt source, the fixed values of Timer 1 and Prescaler 12 (Timer $1 = "01_{16}"$, Prescaler $12 = "FF_{16}"$) are automatically reloaded in order for the oscillation to stabilize. The user can inhibit the automatic setting by writing "1" to bit 0 of MISRG (address 0038₁₆). However, by setting this bit to "1", the previous values, set just before the STP instruction was executed, will remain in Timer 1 and Prescaler 12. Therefore, you will need to set an appropriate value to each register, in accordance with the oscillation stabilizing time, before executing the STP instruction.

Reason

Oscillation will restart when an external interrupt is received. However, internal clock ϕ is supplied to the CPU only when Timer 1 starts to underflow. This ensures time for the clock oscillation using the ceramic resonators to be stabilized.

3.3 Notes on use

3.3.14 Notes on programming

(1) Processor status register

① Initializing of processor status register

Flags which affect program execution must be initialized after a reset. In particular, it is essential to initialize the T and D flags because they have an important effect on calculations.

Reason

After a reset, the contents of the processor status register (PS) are undefined except for the I flag which is "1".

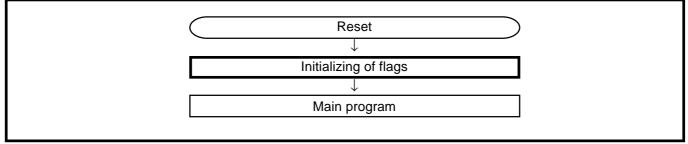
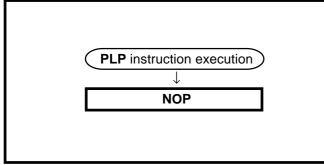


Fig. 3.3.4 Initialization of processor status register

② How to reference the processor status register

To reference the contents of the processor status register (PS), execute the **PHP** instruction once then read the contents of (S+1). If necessary, execute the **PLP** instruction to return the PS to its original status.

A NOP instruction should be executed after every PLP instruction.





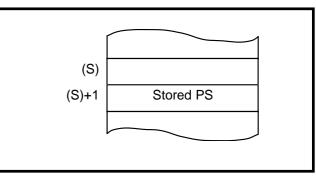


Fig. 3.3.6 Stack memory contents after PHP instruction execution

(2) BRK instruction

① Interrupt priority level

When the BRK instruction is executed with the following conditions satisfied, the interrupt execution is started from the address of interrupt vector which has the highest priority.

- Interrupt request bit and interrupt enable bit are set to "1".
- Interrupt disable flag (I) is set to "1" to disable interrupt.

(3) Decimal calculations

① Execution of decimal calculations

The **ADC** and **SBC** are the only instructions which will yield proper decimal notation, set the decimal mode flag (D) to "1" with the **SED** instruction. After executing the **ADC** or **SBC** instruction, execute another instruction before executing the **SEC**, **CLC**, or **CLD** instruction.

2 Notes on status flag in decimal mode

When decimal mode is selected, the values of three of the flags in the status register (the N, V, and Z flags) are invalid after a **ADC** or **SBC** instruction is executed.

The carry flag (C) is set to "1" if a carry is generated as a result of the calculation, or is cleared to "0" if a borrow is generated. To determine whether a calculation has generated a carry, the C flag must be initialized to "0" before each calculation. To check for a borrow, the C flag must be initialized to "1" before each calculation.

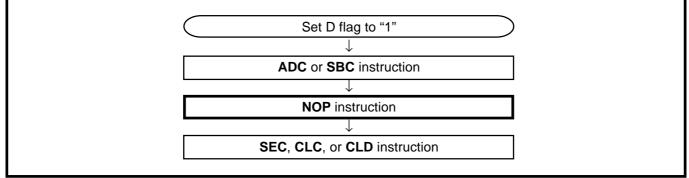


Fig. 3.3.7 Status flag at decimal calculations

(4) JMP instruction

When using the **JMP** instruction in indirect addressing mode, do not specify the last address on a page as an indirect address.

(5) Multiplication and Division Instructions

- The index X mode (T) and the decimal mode (D) flags do not affect the MUL and DIV instruction.
- The execution of these instructions does not change the contents of the processor status register.

(6) Ports

The contents of the port direction registers cannot be read. The following cannot be used:

- The data transfer instruction (LDA, etc.)
- The operation instruction when the index X mode flag (T) is "1"
- The addressing mode which uses the value of a direction register as an index
- The bit-test instruction (BBC or BBS, etc.) to a direction register

• The read-modify-write instructions (ROR, CLB, or SEB, etc.) to a direction register.

Use instructions such as LDM and STA, etc., to set the port direction registers.

(7) Instruction Execution Time

The instruction execution time is obtained by multiplying the frequency of the internal clock ϕ by the number of cycles needed to execute an instruction.

The number of cycles required to execute an instruction is shown in the list of machine instructions. The frequency of the internal clock ϕ is half of the X_{IN} frequency in high-speed mode.

3.3 Notes on use

3.3.15 EPROM Version/One Time PROM Version/Flash Memory Version

The CNVss pin is connected to the internal memory circuit block by a low-ohmic resistance, since it has the multiplexed function to be a programmable power source pin (V_{PP} pin) as well.

To improve the noise reduction, connect a track between CNVss pin and Vss pin or Vcc pin with 1 to 10 $k\Omega$ resistance.

The mask ROM version track of CNVss pin has no operational interference even if it is connected to Vss pin or Vcc pin via a resistor.

3.3.16 Handling of Source Pins

In order to avoid a latch-up occurrence, connect a capacitor suitable for high frequencies as bypass capacitor between power source pin (V_{cc} pin) and GND pin (V_{ss} pin) and between power source pin (V_{cc} pin) and analog power source input pin (AV_{ss} pin). Besides, connect the capacitor to as close as possible. For bypass capacitor which should not be located too far from the pins to be connected, a ceramic capacitor of 0.01 μ F-0.1 μ F is recommended.

3.3.17 Differences between 3850 group (standard) and 3850 group (spec. H)

- (1) The absolute maximum ratings of 3850 group (spec. H) is smaller than that of 3850 group (standard).
 •Power source voltage Vcc = 0.3 to 6.5 V
 - •CNVss input voltage V = -0.3 to Vcc +0.3 V
- (2) The oscillation circuit constants of XIN-XOUT, XCIN-XCOUT may be some differences between 3850 group (standard) and 3850 group (spec. H).
- (3) Do not write any data to the reserved area and the reserved bit. (Do not change the contents after reset.)
- (4) Fix bit 3 of the CPU mode register to "1".
- (5) Be sure to perform the termination of unused pins.

3.4 Countermeasures against noise

3.4 Countermeasures against noise

Countermeasures against noise are described below. The following countermeasures are effective against noise in theory, however, it is necessary not only to take measures as follows but to evaluate before actual use.

3.4.1 Shortest wiring length

The wiring on a printed circuit board can function as an antenna which feeds noise into the microcomputer. The shorter the total wiring length (by mm unit), the less the possibility of noise insertion into a microcomputer.

(1) Package

Select the smallest possible package to make the total wiring length short.

Reason

The wiring length depends on a microcomputer package. Use of a small package, for example QFP and not DIP, makes the total wiring length short to reduce influence of noise.

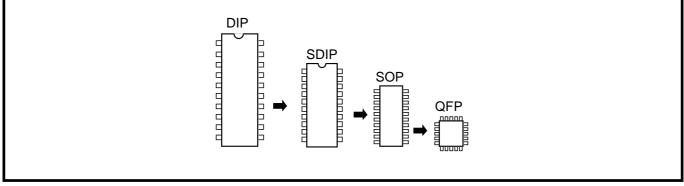


Fig. 3.4.1 Selection of packages

(2) Wiring for RESET pin

Make the length of wiring which is connected to the RESET pin as short as possible. Especially, connect a capacitor across the RESET pin and the Vss pin with the shortest possible wiring (within 20mm).

Reason

The width of a pulse input into the RESET pin is determined by the timing necessary conditions. If noise having a shorter pulse width than the standard is input to the RESET pin, the reset is released before the internal state of the microcomputer is completely initialized. This may cause a program runaway.

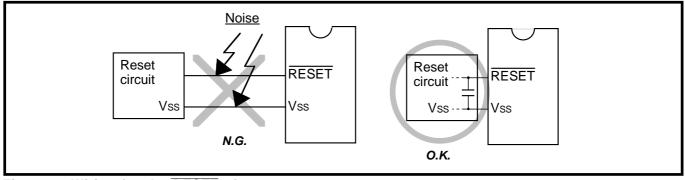


Fig. 3.4.2 Wiring for the RESET pin

3.4 Countermeasures against noise

(3) Wiring for clock input/output pins

- Make the length of wiring which is connected to clock I/O pins as short as possible.
- Make the length of wiring (within 20 mm) across the grounding lead of a capacitor which is connected to an oscillator and the Vss pin of a microcomputer as short as possible.
- Separate the Vss pattern only for oscillation from other Vss patterns.

Reason

If noise enters clock I/O pins, clock waveforms may be deformed. This may cause a program failure or program runaway. Also, if a potential difference is caused by the noise between the Vss level of a microcomputer and the Vss level of an oscillator, the correct clock will not be input in the microcomputer.

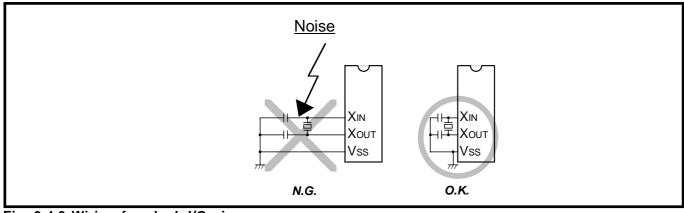


Fig. 3.4.3 Wiring for clock I/O pins

(4) Wiring to CNVss pin

Connect the CNVss pin to the Vss pin with the shortest possible wiring.

Reason

The processor mode of a microcomputer is influenced by a potential at the CNVss pin. If a potential difference is caused by the noise between pins CNVss and Vss, the processor mode may become unstable. This may cause a microcomputer malfunction or a program runaway.

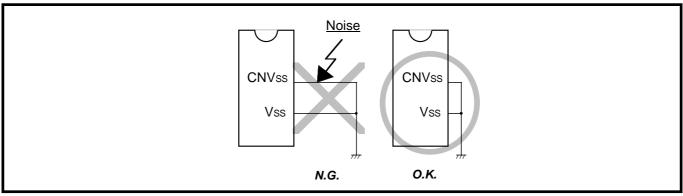


Fig. 3.4.4 Wiring for CNVss pin

(5) Wiring to VPP pin of One Time PROM version, EPROM version, and Flash memory version Connect an approximately 5 k Ω resistor to the VPP pin the shortest possible in series and also to the Vss pin. When not connecting the resistor, make the length of wiring between the VPP pin and the Vss pin the shortest possible.

Note: Even when a circuit which included an approximately 5 k Ω resistor is used in the Mask ROM version, the microcomputer operates correctly.

Reason

The VPP pin of the One Time PROM, the EPROM version, and the flash memory version is the power source input pin for the built-in PROM. When programming in the built-in PROM, the impedance of the VPP pin is low to allow the electric current for writing flow into the PROM. Because of this, noise can enter easily. If noise enters the VPP pin, abnormal instruction codes or data are read from the built-in PROM, which may cause a program runaway.

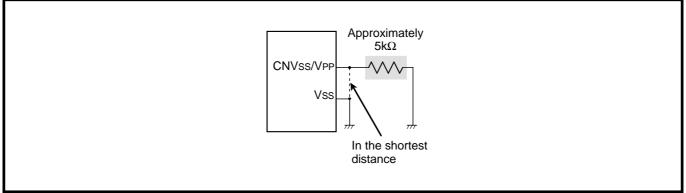


Fig. 3.4.5 Wiring for the VPP pin of the One Time PROM version, the EPROM version, and the flash memory version

3.4.2 Connection of bypass capacitor across Vss line and Vcc line

Connect an approximately 0.1 μ F bypass capacitor across the Vss line and the Vcc line as follows:

- Connect a bypass capacitor across the Vss pin and the Vcc pin at equal length.
- Connect a bypass capacitor across the Vss pin and the Vcc pin with the shortest possible wiring.
- Use lines with a larger diameter than other signal lines for Vss line and Vcc line.
- Connect the power source wiring via a bypass capacitor to the Vss pin and the Vcc pin.

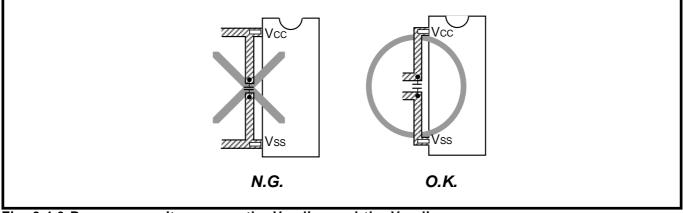


Fig. 3.4.6 Bypass capacitor across the Vss line and the Vcc line

3.4 Countermeasures against noise

3.4.3 Wiring to analog input pins

- Connect an approximately 100 Ω to 1 k Ω resistor to an analog signal line which is connected to an analog input pin in series. Besides, connect the resistor to the microcomputer as close as possible.
- Connect an approximately 1000 pF capacitor across the Vss pin and the analog input pin. Besides, connect the capacitor to the Vss pin as close as possible. Also, connect the capacitor across the analog input pin and the Vss pin at equal length.

Reason

Signals which is input in an analog input pin (such as an A-D converter/comparator input pin) are usually output signals from sensor. The sensor which detects a change of event is installed far from the printed circuit board with a microcomputer, the wiring to an analog input pin is longer necessarily. This long wiring functions as an antenna which feeds noise into the microcomputer, which causes noise to an analog input pin.

If a capacitor between an analog input pin and the V_{SS} pin is grounded at a position far away from the V_{SS} pin, noise on the GND line may enter a microcomputer through the capacitor.

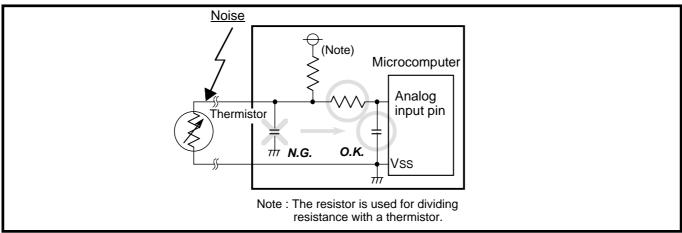


Fig. 3.4.7 Analog signal line and a resistor and a capacitor

3.4.4 Oscillator concerns

Take care to prevent an oscillator that generates clocks for a microcomputer operation from being affected by other signals.

(1) Keeping oscillator away from large current signal lines

Install a microcomputer (and especially an oscillator) as far as possible from signal lines where a current larger than the tolerance of current value flows.

Reason

In the system using a microcomputer, there are signal lines for controlling motors, LEDs, and thermal heads or others. When a large current flows through those signal lines, strong noise occurs because of mutual inductance.

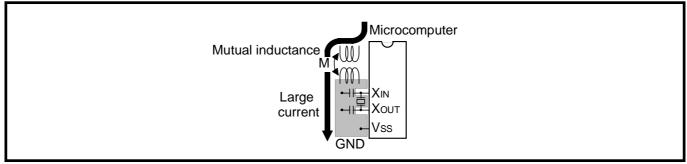


Fig. 3.4.8 Wiring for a large current signal line

- (2) Installing oscillator away from signal lines where potential levels change frequently Install an oscillator and a connecting pattern of an oscillator away from signal lines where potential levels change frequently. Also, do not cross such signal lines over the clock lines or the signal lines which are sensitive to noise.
 - Reason

Signal lines where potential levels change frequently (such as the CNTR pin signal line) may affect other lines at signal rising edge or falling edge. If such lines cross over a clock line, clock waveforms may be deformed, which causes a microcomputer failure or a program runaway.

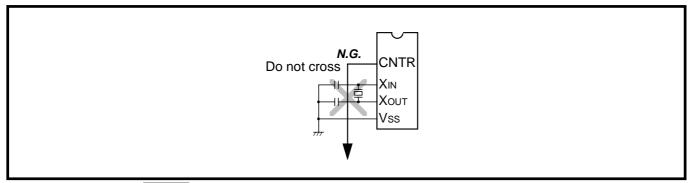


Fig. 3.4.9 Wiring of RESET pin

3.4 Countermeasures against noise

(3) Oscillator protection using Vss pattern

As for a two-sided printed circuit board, print a VSS pattern on the underside (soldering side) of the position (on the component side) where an oscillator is mounted.

Connect the Vss pattern to the microcomputer Vss pin with the shortest possible wiring. Besides, separate this Vss pattern from other Vss patterns.

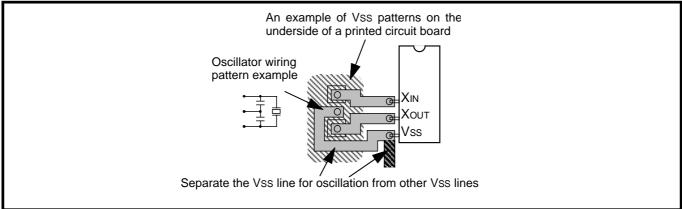


Fig. 3.4.10 Vss pattern on the underside of an oscillator

3.4.5 Setup for I/O ports

Setup I/O ports using hardware and software as follows:

<Hardware>

- Connect a resistor of 100 Ω or more to an I/O port in series.

<Software>

- As for an input port, read data several times by a program for checking whether input levels are equal or not.
- As for an output port, since the output data may reverse because of noise, rewrite data to its port latch at fixed periods.
- Rewrite data to direction registers at fixed periods.
- **Note:** When a direction register is set for <u>input port</u> again at fixed periods, a several-nanosecond short pulse may be output from this port. If this is undesirable, connect a capacitor to this port to remove the noise pulse.

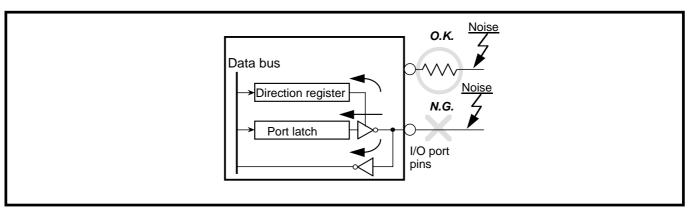


Fig. 3.4.11 Setup for I/O ports

3.4 Countermeasures against noise

3.4.6 Providing of watchdog timer function by software

If a microcomputer runs away because of noise or others, it can be detected by a software watchdog timer and the microcomputer can be reset to normal operation. This is equal to or more effective than program runaway detection by a hardware watchdog timer. The following shows an example of a watchdog timer provided by software.

In the following example, to reset a microcomputer to normal operation, the main routine detects errors of the interrupt processing routine and the interrupt processing routine detects errors of the main routine. This example assumes that interrupt processing is repeated multiple times in a single main routine processing.

<The main routine>

• Assigns a single byte of RAM to a software watchdog timer (SWDT) and writes the initial value N in the SWDT once at each execution of the main routine. The initial value N should satisfy the following condition:

 $N+1 \ge$ (Counts of interrupt processing executed in each main routine)

As the main routine execution cycle may change because of an interrupt processing or others, the initial value N should have a margin.

- Watches the operation of the interrupt processing routine by comparing the SWDT contents with counts of interrupt processing after the initial value N has been set.
- Detects that the interrupt processing routine has failed and determines to branch to the program initialization routine for recovery processing in the following case:

If the SWDT contents do not change after interrupt processing.

<The interrupt processing routine>

- Decrements the SWDT contents by 1 at each interrupt processing.
- Determines that the main routine operates normally when the SWDT contents are reset to the initial value N at almost fixed cycles (at the fixed interrupt processing count).
- Detects that the main routine has failed and determines to branch to the program initialization routine for recovery processing in the following case:

If the SWDT contents are not initialized to the initial value N but continued to decrement and if they reach 0 or less.

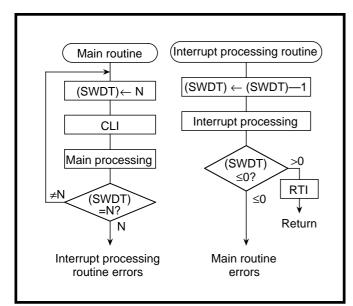
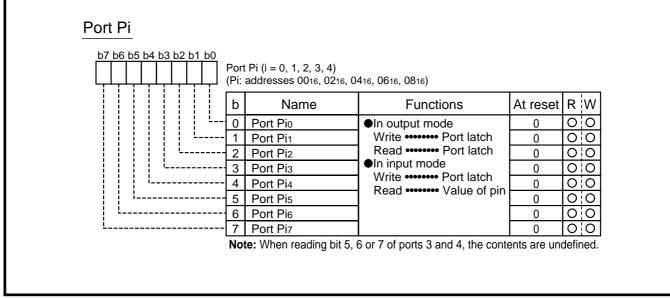
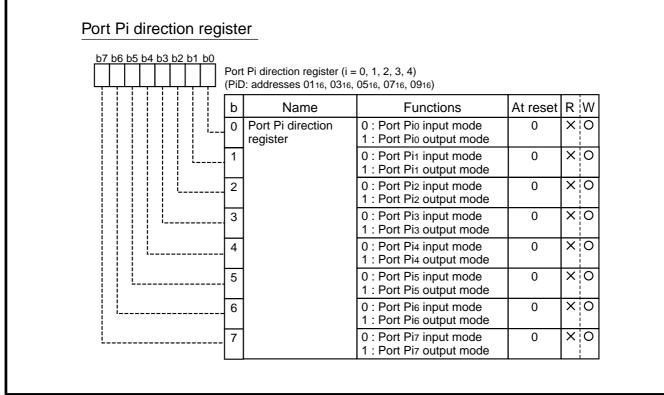


Fig. 3.4.12 Watchdog timer by software

3.5 List of registers









3.5 List of registers

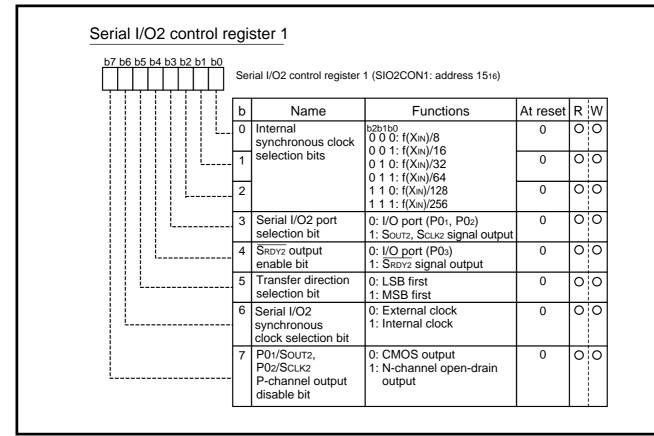


Fig. 3.5.3 Structure of Serial I/O2 control register 1

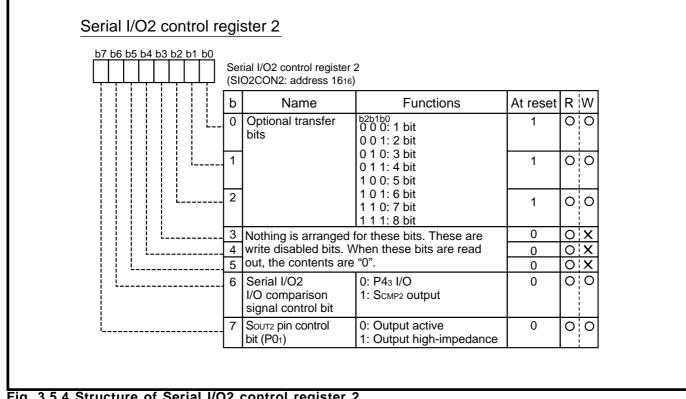
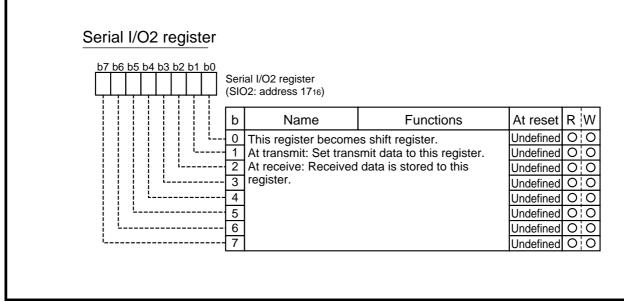
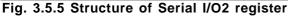


Fig. 3.5.4 Structure of Serial I/O2 control register 2





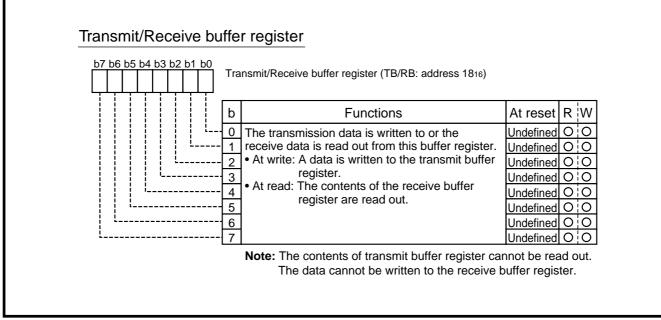
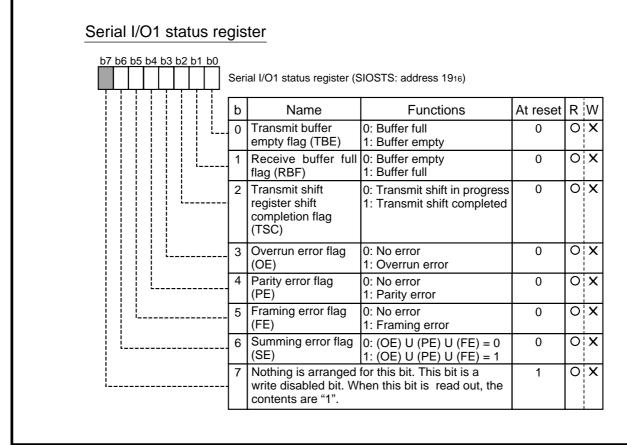
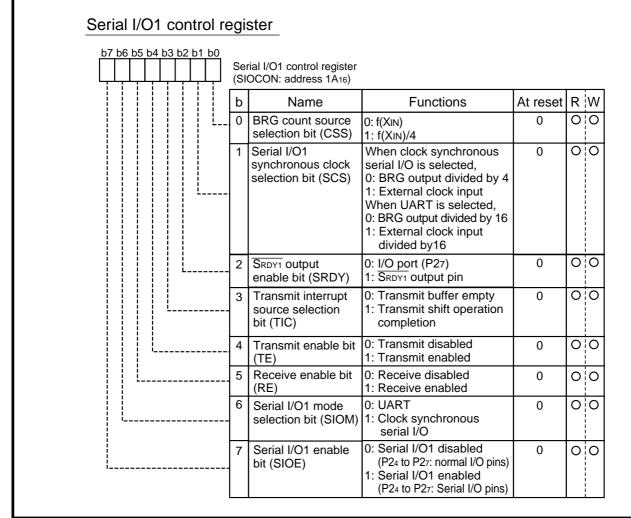


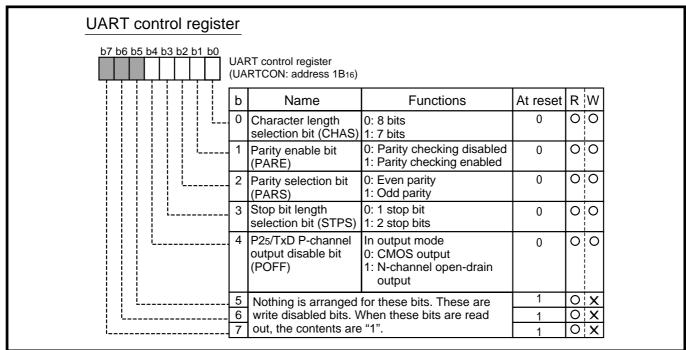
Fig. 3.5.6 Structure of Transmit/Receive buffer register



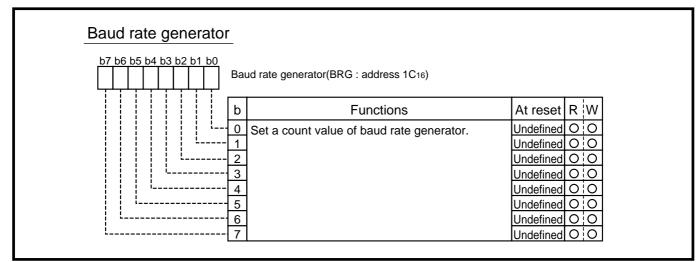


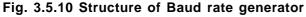


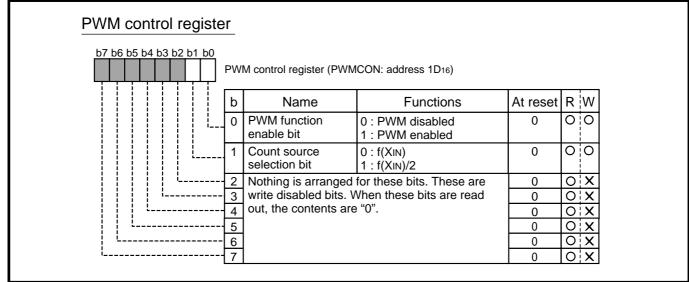


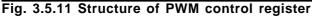












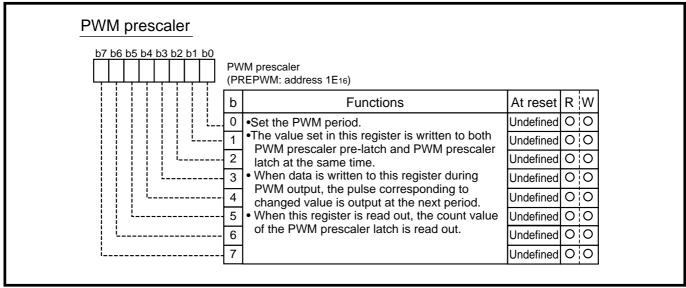


Fig. 3.5.12 Structure of PWM prescaler

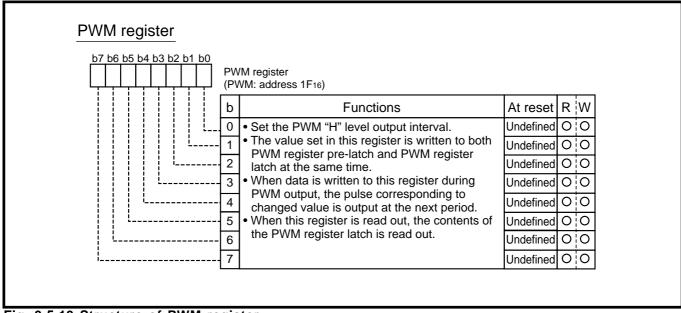


Fig. 3.5.13 Structure of PWM register

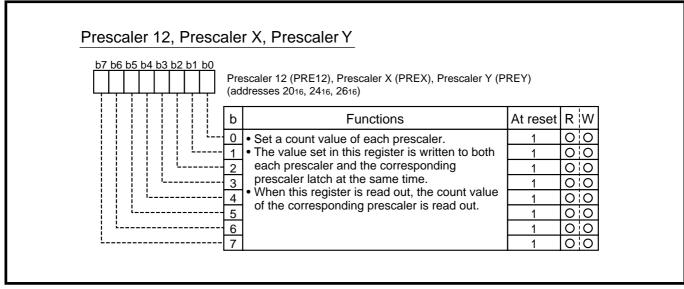
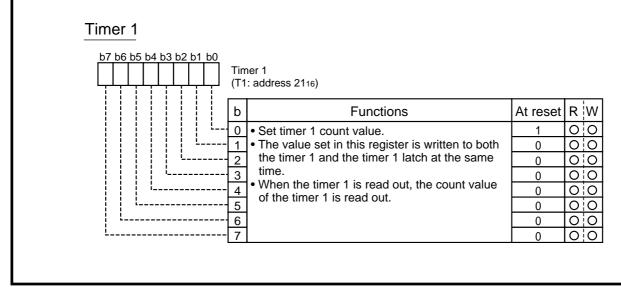
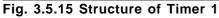


Fig. 3.5.14 Structure of Prescaler 12, Prescaler X, Prescaler Y





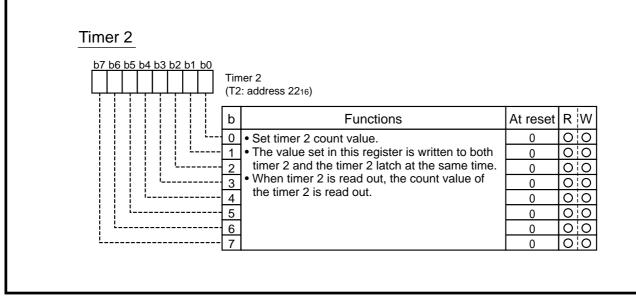


Fig. 3.5.16 Structure of Timer 2

3.5 List of registers

Timer XY mode register

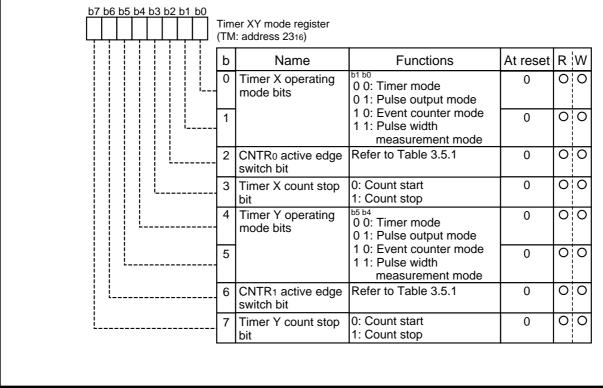
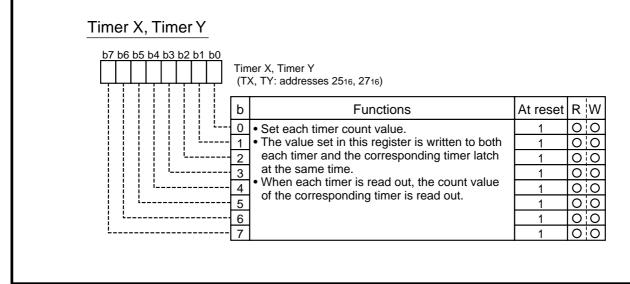


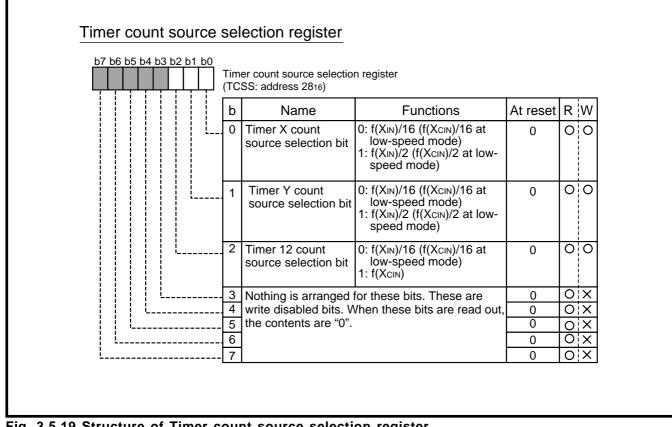
Fig. 3.5.17 Structure of Timer XY mode register

Table 3.5.1 CNTR₀/CNTR₁ active edge switch bit function

Timer X /Timer Y	Set	Timer function	CNTR ₀ / CNTR ₁ interrupt request
operation modes	value		occurrence source
Timer mode	"0"	No influence to timer count	CNTR ₀ /CNTR ₁ input signal falling edge
	"1"	No influence to timer count	CNTR ₀ /CNTR ₁ input signal rising edge
Pulse output	"0"	Pulse output start: Beginning	Output signal falling edge count
mode		at "H" level	
	"1"	Pulse output start: Beginning	Output signal rising edge count
		at "L" level	
Event counter	"0"	Rising edge count	Input signal falling edge count
mode	"1"	Falling edge count	Input signal rising edge count
Pulse width	"0"	"H" level width measurement	Input signal falling edge count
measurement mode	"1"	"L" level width measurement	Input signal rising edge count









3.5 List of registers

A-D control register b7 b6 b5 b4 b3 b2 b1 b0 A-D control register (ADCON: address 3416) RW b At reset Name **Functions** 0 00 0 Analog input pin b2 b1 b0 selection bits 0 0 0: P30/AN0 0 0 1: P31/AN1 1 0 0 0 0 1 0: P32/AN2 0 1 1: P33/AN3 0 010 2 1 0 0: P34/AN4 3 Nothing is arranged for this bit. This is a write 0 0 × disabled bit. When this bit is read out, the contents are "0". 0: Conversion in progress 4 AD conversion 1 010 1: Conversion completed completion bit 0 5 Nothing is arranged for these bits. These are Ο Х write disabled bits. When these bist are read 6 0 0 Х out, the contents are "0". 7 O 0 ×

Fig. 3.5.20 Structure of A-D control register

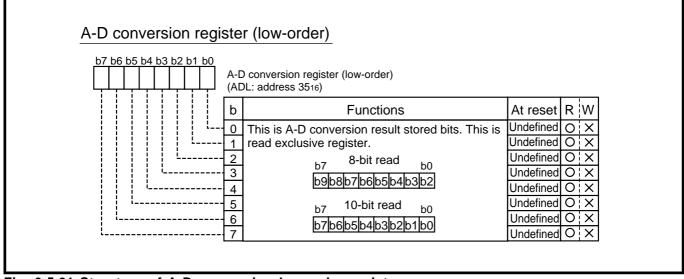


Fig. 3.5.21 Structure of A-D conversion low-order register

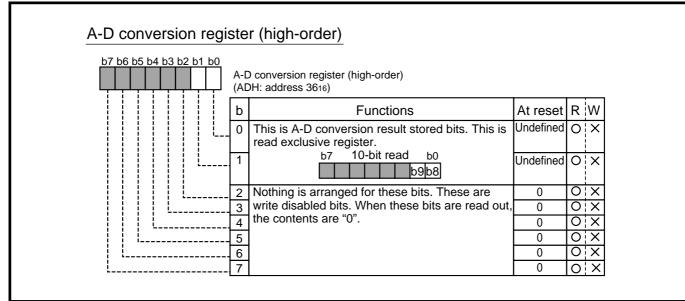
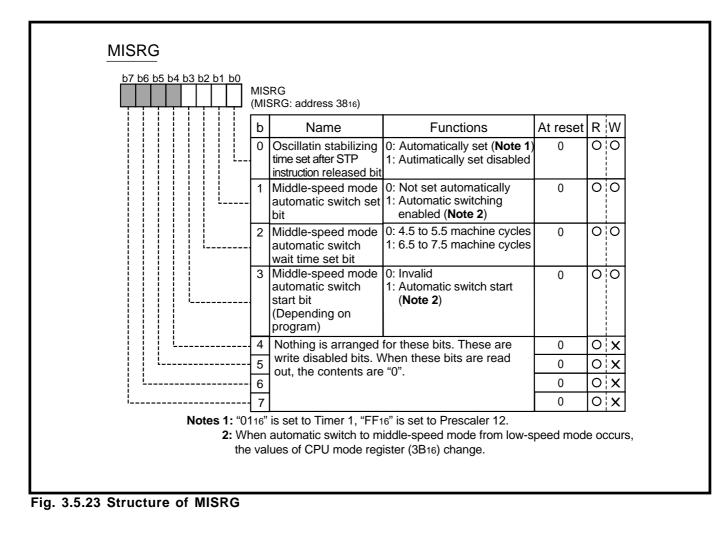
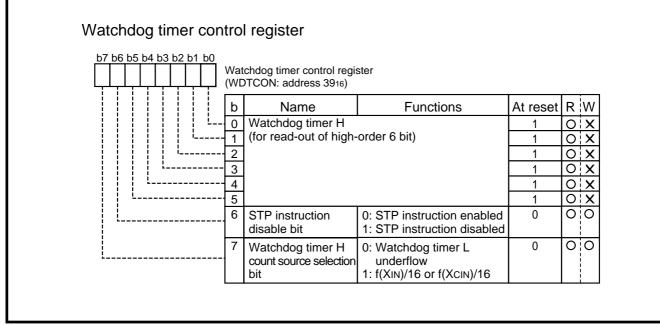
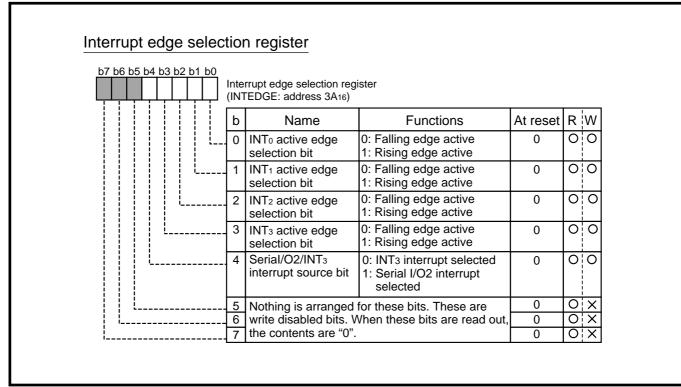


Fig. 3.5.22 Structure of A-D conversion high-order register











3.5 List of registers

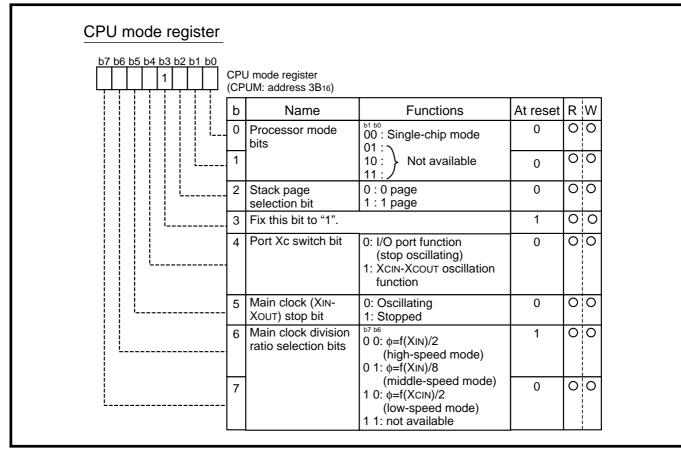
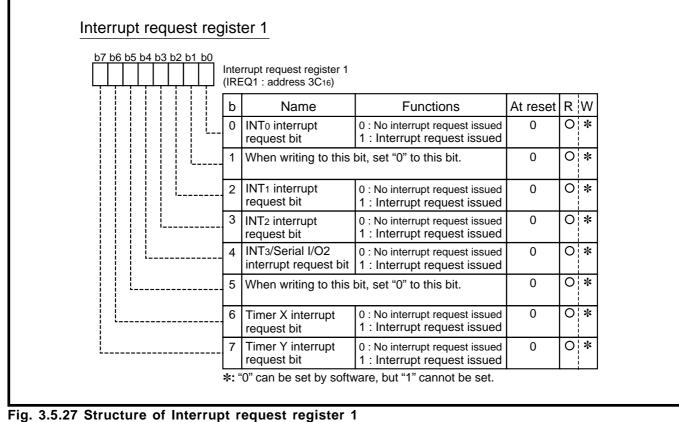
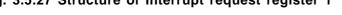
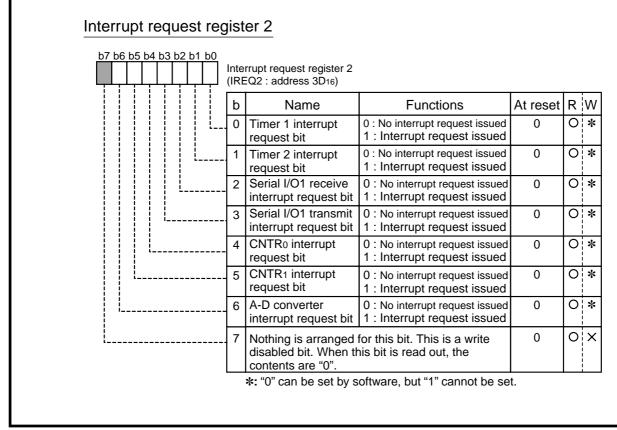


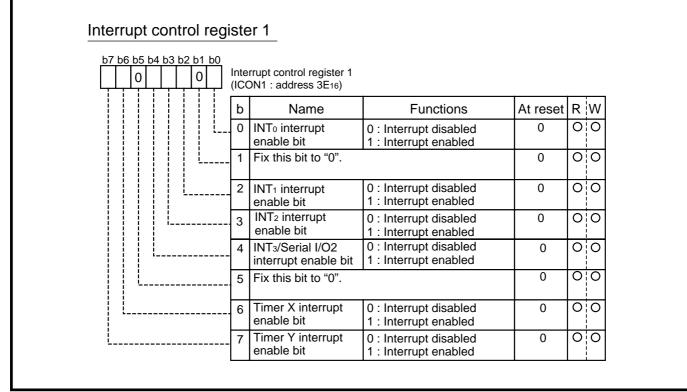
Fig. 3.5.26 Structure of CPU mode register













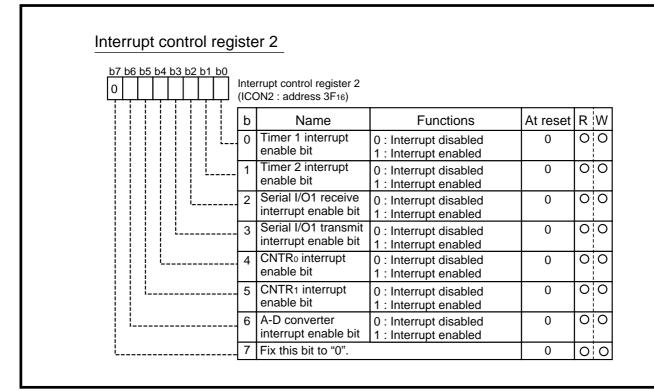
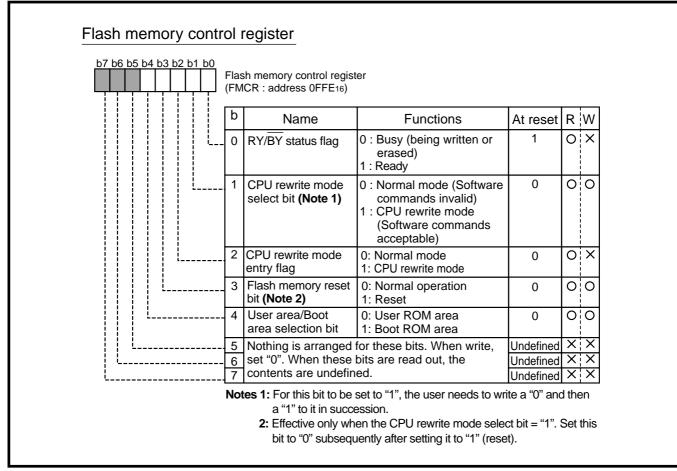
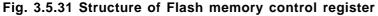


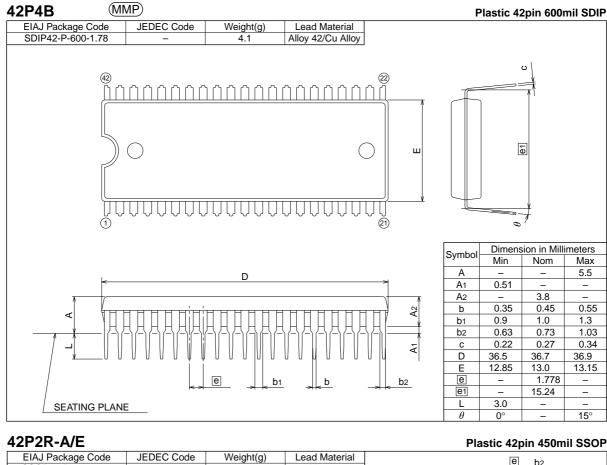
Fig. 3.5.30 Structure of Interrupt control register 2



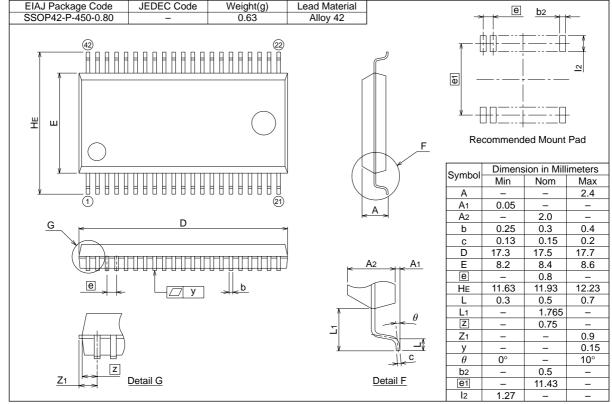


3.6 Package outline

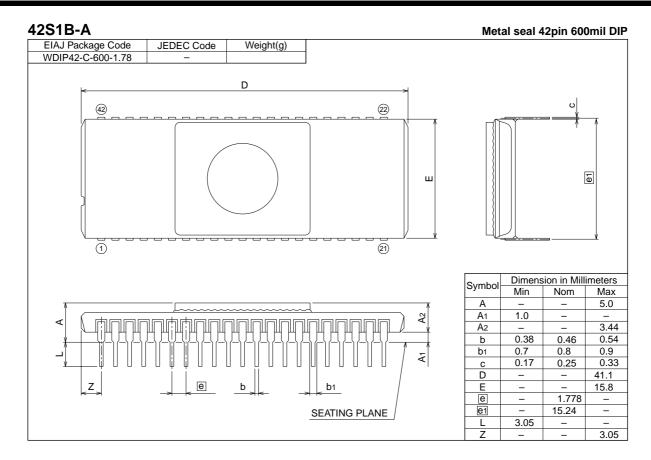
3.6 Package outline







3.6 Package outline



3.7 Machine instructions

3.7 Machine instructions

									_ A	ddr	essi	ng r	noc	le						
Symbol	Function	Details		IMP			IMN	1		А		ВΙ	Г, А	., R		ZP		BIT	, ZP	', R
			OP	n	#	OF	n	#	OP	n	#	OP	n	#	OP	n	#	OP	n	#
ADC (Note 1) (Note 5)	When T = 0 $A \leftarrow A + M + C$ When T = 1 $M(X) \leftarrow M(X) + M + C$	When T = 0, this instruction adds the contents M, C, and A; and stores the results in A and C. When T = 1, this instruction adds the contents of M(X), M and C; and stores the results in M(X) and C. When T=1, the contents of A remain unchanged, but the contents of status flags are changed. M(X) represents the contents of memory where is indicated by X.				69	2	2							65	3	2			
AND (Note 1)	When T = 0 A \leftarrow A \land M When T = 1 M(X) \leftarrow M(X) \land M	When T = 0, this instruction transfers the con- tents of A and M to the ALU which performs a bit-wise AND operation and stores the result back in A. When T = 1, this instruction transfers the con- tents $M(X)$ and M to the ALU which performs a bit-wise AND operation and stores the results back in $M(X)$. When T = 1 the contents of A re- main unchanged, but status flags are changed. M(X) represents the contents of memory where is indicated by X.				29	2	2							25	3	2			
ASL	7 0 [℃←□←0	This instruction shifts the content of A or M by one bit to the left, with bit 0 always being set to 0 and bit 7 of A or M always being contained in C.							0A	2	1				06	5	2			
BBC (Note 4)	Ai or Mi = 0?	This instruction tests the designated bit i of M or A and takes a branch if the bit is 0. The branch address is specified by a relative address. If the bit is 1, next instruction is executed.										1 <u>3</u> 20i	4	2				17 20i	5	3
BBS (Note 4)	Ai or Mi = 1?	This instruction tests the designated bit i of the M or A and takes a branch if the bit is 1. The branch address is specified by a relative address. If the bit is 0, next instruction is executed.										03 20i	4	2				07 20i	5	3
BCC (Note 4)	C = 0?	This instruction takes a branch to the ap- pointed address if C is 0. The branch address is specified by a relative address. If C is 1, the next instruction is executed.																		
BCS (Note 4)	C = 1?	This instruction takes a branch to the ap- pointed address if C is 1. The branch address is specified by a relative address. If C is 0, the next instruction is executed.																		
BEQ (Note 4)	Z = 1?	This instruction takes a branch to the ap- pointed address when Z is 1. The branch address is specified by a relative address. If Z is 0, the next instruction is executed.																		
BIT	A A M	This instruction takes a bit-wise logical AND of A and M contents; however, the contents of A and M are not modified. The contents of N, V, Z are changed, but the contents of A, M remain unchanged.													24	3	2			
BMI (Note 4)	N = 1?	This instruction takes a branch to the ap- pointed address when N is 1. The branch address is specified by a relative address. If N is 0, the next instruction is executed.																		
BNE (Note 4)	Z = 0?	This instruction takes a branch to the ap- pointed address if Z is 0. The branch address is specified by a relative address. If Z is 1, the next instruction is executed.																		

															Ad	Idres	sin	g m	ode															F	Proc	esso	or st	atus	s reg	giste	er
z	ΣP,	Х	Τ	Z	P, ۱	(ABS	3	A	BS,	Х	A	BS,		1	IND		1	P, IN	١D	IN	۱D,	х	IN	٧D,	Y		REL	-		SP		7	1	5	4	3	2	1	0
OP	1	1	0	- 1			OP	n	#	OP	n	#	OP	n	#	OP	n	#	OP	n	#	OP	n	#	OP	n	#	OP	n	#	OP	n	#	N	v	т	в	D	1	z	с
75	4	2					6D	4	3	7D	5	3	79	5	3							61	6	2	71	6	2							N	V	•	•	•	•	Z	С
35	4	2					2D	4	3	3D	5	3	39	5	3							21	6	2	31	6	2							N	•	•	•	•	•	Z	•
16	6	2					0E	6	3	1E	7	3																						N	•	•	•	•	•	Z	С
																																		•	•	•	•	•	•	•	•
																																		•	•	•	•	•	•	•	•
																												90	2	2				•	•	•	•	•	•	•	•
																												В0	2	2				•	•	•	•	•	•	•	•
																												F0	2	2				•	•	•	•	•	•	•	•
							2C	4	3																									M7	M6	•	•	•	•	z	•
																												30	2	2				•	•	•	•	•	•	•	•
																												D0	2	2				•	•	•	•	•	•	•	•

APPENDIX

3.7 Machine instructions

						_			-	Addı	res	sing	mo	de	-					
Symbol	Function	Details		IMF)		IM	Л		А			BIT,	А		ZP		Bľ	т, z	ΈP
			OP	n	#	0	P n	#	OF	'n	#	0	P n	#	OP	n	#	OP	n	#
BPL (Note 4)	N = 0?	This instruction takes a branch to the ap- pointed address if N is 0. The branch address is specified by a relative address. If N is 1, the next instruction is executed.																		
BRA	$PC \gets PC \pm offset$	This instruction branches to the appointed ad- dress. The branch address is specified by a relative address.																		
BRK	$\begin{array}{l} B \leftarrow 1 \\ (PC) \leftarrow (PC) + 2 \\ M(S) \leftarrow PCH \\ S \leftarrow S - 1 \\ M(S) \leftarrow PCL \\ S \leftarrow S - 1 \\ M(S) \leftarrow PS \\ S \leftarrow S - 1 \\ I \leftarrow 1 \\ PCL \leftarrow ADL \\ PCH \leftarrow ADH \end{array}$	When the BRK instruction is executed, the CPU pushes the current PC contents onto the stack. The BADRS designated in the interrupt vector table is stored into the PC.	00	7	1															
BVC (Note 4)	V = 0?	This instruction takes a branch to the appointed address if V is 0. The branch address is specified by a relative address. If V is 1, the next instruction is executed.																		
BVS (Note 4)	V = 1?	This instruction takes a branch to the ap- pointed address when V is 1. The branch address is specified by a relative address. When V is 0, the next instruction is executed.																		
CLB	Ai or Mi ← 0	This instruction clears the designated bit i of A or M.										1E 20	3 2 Ii	1				1F 20i	5	2
CLC	C ← 0	This instruction clears C.	18	2	1															
CLD	D ← 0	This instruction clears D.	D8	2	1															
CLI	I ← 0	This instruction clears I.	58	2	1															
CLT	T ← 0	This instruction clears T.	12	2	1															
CLV	$V \leftarrow 0$	This instruction clears V.	B8	2	1															
CMP (Note 3)	When T = 0 A – M When T = 1 M(X) – M	When T = 0, this instruction subtracts the con- tents of M from the contents of A. The result is not stored and the contents of A or M are not modified.When T = 1, the CMP subtracts the contents of M from the contents of M(X). The result is not stored and the contents of X, M, and A are not modified.M(X) represents the contents of memory where is indicated by X.				C	92	2							C5	3	2			
СОМ	$M \gets M$	This instruction takes the one's complement of the contents of M and stores the result in M.													44	5	2			
СРХ	X – M	This instruction subtracts the contents of M from the contents of X. The result is not stored and the contents of X and M are not modified.				E) 2	2							E4	3	2			
CPY	Y – M	This instruction subtracts the contents of M from the contents of Y. The result is not stored and the contents of Y and M are not modified.				C) 2	2							C4	3	2			
DEC	$A \leftarrow A - 1 \text{ or}$ $M \leftarrow M - 1$	This instruction subtracts 1 from the contents of A or M.							14	2	1				C6	5	2			

														Ad	dres	sing	g ma	ode															F	Proc	esso	or st	atus	reg	giste	er
z	ZP, 2	x	z	۲, ۲	Y		ABS	3	A	BS,	Х	A	BS,	Y		IND		ZF	P, IN	ID	11	۱D,	Х	11	ND,	Y		REL			SP		7	6	5	4	3	2	1	0
OP	n	#	OP	n	#	OP	n	#	OP	n	#	OP	n	#	OP	n	#	OP	n	#	OP	n	#	OP	n	#	OP	n	#	OP	n	#	N	v	т	в	D	I	z	С
																											10	2	2				•	•	•	•	•	•	•	•
																											80	4	2				•	•	•	•	•	•	•	•
																																	•	•	•	1	•	1	•	•
																											50	2	2				•	•	•	•	•	•	•	•
																											70	2	2				•	•	•	•	•	•	•	•
																																	•	•	•	•	•	•	•	•
																																	•	•	•	•	•	•	•	0
																																	•	•	•	•	•	•	•	•
																																	•	•	0	•	•	•	•	•
																																	•	0	•	•	•	•	•	•
D5	4	2				CD	4	3	DD	5	3	D9	5	3							C1	6	2	D1	6	2							N	•	•	•	•	•	z	С
																																	N	•	•	•	•	•	z	•
						EC	4	3																									N	•	•	•	•	•	z	С
						сс	4	3																									N	•	•	•	•	•	Z	с
D6	6	2				CE	6	3	DE	7	3																						N	•	•	•	•	•	Z	•

APPENDIX

3.7 Machine instructions

									A	ddr	ess	ing ı	mod	е				-		
Symbol	Function	Details		IMP	1		IMN			A		В	BIT,	A		ZP		BI	Г, Z	P
			OP		#	OP	n	#	OP	n	#	OP	n	#	OP	n	#	OP	n	#
DEX	X ← X − 1	This instruction subtracts one from the current contents of X.	CA	2	1															
DEY	Y ← Y − 1	This instruction subtracts one from the current contents of Y.	88	2	1															
DIV	$\begin{array}{l} A \leftarrow (M(zz + X + 1), \\ M(zz + X)) \ / \ A \\ M(S) \leftarrow one's \ complement \ of \ Remainder \\ S \leftarrow S - 1 \end{array}$	Divides the 16-bit data in $M(zz+(X))$ (low-order byte) and $M(zz+(X)+1)$ (high-order byte) by the contents of A. The quotient is stored in A and the one's complement of the remainder is pushed onto the stack.																		
EOR (Note 1)	When T = 0 $A \leftarrow A \forall M$ When T = 1 $M(X) \leftarrow M(X) \forall M$	When T = 0, this instruction transfers the con- tents of the M and A to the ALU which performs a bit-wise Exclusive OR, and stores the result in A. When T = 1, the contents of M(X) and M are transferred to the ALU, which performs a bit- wise Exclusive OR and stores the results in M(X). The contents of A remain unchanged, but status flags are changed. M(X) represents the contents of memory where is indicated by X.				49	2	2							45	3	2			
INC	$A \leftarrow A + 1 \text{ or}$ $M \leftarrow M + 1$	This instruction adds one to the contents of A or M.							ЗA	2	1				E6	5	2			
INX	X ← X + 1	This instruction adds one to the contents of X.	E8	2	1															
INY	Y ← Y + 1	This instruction adds one to the contents of Y.	C8	2	1															
JMP	If addressing mode is ABS $PCL \leftarrow ADL$ $PCH \leftarrow ADH$ If addressing mode is IND $PCL \leftarrow M (ADH, ADL)$ $PCH \leftarrow M (ADH, ADL + 1)$ If addressing mode is ZP, IND $PCL \leftarrow M(00, ADL)$ $PCH \leftarrow M(00, ADL + 1)$	This instruction jumps to the address desig- nated by the following three addressing modes: Absolute Indirect Absolute Zero Page Indirect Absolute																		
JSR	$\begin{array}{l} M(S) \leftarrow PCH \\ S \leftarrow S-1 \\ M(S) \leftarrow PCL \\ S \leftarrow S-1 \\ After \ executing the above, \\ if \ addressing mode is \ ABS, \\ PCL \leftarrow ADL \\ PCH \leftarrow ADH \\ if \ addressing mode is \ SP, \\ PCL \leftarrow ADL \\ PCH \leftarrow FF \\ If \ addressing mode is \ ZP, \ IND, \\ PCL \leftarrow M(00, \ ADL) \\ PCH \leftarrow M(00, \ ADL + 1) \end{array}$	This instruction stores the contents of the PC in the stack, then jumps to the address desig- nated by the following addressing modes: Absolute Special Page Zero Page Indirect Absolute																		
LDA (Note 2)	When T = 0 $A \leftarrow M$ When T = 1 $M(X) \leftarrow M$	When T = 0, this instruction transfers the con- tents of M to A. When T = 1, this instruction transfers the con- tents of M to (M(X)). The contents of A remain unchanged, but status flags are changed. M(X) represents the contents of memory where is indicated by X.				A9	2	2							A5	3	2			
LDM	M ← nn	This instruction loads the immediate value in M.													зC	4	3			
LDX	$X \leftarrow M$	This instruction loads the contents of M in X.				A2	2	2							A6	3	2			L
LDY	$Y \leftarrow M$	This instruction loads the contents of M in Y.				A0	2	2							A4	3	2			

						_						_		Ad	dres	sing	g m	ode									-			_			P	roc	esso	or st	atus	s reg	giste	r
Z	P,)	<	z	ΣΡ, `	Y	,	ABS	6	A	BS,	х	A	BS,	Y		IND		ZF	P, IN	ID	١N	ID, İ	х	١N	۱D, ۲	Y	1	REL	-		SP		7	6	5	4	3	2	1	0
ΟP	n	#	OP	n	#	OP	n	#	OP	n	#	OP	n	#	OP	n	#	OP	n	#	OP	n	#	OP	n	#	OP	n	#	OP	n	#	Ν	V	т	в	D	Т	z	С
																																	Ν	•	•	•	•	•	Z	•
																																	Ν	٠	•	•	•	•	Z	•
E2	16	2																															•	•	•	•	•	•	•	•
55	4	2				4D	4	3	5D	5	3	59	5	3							41	6	2	51	6	2							N	•	•	•	•	•	Z	•
F6	6	2				EE	6	3	FE	7	3																						N	•	•	•	•	•	z	•
																																	N	•	•	•	•	•	z	•
																																	N	•	•	•	•	•	Z	•
						4C	3	3							6C	5	3	B2	4	2													•	•	•	•	•	•	•	•
						20	6	3										02	7	2										22	5	2	•	•	•	•	•	•	•	•
B5	4	2				AD	4	3	BD	5	3	В9	5	3							A1	6	2	B1	6	2							N	•	•	•	•	•	Z	•
																																	•	•	•	•	•	•	•	•
			B6	4	2	AE	4	3				BE	5	3																			Ν	•	•	•	•	•	Z	•
B4	4	2				AC	4	3	вс	5	3																						Ν	•	•	•	•	•	Z	•

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3.7 Machine instructions

									A	١ddr	ess	ing ı	mod	e						
Symbol	Function	Details		IMF)		IMN	1		А		в	IT,	A		ΖP		BI	Г, Z	P
			OP	n	#	OF	n	#	OP	n	#	OP	n	#	OP	n	#	OP	n	#
LSR		This instruction shifts either A or M one bit to the right such that bit 7 of the result always is set to 0, and the bit 0 is stored in C.							4A	2	1				46	5	2			
MUL	$\begin{array}{l} M(S) \bullet A \leftarrow A * M(zz + X) \\ S \leftarrow S - 1 \end{array}$	Multiplies Accumulator with the memory speci- fied by the Zero Page X address mode and stores the high-order byte of the result on the Stack and the low-order byte in A.																		
NOP	PC ← PC + 1	This instruction adds one to the PC but does no otheroperation.	ΕA	2	1															
ORA (Note 1)	When T = 0 A \leftarrow A V M When T = 1 M(X) \leftarrow M(X) V M	When T = 0, this instruction transfers the con- tents of A and M to the ALU which performs a bit-wise "OR", and stores the result in A. When T = 1, this instruction transfers the con- tents of $M(X)$ and the M to the ALU which performs a bit-wise OR, and stores the result in $M(X)$. The contents of A remain unchanged, but status flags are changed. M(X) represents the contents of memory where is indicated by X.				09	2	2							05	3	2			
PHA	S ← S – 1	This instruction pushes the contents of A to the memory location designated by S, and decrements the contents of S by one.	48	3	1															
PHP	$\begin{array}{l} M(S) \leftarrow PS \\ S \leftarrow S - 1 \end{array}$	This instruction pushes the contents of PS to the memory location designated by S and dec- rements the contents of S by one.	08	3	1															
PLA	$\begin{array}{l} S \leftarrow S + 1 \\ A \leftarrow M(S) \end{array}$	This instruction increments S by one and stores the contents of the memory designated by S in A.	68	4	1															
PLP	$S \leftarrow S + 1$ PS $\leftarrow M(S)$	This instruction increments S by one and stores the contents of the memory location designated by S in PS.	28	4	1															
ROL	7 0 ←□□←℃←	This instruction shifts either A or M one bit left through C. C is stored in bit 0 and bit 7 is stored in C.							2A	2	1				26	5	2			
ROR		This instruction shifts either A or M one bit right through C. C is stored in bit 7 and bit 0 is stored in C.							6A	2	1				66	5	2			
RRF		This instruction rotates 4 bits of the M content to the right.													82	8	2			
RTI	$\begin{array}{l} S \leftarrow S+1 \\ PS \leftarrow M(S) \\ S \leftarrow S+1 \\ PCL \leftarrow M(S) \\ S \leftarrow S+1 \\ PCH \leftarrow M(S) \end{array}$	This instruction increments S by one, and stores the contents of the memory location designated by S in PS. S is again incremented by one and stores the contents of the memory location designated by S in PCL. S is again incremented by one and stores the contents of memory location designated by S in PCH.	40	6	1															
RTS	$\begin{array}{l} S \leftarrow S + 1 \\ PCL \leftarrow M(S) \\ S \leftarrow S + 1 \\ PCH \leftarrow M(S) \\ (PC) \leftarrow (PC) + 1 \end{array}$	This instruction increments S by one and stores the contents of the memory location designated by S in PCL. S is again incremented by one and the contents of the memory location is stored in PCH. PC is incremented by 1.	60	6	1															

												_		Ad	dres	sin	g m	ode			-						-						F	roc	esso	or st	atus	reç	giste	ər
Z	P,)			<u>P,</u>	1		ABS	-		BS,	-		BS,	1		IND	-		P, IN	r —		ID, İ			ND,			REL			SP	1	7	6	5	4	3	2	1	(
OP	n		OP	n	#	OP		-	OP		-	OP	n	#	OP	n	#	OP	n	#	OP	n	#	OP	n	#	OP	n	#	OP	n	#		V	Т	В	D	I	Z	-
56	6	2				4E	6	3	5E	7	3																						0	•	•	•	•	•	Z	0
62	15	2																															•	•	•	•	•	•	•	•
																																	•	•	•	•	•	•	•	•
15	4	2				0D	4	3	1D	5	3	19	5	3							01	6	2	11	6	2							N	•	•	•	•	•	z	•
																																	•	•	•	•	•	•	•	•
																																	•	•	•	•	•	•	•	
																																	N	•	•	•	•	•	z	
																																		(Va	ue	save	ed ir	n sta	ack)	+)
																																								_
36	6	2				2E	6	3	3E	7	3																						N	•	•	•	•	•	z	1
76	6	2				6E	6	3	7E	7	3																						N	•	•	•	•	•	z	,
																																	•	•	•	•	•	•	•	
																																		(Va	ue	save	ed ir	n sta	ack)	
																																	•	•	•	•	•	•	•	

APPENDIX

3.7 Machine instructions

						-			-	Addı	633	ing i	1100	ie	_					
Symbol	Function	Details		IMP	•		IMI	N		А		В	IT,	A		ΖP		Bľ	T, Z	P
			OP	n	#	OF	P n	#	OF	P n	#	OP	n	#	OP	n	#	OP	n	#
SBC (Note 1) (Note 5)	When T = 0 $A \leftarrow A - M - C$ When T = 1 $M(X) \leftarrow M(X) - M - C$	When T = 0, this instruction subtracts the value of M and the complement of C from A, and stores the results in A and C. When T = 1, the instruction subtracts the contents of M and the complement of C from the contents of M(X), and stores the results in $M(X)$ and C. A remain unchanged, but status flag are changed. M(X) represents the contents of memory where is indicated by X.				E	92	2							E5	3	2			
SEB	Ai or Mi ← 1	This instruction sets the designated bit i of A or M.										0₿ 20i	2	1				0F 20i	5	2
SEC	C ← 1	This instruction sets C.	38	2	1															
SED	D ← 1	This instruction set D.	F8	2	1															
SEI	← 1	This instruction set I.	78	2	1															
SET	T ← 1	This instruction set T.	32	2	1															
STA	$M \gets A$	This instruction stores the contents of A in M. The contents of A does not change.													85	4	2			
STP		This instruction resets the oscillation control F/ F and the oscillation stops. Reset or interrupt input is needed to wake up from this mode.	42	2	1															
STX	$M \leftarrow X$	This instruction stores the contents of X in M. The contents of X does not change.													86	4	2			
STY	$M \gets Y$	This instruction stores the contents of Y in M. The contents of Y does not change.													84	4	2			
ТАХ	$X \leftarrow A$	This instruction stores the contents of A in X. The contents of A does not change.	AA	2	1															
TAY	$A \to A$	This instruction stores the contents of A in Y. The contents of A does not change.	A8	2	1															
TST	M = 0?	This instruction tests whether the contents of M are "0" or not and modifies the N and Z.													64	3	2			
TSX	$X \leftarrow S$	This instruction transfers the contents of S in X.	BA	2	1															
ТХА	$A \leftarrow X$	This instruction stores the contents of X in A.	8A	2	1															
TXS	$S \leftarrow X$	This instruction stores the contents of X in S.	9A	2	1															
ΤΥΑ	$A \gets Y$	This instruction stores the contents of Y in A.	98	2	1	T				1										
WIT		The WIT instruction stops the internal clock but not the oscillation of the oscillation circuit is not stopped. CPU starts its function after the Timer X over flows (comes to the terminal count). All regis- ters or internal memory contents except Timer X will not change during this mode. (Of course needs VDD).	C2	2	1															

														Ad	dres	sin	g m	ode															F	Proc	esso	or st	atus	s reg	giste	r
z	2P, 2	x		ZP,	Y		ABS	3	A	BS	, X	A	BS,		1	IND		r –	P, IN	١D	11	ND,	х	11	ND,	Y		REL			SP		7	6	5	4		2	1	0
OP	n	#		n	1	OP	n	#	ОР	n	#	ОР	n	#	OP	n	#	OP	-	-	OP	n	#	OP	n	#	OP	n	#	OP	n	#	N	v	т	в		1	z	С
F5	4	2				ED	4	3	FD	5	3	F9	5	3							E1	6	2	F1	6	2							N	V	•	•	•	•	Z	С
																																	•	•	•	•	•	•	•	•
																																	•	•	•	•	•	•	•	1
																																	•	•	•	•	1	•	•	•
																																	•	•	1	•	•	•	•	•
95	5	2				8D	5	3	9D	6	3	99	6	3							81	7	2	91	7	2							•	•	•	•	•	•	•	•
																																	•	•	•	•	•	•	•	•
			96	5	2	8E	5	3																									•	•	•	•	•	•	•	•
94	5	2	50			8C		3																									•	•	•	•	•	•	•	•
		_																															N	•	•	•	•	•	z	•
																																	N	•	•	•	•	•	z	•
																																	N	•	•	•	•	•	z	•
																																	N	•	•	•	•	•	z	•
																																	N	•	•	•	•	•	z	•
																																	•	•	•	•	•	•	•	•
																																	N	•	•	•	•	•	z	•
																																	•	•	•	•	•	•	•	•

Notes 1 : The number of cycles "n" is increased by 3 when T is 1.
2 : The number of cycles "n" is increased by 2 when T is 1.
3 : The number of cycles "n" is increased by 1 when T is 1.
4 : The number of cycles "n" is increased by 2 when branching has occurred.
5 : N, V, and Z flags are invalid in decimal operation mode.

APPENDIX

Symbol	Contents	Symbol	Contents
IMP	Implied addressing mode	+	Addition
IMM	Immediate addressing mode	-	Subtraction
A	Accumulator or Accumulator addressing mode	*	Multiplication
BIT, A	Accumulator bit addressing mode	/	Division
BIT, A, R	Accumulator bit relative addressing mode	Λ	Logical OR
ZP	Zero page addressing mode	V	Logical AND
BIT, ZP	Zero page bit addressing mode	\forall	Logical exclusive OR
BIT, ZP, R	Zero page bit relative addressing mode	_	Negation
ZP, X	Zero page X addressing mode	\leftarrow	Shows direction of data flow
ZP, Y	Zero page Y addressing mode	Х	Index register X
ABS	Absolute addressing mode	Y	Index register Y
ABS, X	Absolute X addressing mode	S	Stack pointer
ABS, Y	Absolute Y addressing mode	PC	Program counter
IND	Indirect absolute addressing mode	PS	Processor status register
		РСн	8 high-order bits of program counter
ZP, IND	Zero page indirect absolute addressing mode	PCL	8 low-order bits of program counter
		ADH	8 high-order bits of address
IND, X	Indirect X addressing mode	ADL	8 low-order bits of address
IND, Y	Indirect Y addressing mode	FF	FF in Hexadecimal notation
REL	Relative addressing mode	nn	Immediate value
SP	Special page addressing mode	zz	Zero page address
С	Carry flag	м	Memory specified by address designation of any ad-
Z	Zero flag		dressing mode
I	Interrupt disable flag	M(X)	Memory of address indicated by contents of index
D	Decimal mode flag		register X
В	Break flag	M(S)	Memory of address indicated by contents of stack
т	X-modified arithmetic mode flag		pointer
V	Overflow flag	M(ADH, ADL)	Contents of memory at address indicated by ADH and
N	Negative flag		ADL, in ADH is 8 high-order bits and ADL is 8 low-or-
			der bits.
		M(00, ADL)	Contents of address indicated by zero page ADL
		Ai	Bit i (i = 0 to 7) of accumulator
		Mi	Bit i (i = 0 to 7) of memory
		OP	Opcode
		n	Number of cycles
		#	Number of bytes

3.8 List of instruction code

3.8 List of instruction code

\sim																	
	D3 – D0	0000	0001	0010	0011	0100	0101	0110	0111	1000	1001	1010	1011	1100	1101	1 110	1111
D7 – D4	xadecimal notation	0	1	2	3	4	5	6	7	8	9	A	В	с	D	E	F
0000	0	BRK	ORA IND, X	JSR ZP, IND	BBS 0, A	_	ORA ZP	ASL ZP	BBS 0, ZP	PHP	ORA IMM	ASL A	SEB 0, A	_	ORA ABS	ASL ABS	SEB 0, ZP
0001	1	BPL	ORA IND, Y	CLT	BBC 0, A	_	ORA ZP, X	ASL ZP, X	BBC 0, ZP	CLC	ORA ABS, Y	DEC A	CLB 0, A	_	ORA ABS, X	ASL ABS, X	CLB 0, ZP
0010	2	JSR ABS	AND IND, X	JSR SP	BBS 1, A	BIT ZP	AND ZP	ROL ZP	BBS 1, ZP	PLP	AND IMM	ROL A	SEB 1, A	BIT ABS	AND ABS	ROL ABS	SEB 1, ZP
0011	3	BMI	AND IND, Y	SET	BBC 1, A	_	AND ZP, X	ROL ZP, X	BBC 1, ZP	SEC	AND ABS, Y	INC A	CLB 1, A	LDM ZP	AND ABS, X	ROL ABS, X	CLB 1, ZP
0100	4	RTI	EOR IND, X	STP	BBS 2, A	COM ZP	EOR ZP	LSR ZP	BBS 2, ZP	РНА	EOR IMM	LSR A	SEB 2, A	JMP ABS	EOR ABS	LSR ABS	SEB 2, ZP
0101	5	BVC	EOR IND, Y	_	BBC 2, A	_	EOR ZP, X	LSR ZP, X	BBC 2, ZP	CLI	EOR ABS, Y	_	CLB 2, A	_	EOR ABS, X	LSR ABS, X	CLB 2, ZP
0110	6	RTS	ADC IND, X	MUL ZP, X	BBS 3, A	TST ZP	ADC ZP	ROR ZP	BBS 3, ZP	PLA	ADC IMM	ROR A	SEB 3, A	JMP IND	ADC ABS	ROR ABS	SEB 3, ZP
0111	7	BVS	ADC IND, Y	_	BBC 3, A	_	ADC ZP, X	ROR ZP, X	BBC 3, ZP	SEI	ADC ABS, Y	_	CLB 3, A	_	ADC ABS, X	ROR ABS, X	CLB 3, ZP
1000	8	BRA	STA IND, X	RRF ZP	BBS 4, A	STY ZP	STA ZP	STX ZP	BBS 4, ZP	DEY	_	ТХА	SEB 4, A	STY ABS	STA ABS	STX ABS	SEB 4, ZP
1001	9	всс	STA IND, Y	_	BBC 4, A	STY ZP, X	STA ZP, X	STX ZP, Y	BBC 4, ZP	ΤΥΑ	STA ABS, Y	тхѕ	CLB 4, A	-	STA ABS, X	_	CLB 4, ZP
1010	А	LDY IMM	LDA IND, X	LDX IMM	BBS 5, A	LDY ZP	LDA ZP	LDX ZP	BBS 5, ZP	ΤΑΥ	LDA IMM	ТАХ	SEB 5, A	LDY ABS	LDA ABS	LDX ABS	SEB 5, ZP
1011	В	BCS	LDA IND, Y	JMP ZP, IND	BBC 5, A	LDY ZP, X	LDA ZP, X	LDX ZP, Y	BBC 5, ZP	CLV	LDA ABS, Y	тѕх	CLB 5, A	LDY ABS, X	LDA ABS, X	LDX ABS, Y	CLB 5, ZP
1100	с	CPY IMM	CMP IND, X	WIT	BBS 6, A	CPY ZP	CMP ZP	DEC ZP	BBS 6, ZP	INY	CMP IMM	DEX	SEB 6, A	CPY ABS	CMP ABS	DEC ABS	SEB 6, ZP
1101	D	BNE	CMP IND, Y	-	BBC 6, A	_	CMP ZP, X	DEC ZP, X	BBC 6, ZP	CLD	CMP ABS, Y	_	CLB 6, A	_	CMP ABS, X	DEC ABS, X	CLB 6, ZP
1110	E	CPX IMM	SBC IND, X	DIV ZP, X	BBS 7, A	CPX ZP	SBC ZP	INC ZP	BBS 7, ZP	INX	SBC IMM	NOP	SEB 7, A	CPX ABS	SBC ABS	INC ABS	SEB 7, ZP
1111	F	BEQ	SBC IND, Y	_	BBC 7, A	_	SBC ZP, X	INC ZP, X	BBC 7, ZP	SED	SBC ABS, Y	_	CLB 7, A	_	SBC ABS, X	INC ABS, X	CLB 7, ZP

: 3-byte instruction

: 2-byte instruction

: 1-byte instruction

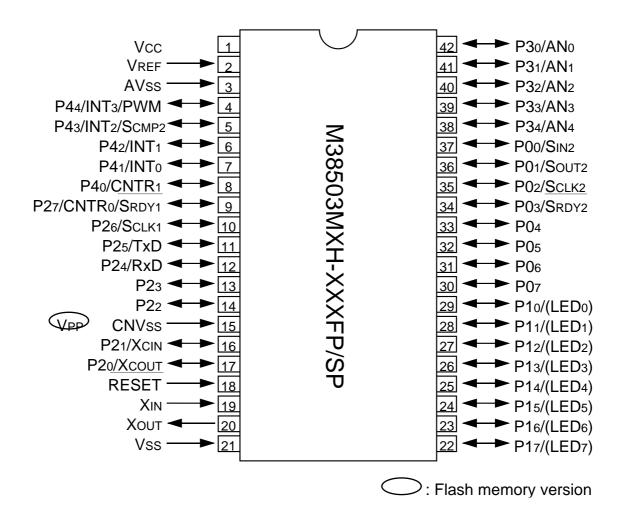
3.9 SFR memory map

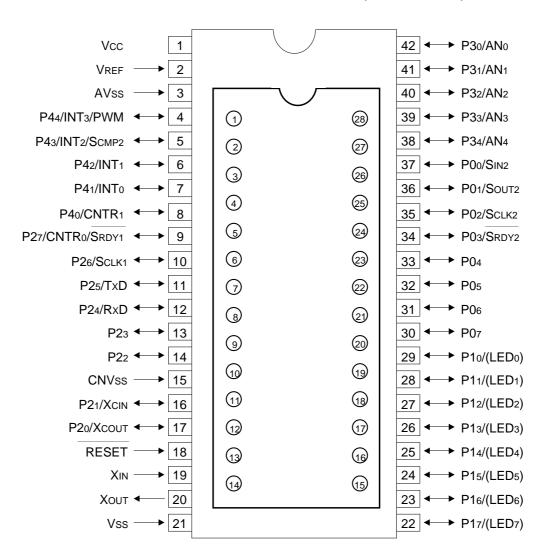
3.9 SFR memory map

000016	Port P0 (P0)	002016	Prescaler 12 (PRE12)
000116	Port P0 direction register (P0D)	002116	Timer 1 (T1)
000216	Port P1 (P1)	002216	Timer 2 (T2)
000316	Port P1 direction register (P1D)	002316	Timer XY mode register (TM)
000416	Port P2 (P2)	002416	Prescaler X (PREX)
000516	Port P2 direction register (P2D)	002516	Timer X (TX)
000616	Port P3 (P3)	002616	Prescaler Y (PREY)
000716	Port P3 direction register (P3D)	002716	Timer Y (TY)
000816	Port P4 (P4)	002816	Timer count source selection register (TCSS)
000916	Port P4 direction register (P4D)	002916	
000A16		002A16	
000B16		002B16	Reserved *
000C16		002C16	Reserved *
000D16		002D16	Reserved *
000E16		002E16	Reserved *
000F16		002F16	Reserved *
001016		003016	Reserved *
001116		003116	Reserved *
001216	Reserved *	003216	
001316	Reserved *	003316	
001416	Reserved *	003416	A-D control register (ADCON)
001516	Serial I/O2 control register 1 (SIO2CON1)	003516	A-D conversion low-order register (ADL)
001616	Serial I/O2 control register 2 (SIO2CON2)	003616	A-D conversion high-order register (ADH)
001716	Serial I/O2 register (SIO2)	003716	Reserved *
001816	Transmit/Receive buffer register (TB/RB)	003816	MISRG
001916	Serial I/O1 status register (SIOSTS)	003916	Watchdog timer control register (WDTCON)
001A16	Serial I/O1 control register (SIOCON)	003A16	Interrupt edge selection register (INTEDGE)
001B16	UART control register (UARTCON)	003B16	CPU mode register (CPUM)
001C16	Baud rate generator (BRG)	003C16	Interrupt request register 1 (IREQ1)
001D16	PWM control register (PWMCON)	003D16	Interrupt request register 2 (IREQ2)
001E16	PWM prescaler (PREPWM)	003E16	Interrupt control register 1 (ICON1)
001F16	PWM register (PWM)	003F16	Interrupt control register 2 (ICON2)
		0FFE16	Flash memory control register (FMCR)

* Reserved : Do not write any data to this addresses, because these areas are reserved.

3.10 Pin configurations





M38517RSS PIN CONFIGURATION (TOP VIEW)

Outline : 42S1M

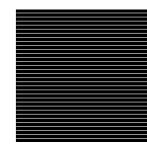
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