SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

DESCRIPTION

The 4570 Group is a 4-bit single-chip microcomputer designed with CMOS technology. Its CPU is that of the 4500 series using a simple, high-speed instruction set. The computer is equipped with a carrier wave output circuit for remote control, an 8-bit timer with a reload register, a 10-bit timer with a reload register, and an 8-bit timer with two reload registers.

The various microcomputers in the 4570 Group include variations of the built-in memory size. The mask ROM version and One Time PROM version of 4570 Group are produced as shown in the table below.

FEATURES

Minimum instruction execution time

When f(XiN)/4 is selected for system clock 2.86 μ s (f(XiN)=4.2 MHz, VDD=2.0 V to 5.5 V)

Supply voltage

 • System clock switch function

f(XIN)/4 or not divided
 Timers

Timer 1... 10-bit timer with a reload register and carrier wave output auto-control function

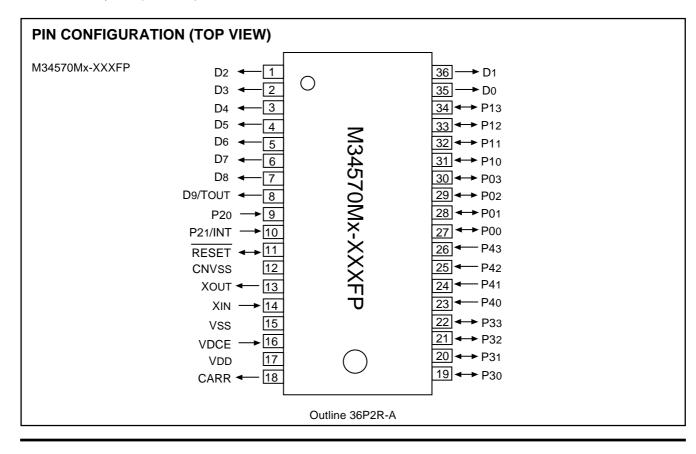
- generation function
- Interrupt 4 sources
- Power-on reset circuit
- Watchdog timer16 bits
- Key-on wakeup function (Ports P0, P1, and P4, ON/OFF of port P4 can be switched)
- Pull-up transistor (Ports P0, P1, and P4, ON/OFF of port P4 can be switched)
- Voltage drop detection circuit
- Clock generating circuit (ceramic resonance)

APPLICATION

Remote control transmitter

Product	ROM (PROM) size (X 10 bits)	RAM size (X 4 bits)	Package	ROM type
M34570M4-XXXFP	4096 words	128 words	36P2R-A	Mask ROM
M34570M8-XXXFP	8192 words	128 words	36P2R-A	Mask ROM
M34570MD-XXXFP	16384 words	128 words	36P2R-A	Mask ROM
M34570E8FP	8192 words	128 words	36P2R-A	One Time PROM
M34570EDFP *	16384 words	128 words	36P2R-A	One Time PROM

*: Under development (Jan. 1999)





Note: PROM 16384 words X 10 bits for the built-in PROM version. 4096 to 16384 words X 10 bits 10 Port D Reset (Voltage drop detection circuit) 128 words X 4 bits Memory Clock generating circuit ROM (Note) RAM XIN –XOUT Port P4 Register B (4 bits) Register E (8 bits) Stack registers SKs (8 levels) Interrupt stack register SDP(1 level) Port P3 4500 Series CPU core ALU(4 bits) Register A (4 bits) Register D (3 bits) Port P2 Internal peripheral functions 2 Port P1 Timers/Carrier wave generation Watchdog timer (16 bits) (Carrier wave generation) Timer 1 (10 bits) Timer 2 (8 bits) Timer 3 (8 bits) Port P0 I/O port

BLOCK DIAGRAM

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

4570 Group

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

Pa	arameter		Function		
Number of basic instructions		ions	99		
Minimum instruction execution time		cution time	1.5 μs (f(XιN) = 2.0 MHz:system clock = f(XιN): VDD = 5.0 V)		
			2.86 μs (f(XiN) = 4.2 MHz:system clock = f(XiN)/4: VDD = 5.0 V)		
Memory sizes	ROM	M34570M4	096 words X 10 bits		
		M34570M8	8192 words X 10 bits		
		M34570MD	16384 words X 10 bits		
		M34570E8	8192 words X 10 bits		
		M34570ED	16384 words X 10 bits		
	RAM		128 words X 4 bits		
Input/Output	Do-D9	Output	Ten independent output ports; port D ₉ is also used as the Tout output pin.		
ports	P00-P03	I/O	4-bit I/O port; every pin of the ports has a key-on wakeup function and a pull-up function.		
	P10-P13	I/O	4-bit I/O port; every pin of the ports has a key-on wakeup function and a pull-up function.		
	P20, P21	Input	2-bit input port, port P21 is also used as INT input pin.		
	P30-P33	I/O	4-bit I/O port		
	P40-P43	Input	4-bit input port; both pull-up function and key-on wakeup function can be switched by software.		
	CARR	Output	1-bit output port (CMOS output)		
	Тоит	Output	1-bit output pin; Tou⊤ output pin is also used as port D₀.		
	INT	Input	1-bit input pin with a key-on wakeup function. INT input pin is also used as port P21.		
Timers	Timer 1		10-bit timer with a reload register and carrier wave output auto-control function		
	Timer 2		8-bit timer with a reload register		
	Timer 3		8-bit timer with two reload registers and carrier wave generation function		
Interrupt	Sources		4 (one for external and three for timer)		
	Nesting		1 level		
Subroutine nes	sting		8 levels (however, only 7 levels can be used when an interrupt is used or the TABP p instruction		
			is executed)		
Device structur	re		CMOS silicon gate		
Package			36-pin plastic molded SSOP		
Operating tem	perature ra	ange	–20 °C to 70 °C		
Supply voltage	•		2.0 V to 5.5 V for mask ROM version (2.5 V to 5.5 V for One Time PROM version)		
Power	at active		1.3 mA ($f(X_{IN}) = 4.2 \text{ MHz}$: system clock = $f(X_{IN})/4$, VDD=5.0 V)		
dissipation			0.5 mA (f(XIN) = 1.0 MHz: system clock = f(XIN), VDD=3.0 V)		
(typical value)	at RAM b	ack-up	0.1 μ A (Ta=25 °C, VDD=5V, typical value)		

PERFORMANCE OVERVIEW

DEFINITION OF CLOCK AND CYCLE

System clock

The system clock is the basic clock for controlling this product. The system clock can be selected by bit 3 of the clock control register MR as shown in the table below.

Table Selection of system clock

MR3	System clock
0	f(XIN)
1	f(XIN)/4

Note: f(XIN)/4 is selected immediately after system is released from reset.

Instruction clock

The instruction clock is the standard clock for controlling CPU. The instruction clock is a signal derived from dividing the system clock by 3. The one cycle of the instruction clock is equivalent to the one machine cycle.

Machine cycle

The machine cycle is the standard cycle required to execute the instruction.

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

PIN DESCRIPTION

Pin	Name	Input/Output	Function		
Vdd	Power supply	_	Connected to a plus power supply.		
Vss	Ground	_	Connected to a 0 V power supply.		
CNVss	CNVss	Input	Connect CNVss to Vss and apply "L" (0V) to CNVss certainly.		
RESET	Reset input	I/O	An N-channel open-drain I/O pin for a system reset. A pull-up transistor		
			capacitor are built-in this pin. When the watchdog timer causes the system to be		
			reset or the low-supply voltage is detected, the RESET pin outputs "L" level.		
Xin	Clock input	Input	I/O pins of the clock generating circuit. Connect a ceramic resonator between XIN		
Хоит	Clock output	Output	pin and Xou⊤ pin. A feedback resistor is built-in between them.		
D0-D9	Output port D	Output	Each pin of port D has an independent 1-bit wide output function. Port D ₉ is also		
			used as Tout output pin. The output structure is N-channel open-drain.		
P00-P03	I/O port P0	I/O	4-bit I/O port. It can be used as an input port when the output latch is set to "1."		
			The output structure is N-channel open-drain. Every pin of the ports has a key-on		
			wakeup function and a pull-up function.		
P10–P13	I/O port P1	I/O	4-bit I/O port. It can be used as an input port when the output latch is set to "1."		
			The output structure is N-channel open-drain. Every pin of the ports has a key-on		
			wakeup function and a pull-up function.		
P20, P21	Input port P2	I/O	2-bit input port. Port P21 is also used as the INT input pin.		
P30-P33	I/O port P3	I/O	4-bit I/O port. It can be used as an input port when the output latch is set to "		
			The output structure is N-channel open-drain.		
P40–P43	Input port P4	Input	4-bit input port. Every pin of the ports has a key-on wakeup function and a pull-up		
			function. Both functions can be switched by software.		
CARR	Carrier wave output	Output	Carrier wave output pin for remote control transmit. The output structure is the		
	for remote control		CMOS circuit.		
INT	Interrupt input	Input	INT input pin accepts an external interrupt and has a key-on wakeup function. INT		
			input pin is also used as port P21.		
Τουτ	Timer output	Output	Tout output pin has the function to output the timer 2 underflow signal divided by		
			2. Tou⊤ output pin is also used as port D ₉ .		
VDCE	Voltage drop	Input	VDCE pin is used to control the operation/stop of the voltage drop detection circuit.		
	detection circuit		The circuit is operating when "H" level is input to the VDCE pin. It is stopped when		
	enable		"L" level is input to this pin.		

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

MULTIFUNCTION

Pin	Multifunction	Pin	Multifunction
D9	Тоит	Τουτ	D9
P21	INT	INT	P21

Notes 1: Pins except above have just single function.

2: The port D9 is the output port and port P21 is the input port.

CONNECTIONS OF UNUSED PINS

Pin	Connection	Pin	Connection
D0-D8	Connect to Vss, or set the output latch to	P30–P33	Connect to Vss, or set the output latch to
D9/Tout	"0" and open.		"0" and open.
P00–P03	Set the output latch to "1" and open.	P40-P43	Connect to Vss (Note 2) or open (Note 3).
P10–P13		CARR	Open.
P20, P21/INT	Connect to Vss (Note 1).		

Notes 1: When the P21/INT pin is connected to Vss pin, set the return level to "H" level by software (interrupt control register I12="1"). When the P21/INT pin is connected to Vss pin while the return level is set to "L" level, system returns from RAM back-up state immediately after system enters the RAM back-up state.

2: In order to connect ports P40–P43 to Vss, turn off their pull-up transistors (pull-up control register PU0i="0") by software and also invalidate the key-on wakeup functions (key-on wakeup control register K0i="0"). When these pins are connected to Vss while the key-on wakeup functions are left valid, the system fails to return from RAM back-up state. In order to make these pins open, turn on their pull-up transistors (register PU0i="1") by software (i = 0, 1, 2, 3).
Descent to explore the key-on wakeup functions and the pull-up transistors (register PU0i="1") by software (i = 0, 1, 2, 3).

Be sure to select the key-on wakeup function and the pull-up function with every one port.

3: In order to make ports $P4_0-P4_3$ open, turn on their pull-up transistors (register $PU_0i = "1"$) by software (i = 0, 1, 2, 3).

(Note in order to set the output latch to "0" or "1" or make pins open)

- After system is released from reset, a port is in a high-impedance state until the output latch of the port is set to "0" by software. Accordingly, the voltage level of pins is undefined and the excess of the supply current may occur.
- To set the output latch periodically is recommended because the value of output latch may change by noise or a program run away (caused by noise).

(Note in order to connect unused pins to Vss)

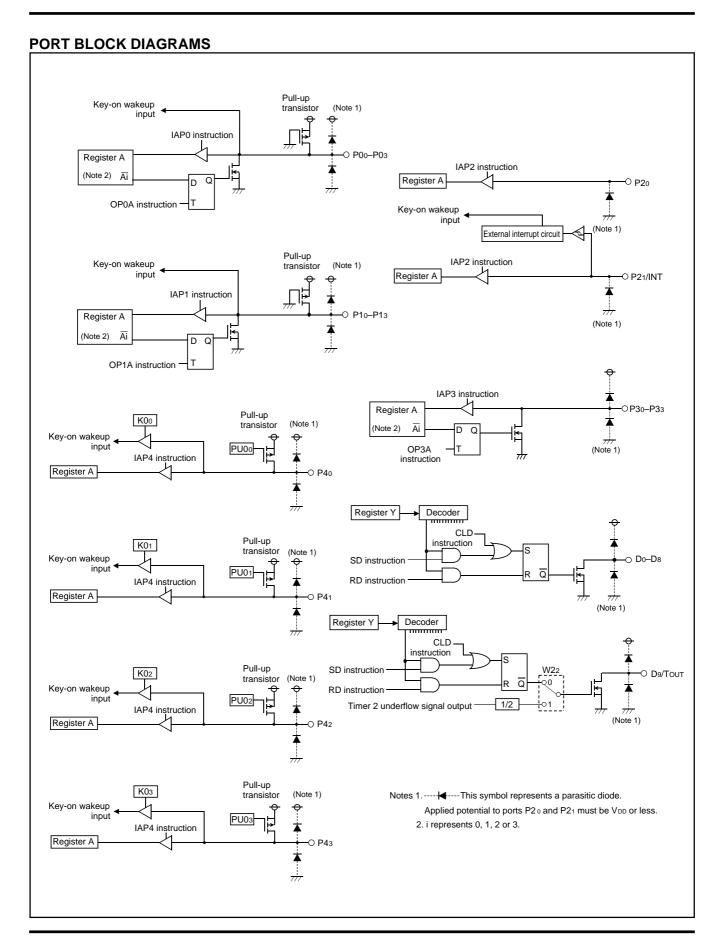
• To avoid noise, connect the unused pins to Vss at the shortest distance using a thick wire.

PORT FUNCTION

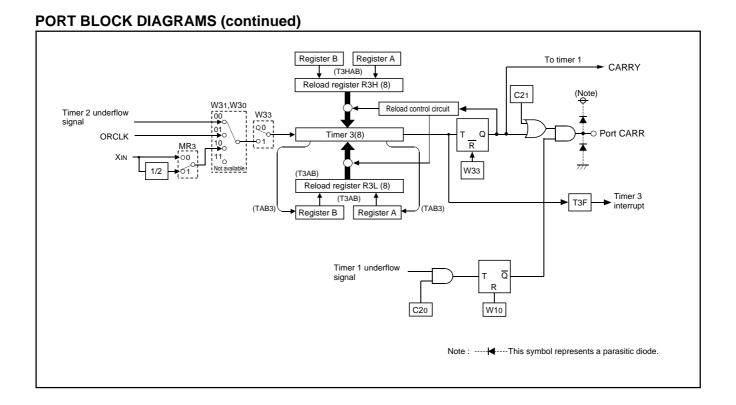
Dort	Port Pin Input/ Output Output structure		Outrast structure	Control	Control	Control	
Pon			Output structure	bits	instructions	registers	Remark
Port D	D0-D8, D9/TOUT	Output	N-channel open-drain	1	SD	W22	W22 controls the switch of D9/
		(10)			RD		Tout pin
					CLD		
Port P0	P00–P03	I/O	N-channel open-drain	4	OP0A		Pull-up functions
		(4)			IAP0		Key-on wakeup functions
Port P1	P10–P13	I/O	N-channel open-drain	4	OP1A		Pull-up functions
		(4)			IAP1		Key-on wakeup functions
Port P2	P20	Input		2	IAP2		
		(2)			SNZI0		
	P21/INT				(Note)		Key-on wakeup function
Port P3	P30–P33	I/O	N-channel open-drain	4	OP3A		
					IAP3		
Port P4	P40–P43	Input		4	IAP4	PU0	Pull-up functions
		(4)				K0	(programmable)
							Key-on wakeup functions
							(programmable)

Note: Level of the P21/INT pin can be examined with the SNZI0 instruction.

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER



SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER



7

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

FUNCTION BLOCK OPERATIONS CPU

(1) Arithmetic logic unit (ALU)

The arithmetic logic unit ALU performs 4-bit arithmetic such as 4-bit data addition, comparison, AND operation, OR operation, and bit manipulation.

(2) Register A and carry flag (CY)

Register A is a 4-bit register used for arithmetic, transfer, exchange, and I/O operation.

Carry flag CY is a 1-bit flag that is set to "1" when there is a carry with the AMC instruction (Figure 1).

It is unchanged with both A n instruction and AM instruction. The value of A_0 is stored in carry flag CY with the RAR instruction (Figure 2).

Carry flag CY can be set to "1" with the SC instruction and cleared to "0" with the RC instruction.

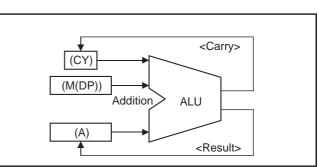
(3) Registers B and E

Register B is a 4-bit register used for temporary storage of 4bit data, and for 8-bit data transfer together with register A. Register E is an 8-bit register. It can be used for 8-bit data transfer with register B used as the high-order 4 bits and register A as the low-order 4 bits (Figure 3).

(4) Register D

Register D is a 3-bit register.

It is used to store a 7-bit ROM address together with register A and is used as a pointer within the specified page when the TABP p, BLA p, or BMLA p instruction is executed (Figure 4).





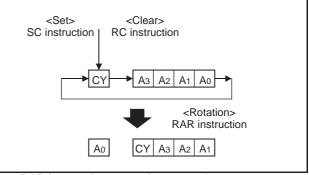


Fig. 2 RAR instruction execution example

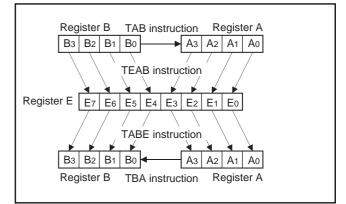


Fig. 3 Registers A, B and register E

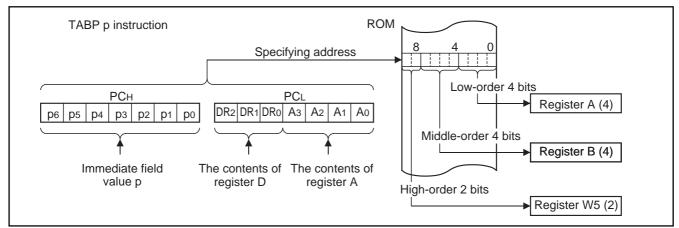


Fig. 4 TABP p instruction execution example

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

(5) Stack registers (SKs) and stack pointer (SP)

Stack registers (SKs) are used to temporarily store the contents of program counter (PC) just before branching until returning to the original routine when;

- branching to an interrupt service routine (referred to as an interrupt service routine),
- · performing a subroutine call, or
- executing the table reference instruction (TABP p).

Stack registers (SKs) are eight identical registers, so that subroutines can be nested up to 8 levels. However, one of stack registers is used when using an interrupt service routine or when executing a table reference instruction. Accordingly, be careful not to stack over when performing these operations together. The contents of registers SKs are destroyed when 8 levels are exceeded.

The register SK nesting level is pointed automatically by 3bit stack pointer (SP). The contents of the stack pointer (SP) can be transferred to register A with the TASP instruction. Figure 5 shows the stack registers (SKs) structure.

Figure 6 shows the example of operation at subroutine call.

(6) Interrupt stack register (SDP)

Interrupt stack register (SDP) is a 1-stage register. When an interrupt occurs, this register (SDP) is used to temporarily store the contents of data pointer, carry flag, skip flag, register A, and register B just before an interrupt until returning to the original routine.

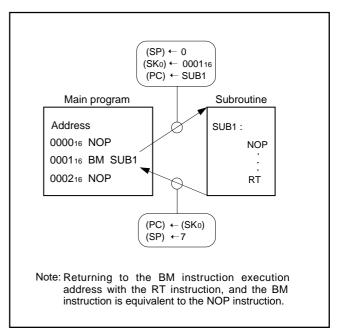
Unlike the stack registers (SKs), this register (SDP) is not used when executing the subroutine call instruction and the table reference instruction.

(7) Skip flag

Skip flag controls skip decision for the conditional skip instructions and continuous described skip instructions. When an interrupt occurs, the contents of skip flag is stored automatically in the interrupt stack register (SDP) and the skip condition is retained.

	Program	cou	nter (PC)		
Executing th call or table instruction	ne subroutine reference	,	Executing the table reference		
		SK0		(SP) = 0	
		SK1		(SP) = 1	
		SK2		(SP) = 2	
		SK3		(SP) = 3	
		(SP) = 4			
		SK5			
		SK6		(SP) = 6	
		SK7		(SP) = 7	
Stack pointer (SP) points "7" at reset or returning from RAM back-up mode. It points "0" by executing the first BM instruction, and the contents of program counter is stored in SKo. When the BM instruction is executed after eight stack registers are used ((SP) = 7), (SP) = 0 and the contents of SKo is destroyed.					







SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

(8) Program counter (PC)

Program counter (PC) is used to specify a ROM address (page and address). It determines a sequence in which instructions stored in ROM are read. It is a binary counter that increments the number of instruction bytes each time an instruction is executed. However, the value changes to a specified address when branch instructions, subroutine call instructions, return instructions, or the table reference instruction (TABP p) is executed.

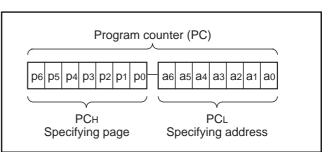
Program counter consists of PCH (most significant bit to bit 7) which specifies to a ROM page and PCL (bits 6 to 0) which specifies an address within a page. After it reaches the last address (address 127) of a page, it specifies address 0 of the next page (Figure 7).

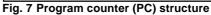
Make sure that the PC $_{\rm H}$ does not specify after the last page of the built-in ROM.

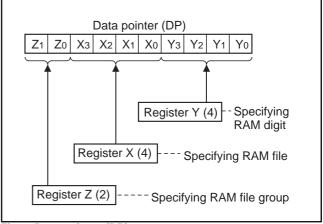
(9) Data pointer (DP)

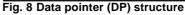
Data pointer (DP) is used to specify a RAM address and consists of registers Z, X, and Y. Register Z specifies a RAM file group, register X specifies a file, and register Y specifies a RAM digit (Figure 8).

Register Y is also used to specify the port D bit position. When using port D, set the port D bit position to register Y certainly and execute the SD or RD instruction (Figure 9).









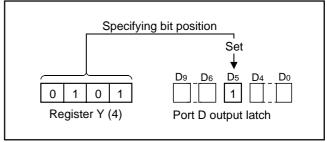


Fig. 9 SD instruction execution example

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

PROGRAM MEMORY (ROM)

1 word of ROM is composed of 10 bits. ROM is separated every 128 words by the unit of page (addresses 0 to 127). Table 1 shows the ROM size and pages. Figure 10 shows the ROM map of M34570M8.

Table	1	ROM	size	and	pages
-------	---	-----	------	-----	-------

Product	ROM size	Pagas		
FIOUUCI	(X 10 bits)	Pages		
M34570M4	4096 words	32 (0 to 31)		
M34570M8	8192 words	64 (0 to 63)		
M34570E8	8192 words	64 (0 to 63)		
M34570MD	16384 words	128 (0 to 127)		
M34570ED	16384 words	128 (0 to 127)		

Note: When the TABP instruction is executed after executing the SBK instruction, data in pages 64 to 127 can be referred. When the TABP instruction is executed after executing the RBK instruction, data in pages 0 to 63 can be referred.

A top part of page 1 (addresses 008016 to 00FF16) is reserved for interrupt addresses (Figure 11). When an interrupt occurs, the address (interrupt address) corresponding to each interrupt is set in the program counter, and the instruction at the interrupt address is executed. When using an interrupt service routine, write the instruction generating the branch to that routine at an interrupt address.

Page 2 (addresses 010016 to 017F16) is the special page for subroutine calls. Subroutines written in this page can be called from any page with the 1-word instruction (BM). Subroutines extending from page 2 to another page can also be called with the BM instruction when it starts on page 2.

ROM pattern (bits 9 to 0) of all addresses can be used as data areas with the TABP $\ensuremath{\mathsf{p}}$ instruction.

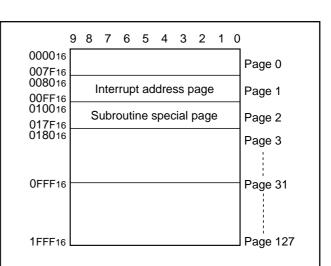


Fig. 10 ROM map of M34570Mx

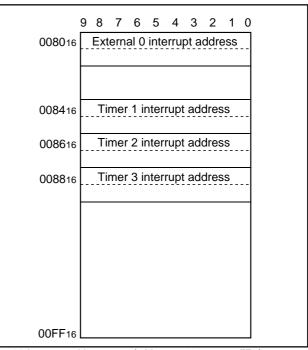


Fig. 11 Interrupt address page (addresses 008016 to 00FF16) structure

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

DATA MEMORY (RAM)

1 word of RAM is composed of 4 bits, but 1-bit manipulation (with the SB j, RB j, and SZB j instructions) is enabled for the entire memory area. A RAM address is specified by a data pointer. The data pointer consists of registers Z, X, and Y. Set a value to the data pointer certainly when executing an instruction to access RAM.

Table 2 shows the RAM size. Figure 12 shows the RAM map.

T	abl	e	2	RA	м	size
	~~~	•	_			0.20

Product	RAM size		
M34570Mx			
M34570Ex	128 words X 4 bits (512 bits)		

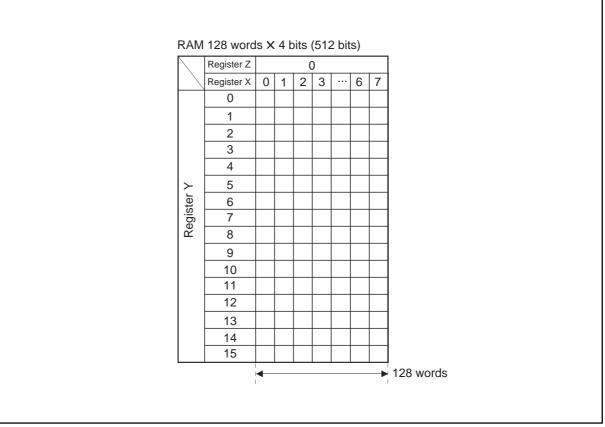


Fig. 12 RAM map

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

# INTERRUPT FUNCTION

The interrupt type is a vectored interrupt branching to an individual address (interrupt address) according to each interrupt source. An interrupt occurs when the following 3 conditions are satisfied.

- Interrupt enable flag (INTE) = "1" (Interrupt enabled)
- Interrupt enable bit = "1" (Interrupt request occurrence enabled)
- An interrupt activated condition is satisfied

(request flag = "1")

Table 3 shows interrupt sources. (Refer to each interrupt request flag for details of activated conditions.)

### (1) Interrupt enable flag (INTE)

The interrupt enable flag (INTE) controls whether the every interrupt enable/disable. Interrupts are enabled when INTE flag is set to "1" with the EI instruction and disabled when INTE flag is cleared to "0" with the DI instruction. When any interrupt occurs, the INTE flag is automatically cleared to "0," so that other interrupts are disabled until the EI instruction is executed.

### (2) Interrupt enable bits (V10-V13, V20-V23)

Use an interrupt enable bit of interrupt control registers V1 and V2 to select the corresponding interrupt request or skip instruction.

Table 4 shows the interrupt request flag, interrupt enable bit and skip instruction.

Table 5 shows the interrupt enable bit function.

### (3) Interrupt request flag

When the activated condition for each interrupt is satisfied, the corresponding interrupt request flag is set to "1." Each interrupt request flag is cleared to "0" when either;

• an interrupt occurs, or

• the next instruction is skipped with a skip instruction.

Each interrupt request flag is set when the activated condition is satisfied even if the interrupt is disabled by the INTE flag or its interrupt enable bit. Once set, the interrupt request flag retains set until a clear condition is satisfied.

Accordingly, an interrupt occurs when the interrupt disable state is released while the interrupt request flag is set.

If more than one interrupt request flag is set when the interrupt disable state is released, the interrupt priority level is as follows shown in Table 3.

### Table 3 Interrupt sources

Priority	Interrupt nome	Activated condition	Interrupt				
level	Interrupt name	Activated condition	address				
1	External 0 interrupt	Level change of	Address 0				
		INT pin	in page 1				
2	Timer 1 interrupt	Timer 1 underflow	Address 4				
			in page 1				
3	Timer 2 interrupt	Timer 2 underflow	Address 6				
			in page 1				
4	Timer 3 interrupt	Timer 3 underflow	Address 8				
			in page 1				

# Table 4 Interrupt request flag, interrupt enable bit and skip instruction

Interrupt name	Request flag	Enable bit	Skip instruction						
External 0 interrupt	EXF0	V10	SNZ0						
Timer 1 interrupt	T1F	V12	SNZT1						
Timer 2 interrupt	T2F	V13	SNZT2						
Timer 3 interrupt	T3F	V20	SNZT3						

### Table 5 Interrupt enable bit function

Interrupt enable bit	Occurrence of interrupt request	Skip instruction	
1	Enabled		
0	Disabled	Valid	

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

#### (4) Internal state during an interrupt

The internal state of the microcomputer during an interrupt is as follows (Figure 14).

- Program counter (PC) An interrupt address is set in program counter. The address to be executed when returning to the main routine is automatically stored in the stack register (SK).
- Interrupt enable flag (INTE) INTE flag is cleared to "0" so that interrupts are disabled.
  Interrupt request flag

Only the request flag for the current interrupt source is cleared to "0."

• Data pointer, carry flag, skip flag, registers A and B The contents of these registers and flags are stored automatically in the interrupt stack register (SDP).

#### (5) Interrupt processing

When an interrupt occurs, a program at an interrupt address is executed after a branch to a sequence for storing data into stack register is performed. Write the branch instruction to an interrupt service routine at an interrupt address.

Use the RTI instruction to return to main routine.

Interrupt enabled by executing the EI instruction is performed after executing 1 instruction (just after the next instruction is executed). Accordingly, when the EI instruction is executed just before the RTI instruction, interrupts are enabled after returning to the main routine. (Refer to Figure 13)

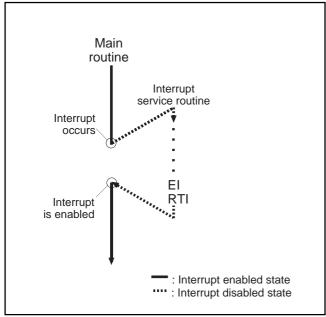


Fig. 13 Program example of interrupt processing

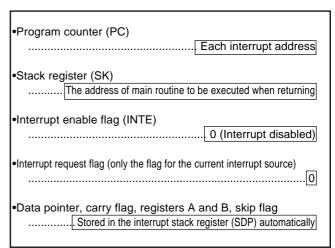
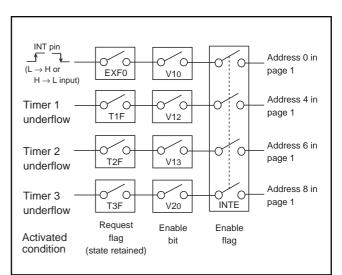
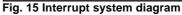


Fig. 14 Internal state when interrupt occurs





SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

# (6) Interrupt control register

Interrupt control register V1

Interrupt enable bits of external 0, timer 1 and timer 2 are assigned to register V1. Set the contents of this register through register A with the TV1A instruction. The TAV1 instruction can be used to transfer the contents of register V1 to register A.

### Table 6 Interrupt control register

Interrupt control register V2

Interrupt enable bit of timer 3 is assigned to register V2. Set the contents of this register through register A with the TV2A instruction. The TAV2 instruction can be used to transfer the contents of register V2 to register A.

	Interrupt control register V1		reset : 00002	RAM back-up : 00002	R/W	
V13	V13 Timer 2 interrupt enable bit		Interrupt disabled (	Interrupt disabled (SNZT2 instruction is valid)		
V 13		1	Interrupt enabled (	Interrupt enabled (SNZT2 instruction is invalid)		
V12	V12 Timer 1 interrupt enable bit	0	Interrupt disabled (SNZT1 instruction is valid)			
VIZ		1	Interrupt enabled (	SNZT1 instruction is invalid)		
V11	Not used	0	This bit has no function, but read/write is enabled.			
VII	Not used	1		clion, but read/write is enabled.		
V10			Interrupt disabled (	SNZ0 instruction is valid)		
VIO	External 0 interrupt enable bit	1	Interrupt enabled (	SNZ0 instruction is invalid)		

Interrupt control register V2		at reset : 00002		at RAM back-up : 00002	R/W		
V23	V23 Not used		This hit has no function, but road/write is enabled				
•=•		1	This bit has no full	This bit has no function, but read/write is enabled.			
V22	/22 Not used						
V 22	V22 Not used	1	This bit has no function, but read/write is enabled.				
V21	Not used	0	This hit has no function, but road/write is enabled				
V21	Not used	1	This bit has no function, but read/write is enabled.				
V20			Interrupt disabled (SNZT3 instruction is valid)				
V20	Timer 3 interrupt enable bit	1	Interrupt enabled (	SNZT3 instruction is invalid)			

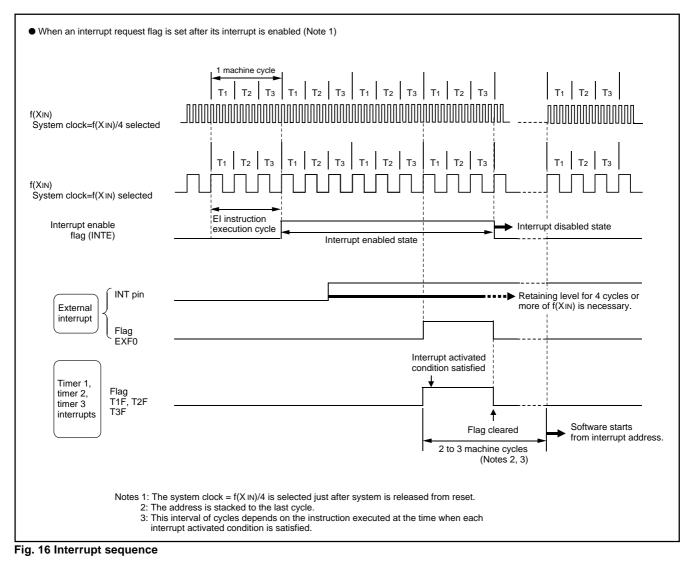
Note: "R" represents read enabled, and "W" represents write enabled.

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

#### (7) Interrupt sequence

Interrupts occur only when the respective INTE flag, interrupt enable bits (V10–V13 and V20–V23), and interrupt request flags (EXF0, T1F, T2F, T3F) are "1." The interrupt actually occurs 2 to 3 machine cycles after the cycle in which all three

conditions are satisfied. The interrupt occurs after 3 machine cycles only when the three interrupt conditions are satisfied on execution of instructions other than one-cycle instructions (Refer to Figure 16).



SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

# **EXTERNAL INTERRUPTS**

An external interrupt request occurs when a valid waveform (= waveform causing the external 0 interrupt) is input to an interrupt input pin (edge detection).

The external 0 interrupt can be controlled with the interrupt control register 11.

#### Table 7 External interrupt activated condition

Name	Input pin	Valid waveform	Valid waveform selection bit (I12)
External 0 interrupt	P21/INT	Falling waveform ("H"→"L")	0
		Rising waveform ("L"→"H")	1

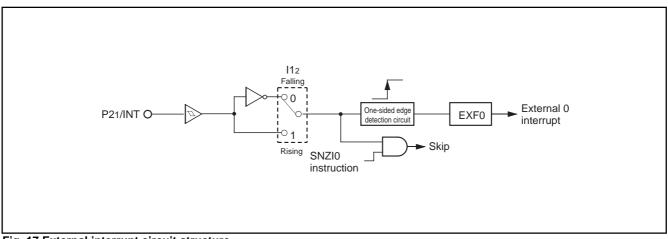


Fig. 17 External interrupt circuit structure

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

#### (1) External 0 interrupt request flag (EXF0)

External 0 interrupt request flag (EXF0) is set to "1" when a valid waveform is input to P21/INT pin.

The valid waveforms causing the interrupt must be retained at their level for 4 cycles or more of the system clock (Refer to Figure 16).

The state of EXF0 flag can be examined with the skip instruction (SNZ0). Use the interrupt control register V1 to select the interrupt or the skip instruction. The EXF0 flag is cleared to "0" when an interrupt occurs or when the next instruction is skipped with the skip instruction.

The P21/INT pin need not be selected the external interrupt input INT function or the normal input port P21 function. However, the EXF0 flag is set to "1" when a valid waveform is input to P21/INT pin even if it is used as an input port P21.

External 0 interrupt activated condition

External 0 interrupt activated condition is satisfied when a valid waveform is input to P21/INT pin.

The valid waveform can be selected from rising waveform or falling waveform. An example of how to use the external 0 interrupt is as follows.

Select the valid waveform with the bit 2 of register I1.

- $\ensuremath{\textcircled{@}}$  Clear the EXF0 flag to "0" with the SNZ0 instruction.
- ③ Set the NOP instruction for the case when a skip is performed with the SNZ0 instruction.
- ④ Set both the external 0 interrupt enable bit (V10) and the INTE flag to "1."

The external 0 interrupt is now enabled. Now when a valid waveform is input to the P21/INT pin, the EXF0 flag is set to "1" and the external 0 interrupt occurs.

### (2) External interrupt control register

Interrupt control register I1

Register 11 controls the valid waveform for the external 0 interrupt, the return level (valid level of wakeup signal) from the RAM back-up and P21/INT pin function. Set the contents of this register through register A with the TI1A instruction. The TAI1 instruction can be used to transfer the contents of register I1 to register A.

Table 8 External interrupt control register
---------------------------------------------

	Interrupt control register I1		reset : 00002	at RAM back-up : state retained	R/W	
11.0	Not used	0	This hit has no fund	tion but road/write is enabled		
113	Not used	1	This bit has no function, but read/write is enabled.			
		0	Falling waveform ("	L" level of INT pin is recognized with t	he SNZI0	
110	Interrupt valid waveform for INT pin/return level selection bit (Note 2)		instruction)/"L" level			
112			Rising waveform ("H" level of INT pin is recognized with the SNZI			
			instruction)/"H" leve	el		
14.	Netwood	0	This hit has no fund	tion but rood/write is enabled		
I11	Not used	1	This bit has no tune	ction, but read/write is enabled.		
14.						
<b>I1</b> 0	Not used	1	This bit has no function, but read/write is enabled.			

Notes 1: "R" represents read enabled, and "W" represents write enabled.

2: Depending on the input state of P21/INT pin, the external interrupt request flag EXF0 may be set to "1" when the contents of I12 is changed. Accordingly, set a value to bit 2 of register I1 and execute the SNZ0 instruction to clear the EXF0 flag after executing at least one instruction.

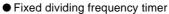
SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

#### TIMERS

The 4570 Group has the programmable timers and a fixed dividing frequency timer.

• Programmable timer

The programmable timer has a reload register and enables the frequency dividing ratio to be set. It is decremented from a set value n. When it underflows (count to n + 1), a timer interrupt request flag is set to "1," new data is loaded from the reload register, and count continues (auto-reload function).



The fixed dividing frequency timer has the fixed frequency dividing ratio (n). An interrupt request flag is set to "1" every n count of a count pulse.

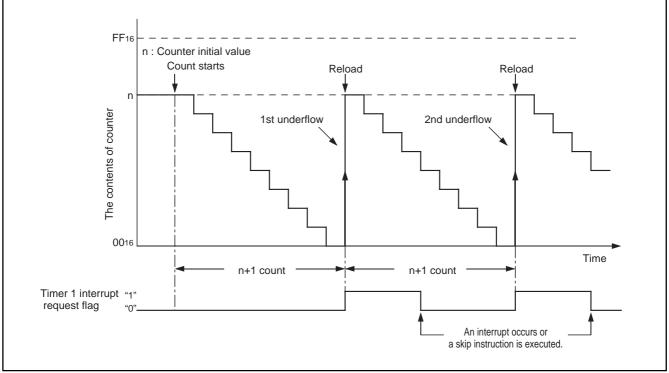


Fig. 18 Auto-reload function

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

The 4570 Group timer consists of the following circuits.

- Prescaler : frequency divider
- Timer 1 : 10-bit programmable timer with the interrupt function and the carrier wave output auto-control function
- Timer 2 : 8-bit programmable timer with the interrupt function
- Timer 3 : 8-bit programmable timer with the interrupt function and the carrier wave generation function
- 16-bit timer

Prescaler, timer 1, timer 2 and timer 3 can be controlled with the timer control registers W1, W2 and W3.

16-bit timer is the free-run counter without the control register. Each function is described below.

Circuit Structure Count source		Count course	Frequency	Use of output signal	Control
Circuit	Siruciure		dividing ratio	Ose of output signal	register
Prescaler	Frequency divider	Instruction clock	4, 8	• Timer 1, 2 and 3 count sources	W1
Timer 1	10-bit programmable	Prescaler output (ORCLK)	1 to 1024	Timer 1 interrupt	W1
	binary down counter	• Carrier wave generating circuit		Carrier wave output auto-control	(W5)
		output (CARRY)		Timer 2 count source	
Timer 2	8-bit programmable	Prescaler output (ORCLK)	1 to 256	Timer 2 interrupt	W2
	binary down counter	Timer 1 underflow		Timer 3 count source	
		<ul> <li>Instruction clock</li> </ul>		• Tout output	
		<ul> <li>16-bit timer underflow</li> </ul>			
Timer 3	8-bit programmable	Prescaler output (ORCLK)	1 to 256	Timer 3 interrupt	W3
	binary down counter	Timer 2 underflow		Timer 1 count source	
		• f(Xin) or f(Xin)/2		Carrier wave	
16-bit timer	16-bit fixed	Instruction clock	65536	Watchdog timer	
	dividing frequency			(15-th bit output is counted	
				twice.)	
				Timer 2 count source	
				(16-bit timer underflow)	

#### **Table 9 Function related timers**

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

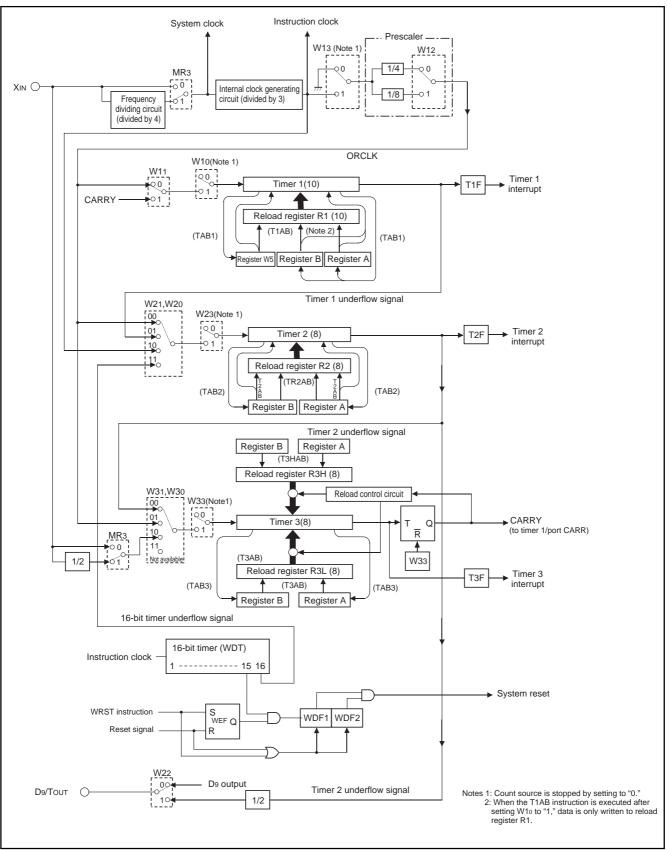


Fig. 19 Timers structure

#### SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

Table 10	Timer control registers				
	Timer control register W1		t reset : 00002	at RAM back-up : 00002	R/W
10/4	W13 Prescaler control bit		Stop (prescaler state initialized)		
VV 13			Operating		
14/4		0	Instruction clock divided by 4		
W12	Prescaler dividing ratio selection bit	1	Instruction clock divided by 8		
14/4	Time and accurate accuracy and actions bit	0	Prescaler output (C	RCLK)	
W11	Timer 1 count source selection bit	1	Carrier output (CARRY)		
14/4	Time on 4 countred with	0	Stop (state retained	1)	
W10	Timer 1 control bit	1	Operating		

	Timer control register W2	at		reset : 00002	at RAM back-up : state retained	R/W
W23	Timer 2 control bit			Stop (state retained	)	
VVZ3				Operating	Operating	
W22	Port D9/Tout pin function selection bit	0		Port D ₉		
VVZ2	Tort Day 1001 pir function selection bit	1		Tout pin		
		W21	W20		Count source	
W21		0	0	Prescaler output (C	PRCLK)	
	Timer 2 count source selection bits	0	1	Timer 1 underflow s	signal	
W20		1	0	Instruction clock		
		1	1	16-bit timer underfle	ow signal	

	Timer control register W3		at reset : 00002		at RAM back-up : state retained	R/W
W33	Timer 3 control bit			Stop (state retained	)	
VV 33				Operating		
W32	Not used	0		This bit has no function, but read/write is enabled.		
			1			
		W31	W30	Count source		
W31		0	0	Timer 2 underflow s	Timer 2 underflow signal	
	Timer 3 count source selection bits	0	1	Prescaler output (O	Prescaler output (ORCLK)	
W30		1	0	f(XIN) or f(XIN)/2		
		1	1	Not available		

Timer count value store register W5	at reset : 002	at RAM back-up : state retained	R/W			
2-bit register. The contents of the high-order 2 bits (bits 9 and 8) of the 10-bit ROM pattern at address (D2D1D0A3A2A1A0) in page						
p specified by registers D and A is stored in this register W5 with the TABP p instruction.						
In addition, data can be transferred between the low-order 2 bits of register A and this register W5 with the TW5A or TAW5						
instruction. Data can be read/written to/from the high-order 2 bits of timer 1 with the T1AB or TAB1 instruction.						
Note: "B" represents read enabled, and "W" represents write enabled						

Note: "R" represents read enabled, and "W" represents write enabled.

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

#### (1) Timer control registers

• Timer control register W1

Register W1 controls the count source and count operation of timer 1, the frequency dividing ratio and count operation of prescaler. Set the contents of this register through register A with the TW1A instruction. The TAW1 instruction can be used to transfer the contents of register W1 to register A.

Timer control register W2

Register W2 controls the count operation and count source of timer 2 and D₉/TouT pin function. Set the contents of this register through register A with the TW2A instruction. The TAW2 instruction can be used to transfer the contents of register W2 to register A.

Timer control register W3

Register W3 controls the count operation and count source of timer 3. Set the contents of this register through register A with the TW3A instruction. The TAW3 instruction can be used to transfer the contents of register W3 to register A.

• Timer count value store register W5

2-bit register. The contents of the high-order 2 bits (bits 9 and 8) of the 10-bit ROM pattern at address in page p specified by registers D and A is stored in this register W5 with the TABP p instruction.

In addition, data can be transferred between the low-order 2 bits of register A and this register W5 with the TW5A or TAW5 instruction. Data can be read/written to/from the high-order 2 bits of timer 1 with the T1AB or TAB1 instruction.

#### (2) Precautions

Note the following for the use of timers.

Prescaler

Stop the prescaler operation to change its frequency dividing ratio.

- Count source
   Stop timer 1, 2 or 3 counting to chance
- Stop timer 1, 2 or 3 counting to change its count source. Reading the timer count value
- Stop each of the timers and then execute the TAB1, TAB2 or TAB3 instruction to read timer 1, 2 or 3 data.
- Writing to reload register R1
   When writing data to reload register R1 while timer 1 is operating, avoid a timing when timer 1 underflows.
- Writing to reload register R3H When writing data to reload register R3H while timer 3 is operating, avoid a timing when timer 3 underflows.

#### (3) Prescaler

Prescaler is a frequency divider. Its frequency dividing ratio can be selected. The count source of prescaler is the instruction clock.

Use the bit 2 of register W1 to select the prescaler dividing ratio and the bit 3 to start and stop its operation. When the bit 3 of register W1 is cleared to "0," prescaler is initialized, and the output signal (ORCLK) stops.

#### (4) Timer 1 (interrupt function)

Timer 1 is a 10-bit binary down counter with the timer 1 reload register (R1). The 10-bit data can be set in timer 1 through registers A, B and W5. Set bits 0 to 3 to register A, bits 4 to 7 to register B and bits 8 to 9 to register W5 to set data to timer 1. Also, ROM pattern (bits 0 to 9) can be set to registers A, B and W5 with the TABP p instruction. Execute the T1AB instruction to set data in timer 1.

When timer 1 stops, 10-bit data can be set simultaneously in timer 1 and the reload register (R1) with the T1AB instruction. When timer 1 is operating, data can be set only in the reload register (R1) with the T1AB instruction.

When setting the next count data to reload register R1 while timer 1 is operating, be sure to set data before timer 1 underflows.

Timer 1 starts counting after the following process;

- ① set data in timer 1,
- 2 select the count source with bit 1 of register W1,

3 set the bit 0 of register W1 to "1."

Once count is started, when timer 1 underflows (the next count pulse is input after the contents of timer 1 becomes "0"), the timer 1 interrupt request flag (T1F) is set to "1," new data is loaded from reload register R1, and count continues (auto-reload function).

When a value set in reload register R1 is n, timer 1 divides the count source signal by n + 1 (n = 0 to 1023).

Data can be read from timer 1 to registers A, B and W5. Stop counting and then execute the TAB1 instruction to read its data.

#### (5) Timer 2 (interrupt function)

Timer 2 is an 8-bit binary counter with the timer 2 reload register (R2). Data can be set simultaneously in timer 2 and the reload register (R2) with the TAB2 instruction. Also, data can be set only in the reload register (R2) with the TR2AB instruction.

Timer 2 starts counting after following process;

① set data in timer 2,

 $\ensuremath{\textcircled{}}$  select the count source with bits 0 and 1 of register W2,

③ set the bit 3 of register W2 to "1."

Once count is started, when timer 2 underflows (the next count pulse is input after the contents of timer 2 becomes "0"), the timer 2 interrupt request flag (T2F) is set to "1," new data is loaded from reload register R2, and count continues (auto-reload function).

When a value set in reload register R2 is n, timer 2 divides the count source signal by n+1 (n = 0 to 255).

Data can be read from timer 2 to registers A and B with the TAB2 instruction. Stop counting and then execute the TAB2 instruction to read its data.

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

#### (6) Timer 3

Timer 3 is an 8-bit binary down counter with the timer 3 reload registers (R3H, R3L). Data can be set simultaneously in timer 3 and the reload register (R3L) with the T3AB instruction. Data can be set in reload register R3H with the T3HAB instruction.

Timer 3 starts counting after the following process;

- ① set data in timer 3,
- select the count source with the bits 1 and 0 of register W3,
- ③ set the bit 3 of register W3 to "1."

The f(XIN) or f(XIN)/2 is selected as the count source by setting W31 to "1" and W30 to "0."

When the  $f(X_{IN})$  is selected as the system clock (bit 3 of clock control register MR= "0"),  $f(X_{IN})$  is selected as the count source.

When the  $f(X_{IN})/4$  is selected as the system clock (bit 3 of clock control register MR= "1"),  $f(X_{IN})/2$  is selected as the count source.

Once count is started, when timer 3 underflows (the next count pulse is input after the contents of timer 3 become "0"), the timer 3 interrupt request flag (T3F) is set to "1," new data is loaded from reload register R3H, and count coutinues (autoreload function).

When the timer 3 underflows again after auto-reload is performed, the timer 3 interrupt request flag (T3F) is set to "1" and new data is reloaded from the reload register R3L and count continues. Timer 3 reloads data from reload register R3H or R3L alternately every underflow.

When the T3AB instruction is executed while timer 3 is operating, new data is set in timer 3 and reload register R3L, count is started again at the next machine cycle. At the next underflow, data is reloaded from R3H and count continues regardless that auto-reload is performed from reload register R3H or R3L at the previous underflow.

Data can be read from timer 3 through registers A and B. Stop counting and then execute the TAB3 instruction to read its data. Timer 3 can be also used as the carrier wave generating circuit.

#### (7) Timer output pin (D9/Tout)

Timer output pin (D $_9$ /TouT) is used to output the timer 2 underflow signal.

The D₉/Tout pin function can be selected by the bit 2 of register W2.

#### (8) Timer interrupt request flags (T1F, T2F, T3F)

Each timer interrupt request flag is set to "1" when each timer underflows. The state of these flags can be examined with the skip instructions (SNZT1, SNZT2, SNZT3).

Use the interrupt control registers V1 and V2 to select an interrupt or a skip instruction.

An interrupt request flag is cleared to "0" when an interrupt occurs or when the next instruction is skipped with a skip instruction.

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

### WATCHDOG TIMER

Watchdog timer provides a method to reset the system when a program runs wild. Watchdog timer consists of 16-bit timer (WDT), watchdog timer enable flag (WEF), and watchdog timer flags (WDF1, WDF2).

Timer WDT starts downcounting the instruction clocks as the count source immediately after system is released from reset. The underflow signal is generated when the count value reaches "000016." This underflow signal can be used as the timer 2 count source.

When the WRST instruction is executed after system is released from reset, the WEF flag is set to "1." At this time, the watchdog timer starts operating.

When the count value of timer WDT reaches "BFFF16" or "3FFF16," WDF1 flag is set to "1." Then, if the WRST instruction is not executed while the timer WDT counts 32767, the WDF2 flag is set to "1" and the RESET pin outputs "L" level to reset the microcomputer. In software using the watchdog timer, make sure that the WRST instruction is executed in 32766 machine cycles or less in order to keep the microcomputer operating normally. To prevent the watchdog timer from stopping in the event of misoperation, the WEF flag is designed not to be initialized once the WRST instruction has been executed. Note also that, if the WRST instruction is never executed, the watchdog timer does not start.

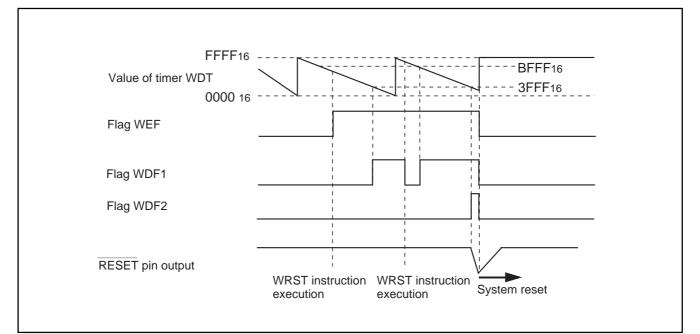
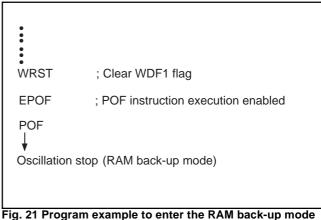


Fig. 20 Watchdog timer function

The contents of the WEF flag, the WDF1 and WDF2 flags and the timer WDT are initialized at the RAM back-up mode. However, if the WDF2 flag is set to "1" at the same time that the microcomputer enters the RAM back-up mode, system reset may be performed.

When using the watchdog timer and the RAM back-up mode, initialize the WDF1 flag with the WRST instruction just before the microcomputer enters the RAM back-up mode (refer to Figure 21).



ig. 21 Program example to enter the RAM back-up mode when using the watchdog timer

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

# **CARRIER WAVE GENERATING CIRCUIT**

The 4570 Group has a carrier wave generating circuit that generates the transfer waveform for various remote control carrier wave.

The carrier wave generating circuit outputs the signal inverted every timer 3 underflow (CARRY) from port CARR.

When using the carrier wave generating circuit, select the f(XIN) or  $f(X_{IN})/2$  for the timer 3 count source (W31="1", W30="0").

When the bit 3 of the clock control register MR is "0" (system  $clock=f(X_{IN})$ ,  $f(X_{IN})$  is selected as the count source.

When the bit 3 of the clock control register MR is "1" (system  $clock=f(X_{IN})/4$ ),  $f(X_{IN})/2$  is selected as the count source.

Set the count value corresponding to "L" interval of carrier wave output to timer 3 reload register R3L.

Set the count value corresponding to "H" interval of carrier wave output to timer 3 reload register R3H.

Also, timer 1 can auto-control the carrier wave output of port CARR by setting the carrier wave output control register (C2).

When timer 3 is stopped, the output level of port CARR is initialized. ("L" level)

### (1) Carrier wave output control register (C2)

Timer 1 can auto-control the output enable interval and the output disable interval of the carrier wave output from port CARR by setting the bit 0 of register C2 to "1." Set the contents of this register through register A with the TC2A instruction.

The setting of the output enable/disable interval is described below.

- ① Validate the carrier wave output auto-control function (C20="1").
- 2 Set the count value ("L" interval of carrier wave output) to timer 3 and reload register R3L.
- 3 Set the count value ("H" interval of carrier wave output) to timer 3 reload register R3H.
- ④ Set the count value (the output enable interval of carrier wave from port CARR) to timer 1.
- ⑤ Select the carrier wave (W11 = "1") as the timer 1 count source.
- 6 Operate timer 1 (W1o="1").
- ⑦ Operate timer 3 (W33="1").
- ® Set the next count value (the output disable interval of carrier wave from port CARR) to reload register R1 before timer 1 underflow occurs.

The carrier wave is output from port CARR until the first timer 3 underflow occurs. The output of the carrier wave from port CARR is disabled and the next count value is loaded from reload register R1 to timer 1 by the first timer 1 underflow.

Then, the output of carrier wave is disabled until the second timer 1 underflow occurs. Also, the next enable interval of the carrier wave output can be set by setting the third count value to timer 1 reload register R1 before the second timer 1 underflow occurs.

If the carrier wave output auto-control function is invalidated (C20="0") while the carrier wave output is auto-controlled, the output of port CARR retains the state when the auto-control is invalidated regardless of timer 1 underflow. This state can be terminated by timer 1 stop (W10="0").

When the carrier wave output auto-control function is validated (C20="1") again after it is invalidated (C20="0"), the autocontrol of carrier wave output is started again when the next timer 1 underflow occurs.

Stop the timer 3 and invalidate the auto-control function by timer 1 to use the port CARR output contorl bit (C21).

### (2) Notes when using the carrier wave output auto-control function

- Set the timer 1 and register C2 before timer 3 is started to operate (W33="1").
- Stop the timer 1 (W10="0") after stopping the timer 3 (W33="0") while the carrier wave output is disabled in order to stop the carrier wave output auto-control operation.
- If the carrier wave output auto-control function is invalidated (C20="0") while the carrier wave output is auto-controlled, the output of port CARR retains the state when the autocontrol is invalidated regardless of timer 1 underflow. When the carrier wave output auto-control function is validated (C20="1") again after it is invalidated (C20="0"), the auto-control by timer 1 is validated again when the next timer 1 underflow occurs.

However, when the carrier wave output auto-control bit (C20) is changed during timer 1 underflow, the error-operation may occur.

When the carrier wave output auto-control function is selected, use the carrier wave CARRY as the timer 1 count source.

If the ORCLK is used as the count source, a short pulse may occur in port CARR output because ORCLK is not synchronized with the carrier wave.

•When the carrier wave output auto-control function is selected and data is set to reload register R1 while timer 1 is operating, avoid the timing that the contents of timer 1 becomes "0" to execute the T1AB instruction.

Ca	Carrier wave output control register C2 at reset : 002		at reset : 002	at RAM back-up : 002	W
C21 Port CARR output control bit		0	Port CARR "L" level output		-
		1	Port CARR "H" leve	el output	
C20			Auto-control output	by timer 1 is invalid	
<b>C</b> 20	Carrier wave output auto-control bit	1 Auto-control output		by timer 1 is valid	

## Table 11 Carrier wave output control register

Note: "W" represents write enabled.

**MITSUBISHI MICROCOMPUTERS** 

# 4570 Group

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

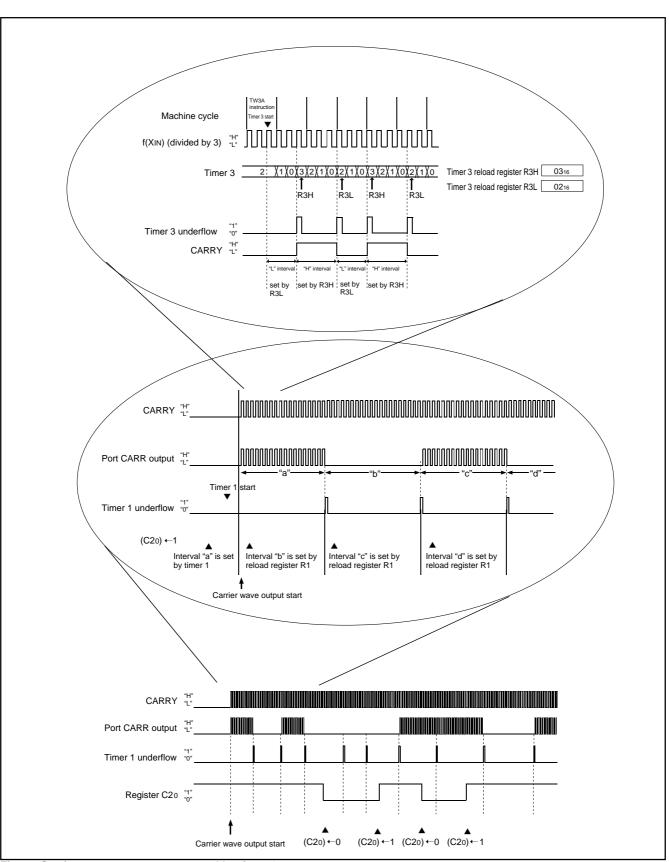


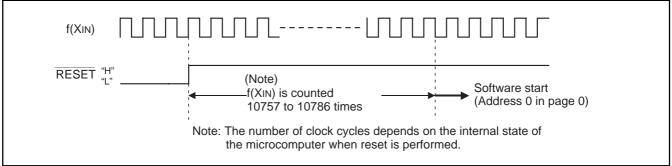
Fig. 22 Carrier wave output auto-control by timer 1

### SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

# **RESET FUNCTION**

System reset is performed by applying "L" level to RESET pin for 1 machine cycle or more when the following condition is satisfied; • the value of supply voltage is the minimum value or more of the recommended operating conditions.

Then when "H" level is applied to  $\overrightarrow{\text{RESET}}$  pin, software starts from address 0 in page 0.



### Fig. 23 Reset release timing

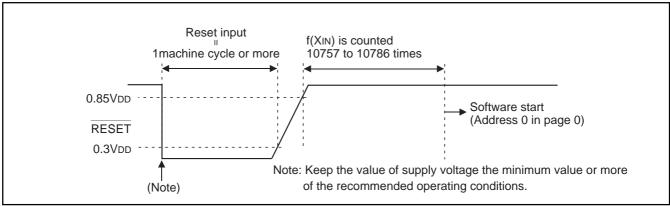


Fig. 24 RESET pin input waveform and reset operation

### (1) Power-on reset

Reset can be automatically performed at power on (poweron reset) by the built-in power-on reset circuit. When the builtin power-on reset circuit is used, the time for the supply voltage to reach the minimum operating voltage must be set to 100  $\mu s$  or less. If the rising time exceeds 100  $\mu s$ , connect a capacitor between the  $\overline{\text{RESET}}$  pin and Vss at the shortest distance, and input "L" level to  $\overline{\text{RESET}}$  pin until the value of supply voltage reaches the minimum operating voltage.

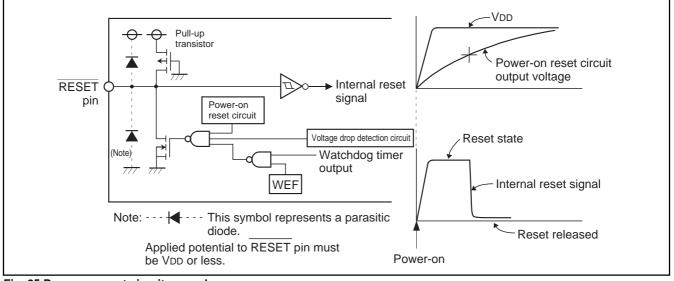


Fig. 25 Power-on reset circuit example

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

The contents of timers, registers, flags and RAM except those

shown in Figure 26 are undefined, so set the initial values to

## (2) Internal state at reset

Table 12 shows port state at reset, and Figure 26 shows internal state at reset (they are retained after system is released from reset).

### Table 12 Port state at reset

Name	Function	State
D0-D8, D9/Tout	D0-D8, D9	High impedance (Note 1)
P00-P03	P00–P03	(11) ()/> laval (blacks 4)
P10-P13	P10–P13	"H" (VDD) level (Note 1)
P20, P21/INT	P20, P21	High impedance
P30–P33	P30–P33	High impedance (Note 1)
P40-P43	P40-P43	High impedance (Note 2)
CARR	CARR	"L" (Vss) level

them.

Notes 1: Output latch is set to "1."

2: The pull-up transistor is turned off.

• Program counter (PC)
Program counter (PC)     0 0 0 0 0 0 0 0 0 0 0 0 0 0
Interrupt enable flag (INTE)     Interrupt enable flag (INTE)
Power down flag (P)
External 0 interrupt request flag (EXF0)
Interrupt control register I1     O O O O
• Timer 1 interrupt request flag (T1F)
• Timer 2 interrupt request flag (T2F)
• Timer 3 interrupt request flag (T3F)
• Watchdog timer flags (WDF1, WDF2)0
• Watchdog timer enable flag (WEF)
• Timer control register W1
• Timer control register W2
• Timer control register W3
Timer count value store register W5
Clock control register MR
• 8-bit general-purpose register SI
Carrier wave output control register C2
Key-on wakeup control register K0
Pull-up control register PU0     0 0 0 0
• Carry flag (CY)0
• Register A
• Register B
• Register D X X X
• Register E
• Register X
• Register Y
• Register Z
• Stack pointer (SP) "X" represents undefined.

Fig. 26 Internal state at reset

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

### **VOLTAGE DROP DETECTION CIRCUIT**

The built-in voltage drop detection circuit is designed to detect a drop in voltage and to reset the microcomputer if the supply voltage drops below a set value.

The voltage drop detection circuit is not operated at the RAM back-up mode.

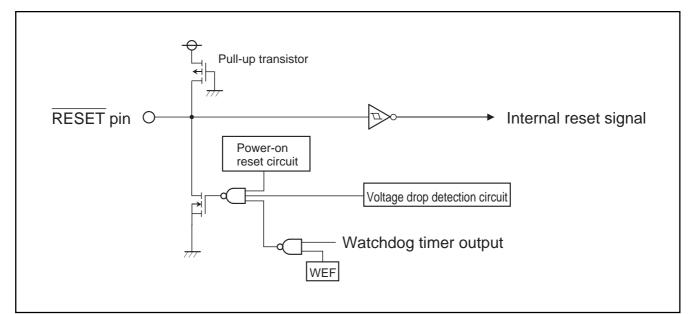


Fig. 27 Voltage drop detection reset circuit

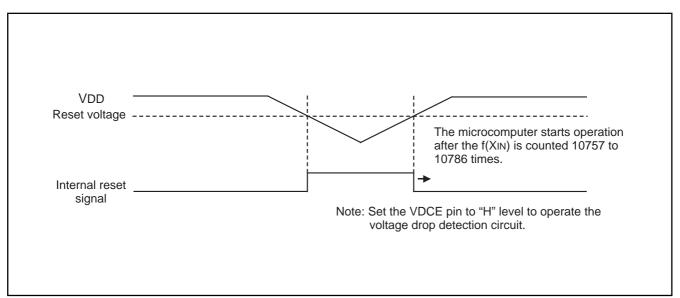


Fig. 28 Voltage drop detection circuit operation waveform

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

# **RAM BACK-UP MODE**

The 4570 Group has the RAM back-up mode.

When the EPOF and POF instructions are executed continuously, system enters the RAM back-up state.

The POF instruction is equivalent to the NOP instruction when the EPOF instruction is not executed before the POF instruction. As oscillation is stopped retaining RAM, the function of reset circuit and states at RAM back-up mode, power dissipation can be reduced without losing the contents of RAM.

Table 13 shows the function and states retained at RAM backup. Figure 29 shows the state transition.

### (1) Identification of the start condition

Warm start (return from the RAM back-up mode) or cold start (return from the normal reset state) can be identified by examining the state of the power down flag (P) with the SNZP instruction.

### (2) Warm start condition

When the external wakeup signal is input after the system enters the RAM back-up mode by executing the EPOF and POF instructions continuously, the CPU starts executing the software from address 0 in page 0. In this case, the P flag is "1."

### (3) Cold start condition

The CPU starts executing the software from address 0 in page 0 when;  $\label{eq:cpu}$ 

- reset pulse is input to  $\overline{\text{RESET}}$  pin, or
- reset by watchdog timer is performed, or
- voltage drop detection circuit detects the voltage drop. In this case, the P flag is "0."

### Table 13 Functions and states retained at RAM back-up

Function	RAM back-up	
Program counter (PC), registers A, B,	×	
carry flag (CY), stack pointer (SP) (Note 2)	×	
Contents of RAM	0	
Port level	0	
Clock control register MR	0	
Timer control register W1	×	
Timer control registers W2, W3	0	
Timer count value store register W5	0	
Interrupt control registers V1, V2	×	
Interrupt control register I1	0	
Carrier wave output control register C2	×	
8-bit general-purpose register SI	0	
Timer 1 function	×	
Timer 2 function	(Note 3)	
Timer 3 function	(Note 3)	
Pull-up control register PU0	0	
Key-on wakeup control register K0	0	
External 0 interrupt request flag (EXF0)	×	
Timer 1 interrupt request flag (T1F)	×	
Timer 2 interrupt request flag (T2F)	(Note 3)	
Timer 3 interrupt request flag (T3F)	(Note 3)	
Watchdog timer flag 1 (WDF1)	X (Note 4)	
Watchdog timer flag 2 (WDF2)	X (Note 4)	
Watchdog timer enable flag (WEF)	X (Note 4)	
16-bit timer (WDT)	X (Note 4)	
Interrupt enable flag (INTE)	×	

Notes 1: "O" represents that the function can be retained, and "X" represents that the function is initialized. Registers and flags other than the above are undefined

at RAM back-up, and set an initial value after returning.2: The stack pointer (SP) points the level of the stack register and is initialized to "1112" at RAM back-up.

- 3: The state of the timer is undefined.
- 4: Initialize the watchdog timer with the WRST instruction, and then execute the EPOF and POF instructions.

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

#### (4) Return signal

An external wakeup signal is used to return from the RAM back-up mode because the oscillation is stopped. Table 14 shows the return condition for each return source.

#### (5) Port P4 control registers

• Key-on wakeup control register K0 Register K0 controls the port P4 key-on wakeup function. Set the contents of this register through register A with the TK0A instruction. In addition, the TAK0 instruction can be used to transfer the contents of register K0 to register A.

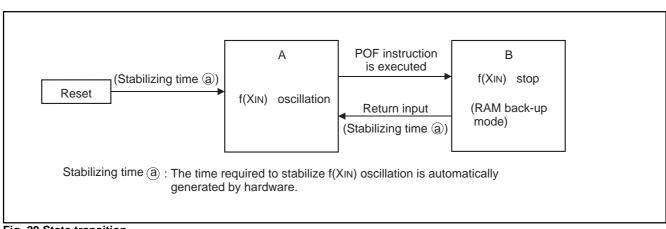
#### Table 14 Return source and return condition

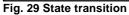
• Pull-up control register PU0

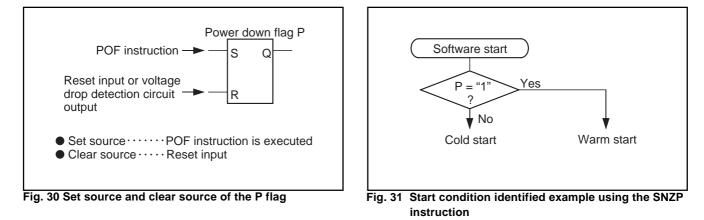
Register PU0 controls the ON/OFF of the port P4 pull-up transistor. Set the contents of this register through register A with the TPU0A instruction. In addition, the TAPU0 instruction can be used to transfer the contents of register PU0 to register A.

Table					
Re	turn source	Return condition	Remarks		
	Ports P0, P1	Return by an external falling edge	Port P0 shares the falling edge detection circuit with ports P1 and P4.		
signal	and P4	input ("H"→"L").	Key-on wakeup functions of ports P0 and P1 are always valid. The key-		
			on wakeup function valid/invalid of port P4 can be controlled with register		
wakeup			K0. Set the port using the key-on wakeup function selected to "H" level		
			before going into the RAM back-up mode.		
External	P21/INT pin	Return by an external "H" level or	Select the return level ("L" level or "H" level) with the bit 2 of register I1		
ter		"L" level input.	according to the external state before going into the RAM back-up mode.		
l û		The EXF0 flag is not set.			

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER







#### Table 15 Key-on wakeup control register and pull-up control register

	Key-on wakeup control register K0		reset : 00002	at RAM back-up : state retained	R/W
K03	Port P43 key-on wakeup	0	0     Key-on wakeup not used       1     Key-on wakeup used		
K03	control bit	1			
KOs	Port P42 key-on wakeup	0	Key-on wakeup not used		
K02	control bit	1	1 Key-on wakeup used		
KO.	Port P41 key-on wakeup	0	0     Key-on wakeup not used       1     Key-on wakeup used		
K01	control bit	1			
KOs	Port P40 key-on wakeup	0	Key-on wakeup not used		
K00	control bit	1	Key-on wakeup use	d	

Pull-up control register PU0		at reset : 00002		at RAM back-up : state retained	R/W
DUID	Port P4 ₃ pull-up transistor	0 Pull-up transistor OFF			
PU03	control bit	1	Pull-up transistor Ol	N	
DU O.	Port P42 pull-up transistor	0	Pull-up transistor OFF Pull-up transistor ON		
PU02	control bit	1			
DUO.	Port P41 pull-up transistor	0	Pull-up transistor OFF		
PU01	control bit	1	Pull-up transistor ON		
PU00	Port P40 and P01 pull-up transistor	0	0 Pull-up transistor OFF		
P000	control bit 1 Pull-up transistor O		N		

Note: "R" represents read enabled, and "W" represents write enabled.

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

# **CLOCK CONTROL**

The clock control circuit consists of the following circuits.

- Clock generating circuit
- Control circuit to stop the clock oscillation
- System clock selection circuit
- Instruction clock generating circuit
- Control circuit to return from the RAM back-up mode

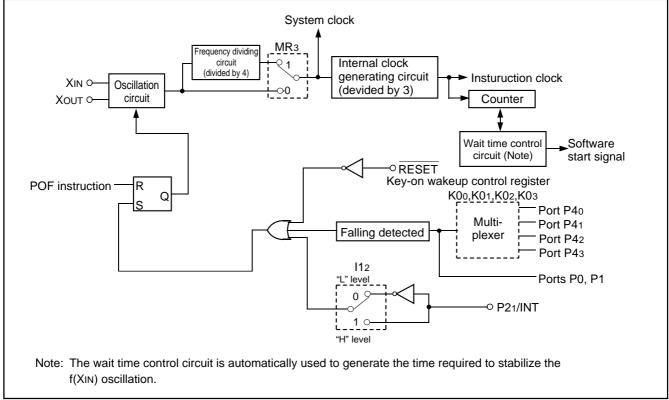


Fig. 32 Clock control circuit structure

Clock signal  $f(X_{IN})$  is obtained by externally connecting a ceramic resonator. Connect this external circuit to pins  $X_{IN}$  and  $X_{OUT}$  at the shortest distance. A feedback resistor is built-in between pins  $X_{IN}$  and  $X_{OUT}$ .

## **ROM ORDERING METHOD**

Please submit the information described below when ordering Mask ROM.

- M34570M4-XXXFP Mask ROM Order Confirmation Form, M34570M8-XXXFP Mask ROM Order Confirmation Form, or M34570MD-XXXFP Mask ROM Order Confirmation Form
- (2) Data to be written into mask ROM ...... EPROM (three sets containing the identical data)
- (3) Mark Specification Form ...... 1

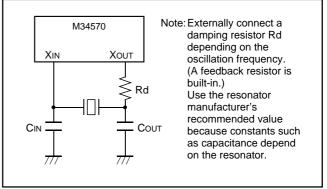


Fig. 33 Ceramic resonator external circuit

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

# LIST OF PRECAUTIONS

① Noise and latch-up prevention

Connect a capacitor on the following condition to prevent noise and latch-up;

- connect a capacitor (approx. 0.1  $\mu\text{F})$  between pins VDD and Vss at the shortest distance,
- · equalize its wiring in width and length, and

use the thickest wire.

In the One Time PROM version, CNVss pin is also used as VPP pin. Accordingly, when using this pin, connect this pin to Vss through a resistor about 5 k $\Omega$  (connect this resistor to CNVss/VPP pin as close as possible).

#### 2 Prescaler

Stop the prescaler operation to change its frequency dividing ratio.

3 Count source

Stop timer 1, timer 2 or timer 3 counting to change its count source.

# ④ Reading the timer count value

Stop each of the timers and then execute the TAB1, TAB2 or TAB3 instruction to read timer 1, 2 or 3 data.

#### ^⑤Writing to reload register R1

When writing the data to reload register R1 while timer 1 is operating, avoid a timing when timer 1 underflows.

#### [®] Writing to reload register R3H

When writing the data to reload register R3H while timer 3 is operating, avoid a timing when timer 3 underflows.

#### ⑦ Notes on timer 3 operation start

Set the timer 1 and register C2 before timer 3 is started to operate (W3₃="1").

® Notes on carrier wave output auto-control operation stop

Stop the timer 1 (W10=" $^{\circ}$ ") after stopping the timer 3 (W33=" $^{\circ}$ ") while the carrier wave output is disabled in order to stop the carrier wave output auto-control operation.

#### Notes on setting carrier wave output control regiter C2

If the carrier wave output auto-control function is invalidated (C2₀="0") while the carrier wave output is auto-controlled, the output of port CARR retains the state when the auto-control is invalidated regardless of timer 1 underflow.

When the carrier wave output auto-control function is validated  $(C2_0="1")$  again after it is invalidated  $(C2_0="0")$ , the auto-control by timer 1 is validated again when the next timer 1 underflow occurs.

However, when the carrier wave output auto-control bit (C2₀) is changed during timer 1 underflow, the error-operation may occur.

#### Notes on timer 1 count source

When the carrier wave output auto-control function is selected, use the carrier wave CARRY as the timer 1 count source. If the ORCLK is used as the count source, a short pulse may occur in port CARR output because ORCLK is not synchronized with the carrier wave.

Notes on writing to reload register R1 when carrier wave output auto-control operation

When the carrier wave output auto-control function is selected and data is set to reload register R1 while timer 1 is operating, avoid the timing that the contents of timer 1 becomes "0" to execute the T1AB instruction.

#### ¹²One Time PROM version

The operating power voltage of the One Time PROM version is within the range of 2.5 V to 5.5 V.

#### ¹³ Multifunction

Note that the port D₉ output function and P2₁ input function can be used even when Tout and INT pin function is selected.

#### ¹ POF instruction

Note that system cannot enter the RAM back-up state when executing only the POF instruction.

Execute the POF instruction immediately after executing the EPOF instruction to enter the RAM back-up.

Be sure to disable interrupts by executing the DI instruction before executing the EPOF instruction.

#### 15 Program counter

Make sure that the PCH does not specify after the last page of the built-in ROM.

#### 16 P21/INT pin

When the interrupt valid waveform of P21/INT pin is changed with the bit 2 of register I1 in software, be careful about the following notes.

- Clear the bit 0 of register V1 to "0" and then change the interrupt valid waveform of P21/INT pin with the bit 2 of register I1 (refer to Figure 34⁽¹⁾).
- Clear the bit 2 of register I1 to "0" and execute the SNZ0 instruction to clear the EXF0 flag after executing at least one instruction (refer to Figure 34⁽²⁾). Depending on the input state of the P21/INT pin, the external 0 interrupt request flag (EXF0) may be set to "1" when the interrupt valid waveform is changed.

:	X : this bit is not related to the setting of INT.
NOP	
SNZ0	;The SNZ0 instruction is executed
NOP	2
TI1A	; Change of the interrupt valid waveform
LA	4
TV1A	; The SNZ0 instruction is valid ①
LA	4 ; ( <b>XXX</b> 02)
:	

Fig. 34 External 0 interrupt program example

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

## SYMBOL

The symbols shown below are used in the following list of instruction function and machine instructions.

Symbol	Contents	Symbol	Contents
A	Register A (4 bits)	WDF1	Watchdog timer flag 1
В	Register B (4 bits)	WDF2	Watchdog timer flag 2
DR	Register D (3 bits)	WEF	Watchdog timer enable flag
E	Register E (8 bits)	INTE	Interrupt enable flag
C2	Carrier wave output control register C2 (2 bits)	EXF0	External 0 interrupt request flag
SI	8-bit general-purpose register SI (8 bits)	Р	Power down flag
V1	Interrupt control register V1 (4 bits)		
V2	Interrupt control register V2 (4 bits)	D	Port D (10 bits)
11	Interrupt control register I1 (4 bits)	P0	Port P0 (4 bits)
W1	Timer control register W1 (4 bits)	P1	Port P1 (4 bits)
W2	Timer control register W2 (4 bits)	P2	Port P2 (2 bits)
W3	Timer control register W3 (4 bits)	P3	Port P3 (4 bits)
W5	Timer count value store register W5 (2 bits)	P4	Port P4 (4 bits)
K0	Key-on wakeup control register K0 (4 bits)	x	Hexadecimal variable
PU0	Pull-up control register PU0 (4 bits)	у	Hexadecimal variable
MR	Clock control register MR (4 bits)	z	Hexadecimal variable
х	Register X (4 bits)	р	Hexadecimal variable
Y	Register Y (4 bits)	n	Hexadecimal constant which represents the
Z	Register Z (2 bits)		immediate value
DP	Data pointer (10 bits)	i	Hexadecimal constant which represents the
	(It consists of registers X, Y, and Z)		immediate value
PC	Program counter (14 bits)	j	Hexadecimal constant which represents the
РСн	High-order 7 bits of program counter		immediate value
PC∟	Low-order 7 bits of program counter	A3A2A1A0	Binary notation of hexadecimal variable A
SK	Stack register (14 bits X 8)		(same for others)
SP	Stack pointer (3 bits)		
CY	Carry flag	<i>←</i>	Direction of data movement
R1	Timer 1 reload register	$\leftrightarrow$	Data exchange between a register and memory
R2	Timer 2 reload register	?	Decision of state shown before "?"
R3H	Timer 3 reload register	()	Contents of registers and memories
R3L	Timer 3 reload register	_	Negate, Flag unchanged after executing
T1	Timer 1		instruction
T2	Timer 2	M(DP)	RAM address pointed by the data pointer
ТЗ	Timer 3	а	Label indicating address a6 a5 a4 a3 a2 a1 a0
T1F	Timer 1 interrupt request flag	р, а	Label indicating address a6 a5 a4 a3 a2 a1 a0
T2F	Timer 2 interrupt request flag		in page p5 p4 p3 p2 p1 p0
	The second state of the se		Hex. C + Hex. number x (also same for others)
T3F	Timer 3 interrupt request flag	C	
T3F	Timer 3 interrupt request flag	+	

Note : The 4570 Group just invalidates the next instruction when a skip is performed. The contents of program counter is not increased by 2. Accordingly, the number of cycles does not change even if skip is not performed. However, the cycle count becomes "1" if the TABP p, RT, or RTS instruction is skipped.

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

## LIST OF INSTRUCTION FUNCTION

	nction Groupir	9 Mnemonic	Function	Grouping	Mnemonic	Function
TAB $(A) \leftarrow (B)$ TBA $(B) \leftarrow (A)$	ansfer	XAMI j	$\begin{array}{l} (A) \leftarrow \rightarrow (M(DP)) \\ (X) \leftarrow (X)EXOR(j) \\ j = 0 \text{ to } 15 \end{array}$		SB j	$(Mj(DP)) \leftarrow 1$ j = 0 to 3
TAY $(B) \leftarrow (A)$ $(A) \leftarrow (Y)$	RAM to register transfer	TMA j	$(Y) \leftarrow (Y) + 1$ $(M(DP)) \leftarrow (A)$	Bit operation	RB j	$(Mj(DP)) \leftarrow 0$ j = 0 to 3
TYA $(Y) \leftarrow (A)$	RAM to		$(M(DF)) \leftarrow (A)$ $(X) \leftarrow (X)EXOR(j)$ $j = 0 \text{ to } 15$	Bit	SZB j	(Mj(DP)) = 0 ? j = 0 to 3
$\begin{array}{ c c c c } TEAB & (E7-E4) \leftarrow \\ (E3-E0) \leftarrow \end{array}$		LA n	(A) ← n n = 0 to 15	иц	SEAM	(A) = (M(DP)) ?
TABE(B) $\leftarrow$ (E7- (A) $\leftarrow$ (E3-		TABP p	$(SP) \leftarrow (SP) + 1$ $(SK(SP)) \leftarrow (PC)$	Comparison operation	SEA n	(A) = n ? n = 0 to 15
TDA (DR2–DR0)	) ← (A2–A0)		$(PCH) \leftarrow p$ $(PCL) \leftarrow (DR_2-DR_0,$		Ва	(PC∟) ← a6–a0
TAD $(A_2-A_0) \leftarrow (A_3) \leftarrow 0$	(DR2–DR0)		$\begin{array}{l} A_{3}-A_{0})\\ (W5)\leftarrow(ROM(PC))_{9 \text{ to } 8}\\ (B)\leftarrow(ROM(PC))_{7 \text{ to } 4}\end{array}$	Branch operation	BL p, a	(РСн)
TAZ $(A_1, A_0) \leftarrow (A_3, A_2) \leftarrow$			$\begin{array}{l} (A) \leftarrow (ROM(PC))_{3 \text{ to } 0} \\ (PC) \leftarrow (SK(SP)) \\ (SP) \leftarrow (SP) - 1 \end{array}$	Branch	BLA p	$(PCH) \leftarrow p$ $(PCL) \leftarrow (DR_2-DR_0)$ $A_3-A_0)$
TAX $(A) \leftarrow (X)$		AM	(A) ← (A) + (M(DP))		BM a	(SP) ← (SP) + 1
TASP $(A_2-A_0) \leftarrow$ $(A_3) \leftarrow 0$	(SP ₂ –SP ₀ ) (SP ₂ –SP ₀ ) = 0 to 15 = 0 to 15	AMC	$(A) \leftarrow (A) + (M(DP)) + (CY)$			$(SK(SP)) \leftarrow (PC)$ $(PCH) \leftarrow 2$ $(PCL) \leftarrow a_6-a_0$
$\begin{array}{c c} LXY \ x, \ y & (X) \leftarrow x, \ x \\ & (Y) \leftarrow y, \ y \end{array}$	= 0 to 15 = 0 to 15 = 0 to 15	An	$(CY) \leftarrow Carry$ $(A) \leftarrow (A) + n$	eration	BML p, a	(SP) ← (SP) + 1 (SK(SP)) ← (PC)
$LZ z$ (Z) $\leftarrow$ z, z	= 0 to 3		n = 0 to 15	Subroutine operation		(РСн) ← р (РСL) ← а6–ао
$ INY   (Y) \leftarrow (Y) -$	+ 1	AND	$(A) \leftarrow (A)AND(M(DP))$	Subrou	BMLA p	(SP) ← (SP) + 1
DEY $(Y) \leftarrow (Y)$ -	- 1	OR	$(A) \gets (A)OR(M(DP))$			(SK(SP)) ← (PC) (РСн) ← р
TAM j $(A) \leftarrow (M(D))$ $(X) \leftarrow (X)E$		SC	(CY) ← 1			(PCL) ← (DR2–DR0, A3–A0)
j = 0 to 15		RC	$(CY) \leftarrow 0$		RTI	(PC) ← (SK(SP))
XAM j(A) $\leftarrow \rightarrow$ (I(X) $\leftarrow$ (X)E		SZC	(CY) = 0 ?			(SP) ← (SP) − 1
j = 0 to 15		СМА	$(\overline{A}) \rightarrow (\overline{A})$	peratior	RT	$(PC) \leftarrow (SK(SP))$ $(SP) \leftarrow (SP) - 1$
(X) ← (X)E j = 0 to 15	XOR(j)	RAR	→CY→A3A2A1A0	Return op	RTS	$(PC) \leftarrow (SK(SP))$ $(SP) \leftarrow (SP) - 1$
XAMD j (A	A) ← → (I <) ← (X)E = 0 to 15	$A) \leftarrow \rightarrow (M(DP))$ $A) \leftarrow (X)EXOR(j)$	$ \begin{array}{c} A) \leftarrow \rightarrow (M(DP)) \\ A) \leftarrow (X)EXOR(j) \\ = 0 \text{ to } 15 \end{array} \end{array} $	$A) \leftarrow \rightarrow (M(DP))$ $A) \leftarrow (X)EXOR(j)$ = 0  to  15 $RAR \qquad $	$A) \leftarrow \rightarrow (M(DP))$ $A) \leftarrow (X)EXOR(j)$ = 0  to  15 $RAR \rightarrow CY \rightarrow A_{3}A_{2}A_{1}A_{0} \qquad G$	

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

## LIST OF INSTRUCTION FUNCTION (CONTINUED)

Grouping	Mnemonic	Function	Grouping	Mnemonic	Function	Grouping	Mnemonic	Function
	DI	$(INTE) \leftarrow 0$		TAW1	$(A) \gets (W1)$		TAB3	(B) ← (T37–T34)
	EI	(INTE) ← 1		TW1A	(W1) ← (A)			(A) ← (T33–T30)
				10017	(WT) <= (A)		ТЗАВ	(R3L7–R3L4) ← (B)
	SNZ0	(EXF0) = 1 ?		TAW2	(A) ← (W2)		-	(T37−T34) ← (B)
		After skipping the next						$(R3L_3-R3L_0) \leftarrow (A)$
		instruction,		TW2A	$(W2) \gets (A)$			$(T3_3T3_0) \leftarrow (A)$
		(EXF0) ← 0		TANKO	(4) (14(0)		TOLINE	
ç	SNZI0	I12 = 1 : (INT0) = "H" ?		TAW3	$(A) \leftarrow (W3)$		ТЗНАВ	$(R3H7-R3H4) \leftarrow (B)$ $(R3H3-R3H_0) \leftarrow (A)$
ratio	GINZIO	I12 = 0 : (INT0) = "L"?		ТѠЗА	(W3) ← (A)	_		
ope						atior	SNZT1	(T1F) = 1 ?
upt	TAV1	$(A) \leftarrow (V1)$		TAW5	$(A) \leftarrow (0, 0, W51, W50)$	pera		After skipping the nex
Interrupt operation						Timer operation		instruction,
-	TV1A	$(V1) \leftarrow (A)$		TW5A	(W51, W50) ← (A1, A0)	Tim		(T1F) ← 0
	TAV2	(A) ← (V2)		TAB1	(W5) ← (T19–T18)		SNZT2	(T2F) = 1 ?
					(B) ← (T17–T14)			After skipping the nex
	TV2A	$(V2) \leftarrow (A)$			(A) ← (T13–T10)			instruction,
								(T2F) ← 0
	TAI1	$(A) \leftarrow (I1)$		T1AB	at timer 1 stop (W1₀=0) (R19–R1ଃ) ← (W5)		SNZT3	(T3F) = 1 ?
	TI1A	(I1) ← (A)	_		$(T19-T18) \leftarrow (W5)$ $(T19-T18) \leftarrow (W5)$		SINZIS	After skipping the nex
			atio		(R17–R14) ← (B)			instruction,
			Timer operation		$(T17-T14) \gets (B)$			(T3F) ← 0
			ner o		$(R1_3-R1_0) \leftarrow (A)$			
			Ë		(T13-T10) ← (A) At timer 1 operating			
					(W1 ₀ =1),			
					(R19–R18) ← (W5)			
					(R17−R14) ← (B)			
					(R13–R1₀) ← (A)			
				TADO				
				TAB2	$(B) \leftarrow (T27 - T24)$ $(A) \leftarrow (T23 - T20)$			
					(1) (123 120)			
				T2AB	$(R27-R24) \leftarrow (B)$			
					$(T27-T24) \gets (B)$			
					$(R23-R20) \leftarrow (A)$			
					$(T23-T20) \leftarrow (A)$			
				TR2AB	(R27–R24) ← (B)			
					(R23–R20) ← (A)			

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

Grouping	Mnemonic	Function	Grouping	Mnemonic	Function
- Si Supili iy	IAP0	$(A) \leftarrow (P0)$	Siouping	NOP	$(PC) \leftarrow (PC) + 1$
	OP0A	$(P0) \leftarrow (A)$		POF	RAM back-up mode
	IAP1	(A) ← (P1)		EPOF	POF instruction valid
	OP1A	(P1) ← (A)	u	SNZP	(P) = 1 ?
	IAP2	(A1, A0) ← (P21, P20) (A3, A2) ← (0)	Other operation	WRST	$(WDF1) \leftarrow 0, (WEF) \leftarrow 1$
	IAP3	(A) ← (P3)	Other	TAMR	$(A) \leftarrow (MR_3 - MR_0)$
ion	ОРЗА	(P3) ← (A)		TMRA	$(MR_3-MR_0) \leftarrow (A)$
t operat	IAP4	(A) ← (P4)		TABSI	$      (B) \leftarrow (SI_7 - SI_4) \\       (A) \leftarrow (SI_3 - SI_0) $
Input/Output operation	CLD	(D) ← 1		TSIAB	$(SI_7-SI_4) \leftarrow (B)$ $(SI_3-SI_0) \leftarrow (A)$
ndul	RD	$(D(Y)) \leftarrow 0$ (Y) = 0  to  9		SBK	When executing the TABP p instruction,
	SD	(D(Y)) ← 1 (Y) = 0 to 9		RBK	$p_{6} \leftarrow 1$ When executing the
	TK0A	(K0) ← (A)			TABP p instruction, $p_6 \leftarrow 0$
	TAK0	(A) ← (K0)			
	TPU0A	(PU0) ← (A)			
	TAPU0	$(A) \gets (PU0)$			
Carrier wave generating operation	TC2A	(C21, C20) ← (A1, A0)			

## LIST OF INSTRUCTION FUNCTION (CONTINUED)

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

INST	RUC		1 CO	DE T	ABL	E													
	9—D4	000000	000001	000010	000011	000100	000101	000110	000111	001000	001001	001010	001011	001100	001101	001110	001111	1	011000   011111
D3— D0	Hex. notation	00	01	02	03	04	05	06	07	08	09	0A	0B	0C	0D	0E	0F		18—1F
0000	0	NOP	BLA	SZB 0	BMLA	RBK	TASP	A 0	LA 0	TABP 0*	TABP 16*	TABP 32**	TABP 48**	BML	BML	BL	BL	вм	в
0001	1	_	CLD	SZB 1	_	SBK	TAD	A 1	LA 1	TABP 1*	TABP 17*	TABP 33**	TABP 49**	BML	BML	BL	BL	BM	В
0010	2	POF	_	SZB 2	_	_	ТАХ	A 2	LA 2	TABP 2*	TABP 18*	TABP 34**	TABP 50**	BML	BML	BL	BL	вм	в
0011	3	SNZP	INY	SZB 3	_	_	TAZ	A 3	LA 3	TABP 3*	TABP 19*	TABP 35**	TABP 51**	BML	BML	BL	BL	вм	в
0100	4	DI	RD	_	_	RT	TAV1	A 4	LA 4	TABP 4*	TABP 20*	TABP 36**	TABP 52**	BML	BML	BL	BL	вм	В
0101	5	EI	SD	SEAn	_	RTS	TAV2	A 5	LA 5	TABP 5*	TABP 21*	TABP 37**	TABP 53**	BML	BML	BL	BL	ВМ	В
0110	6	RC	_	SEAM	_	RTI	_	A 6	LA 6	TABP 6*	TABP 22*	TABP 38**	TABP 54**	BML	BML	BL	BL	вм	В
0111	7	SC	DEY	_	_	-	_	A 7	LA 7	TABP 7*	TABP 23*	TABP 39**	TABP 55**	BML	BML	BL	BL	вм	в
1000	8	_	AND	_	SNZ0	LZ 0	_	A 8	LA 8	TABP 8*	TABP 24*	TABP 40**	TABP 56**	BML	BML	BL	BL	вм	в
1001	9	_	OR	TDA	_	LZ 1	_	A 9	LA 9	TABP 9*	TABP 25*	TABP 41**	TABP 57**	BML	BML	BL	BL	вм	в
1010	А	AM	TEAB	TABE	SNZI0	LZ 2	—	A 10	LA 10	TABP 10*	TABP 26*	TABP 42**	TABP 58**	BML	BML	BL	BL	вм	в
1011	в	AMC	_	_	_	LZ 3	EPOF	A 11	LA 11	TABP 11*	TABP 27*	TABP 43**	TABP 59**	BML	BML	BL	BL	вм	В
1100	С	TYA	СМА	_	_	RB 0	SB 0	A 12	LA 12	TABP 12*	TABP 28*	TABP 44**	TABP 60**	BML	BML	BL	BL	BM	в
1101	D	_	RAR	_	_	RB 1	SB 1	A 13	LA 13	TABP 13*	TABP 29*	TABP 45**	TABP 61**	BML	BML	BL	BL	BM	В
1110	E	ТВА	ТАВ	_	TV2A	RB 2	SB 2	A 14	LA 14	TABP 14*	TABP 30*	TABP 46**	TABP 62**	BML	BML	BL	BL	BM	в
1111	F	_	TAY	szc	TV1A	RB 3	SB 3	A 15	LA 15	TABP 15*	TABP 31*	TABP 47**	TABP 63**	BML	BML	BL	BL	ВМ	в

The above table shows the relationship between machine language codes and machine language instructions. D  $_{3}$ —D $_{0}$  show the low-order 4 bits of the machine language code, and D  $_{9}$ —D4 show the high-order 6 bits of the machine language code. The hexadecimal representation of the code is also provided. There are one-word instructions and two-word instructions, but only the first word of each instruction is shown. Do not use code marked "—."

** cannot be used at M34570M4.

For M34570M4/M8/E8, the SBK and RBK instructions cannot be used.

For M34570MD/ED, the pages which is referred with the TABP instruction (*, **) can be switched with the SBK and RBK instructions. After executing the SBK instruction, the pages which can be referred with the TABP instruction are 64 to 127. (ex. TABP 0  $\rightarrow$ TABP 64) After executing the RBK instruction, the pages which can be referred with the TABP instruction are 0 to 63. If the SBK instruction is not executed, the pages which can be referred with the TABP instruction are always 0 to 63.

The codes for the second word of a two-word instruction are described below.

	The second word
BL	1p paaa aaaa
BML	1p paaa aaaa
BLA	1р ррОО рррр
BMLA	1р ррОО рррр
SEA	00 0111 nnnn
SZD	00 0010 1011

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

NST	RUC	TION	I CO	DE T	ABL	E (C(	ΟΝΤΙ	NUE	D)									
$\mathbb{N}$										101000	101001	101010	101011	101100	101101	101110	101111	110000   111111
D3— D0	Hex. notation	20	21	22	23	24	25	26	27	28	29	2A	2B	2C	2D	2E	2F	30—3F
0000	0	_	TW3A	OP0A	T1AB	_		IAP0	TAB1	SNZT1	_	WRST	TMA 0	TAM 0	XAM 0	XAMI 0	XAMD 0	LXY
0001	1	_	_	OP1A	T2AB	_		IAP1	TAB2	SNZT2	_	_	TMA 1	TAM 1	XAM 1	XAMI 1	XAMD 1	LXY
0010	2	_	TW5A	_	ТЗАВ	_	TAMR	IAP2	ТАВЗ	SNZT3	_	_	TMA 2	TAM 2	XAM 2	XAMI 2	XAMD 2	LXY
0011	3	_		ОРЗА		_	TAI1	IAP3	_		_	_	TMA 3	TAM 3	XAM 3	XAMI 3	XAMD 3	LXY
0100	4	_				_	_	IAP4	_		_	_	TMA 4	TAM 4	XAM 4	XAMI 4	XAMD 4	LXY
0101	5					_	_	_	_	_	_	_	TMA 5	TAM 5	XAM 5	XAMI 5	XAMD 5	LXY
0110	6	_	TMRA			_	TAK0	_	_		_	_	TMA 6	TAM 6	XAM 6	XAMI 6	XAMD 6	LXY
0111	7	_	TI1A	_	_	_	TAPU0	_	_	_	_	_	TMA 7	TAM 7	XAM 7	XAMI 7	XAMD 7	LXY
1000	8	_		_	TSIAB	_		_	TABSI	_	_	_	TMA 8	TAM 8	XAM 8	XAMI 8	XAMD 8	LXY
1001	9	_		_		_		_	_		_	TC2A	TMA 9	TAM 9	XAM 9	XAMI 9	XAMD 9	LXY
1010	А	_			TR2AB	_	_	_	_	_	_	_	TMA 10	TAM 10	XAM 10	XAMI 10	XAMD 10	LXY
1011	В	_	TK0A	_	_	TAW1		_	_		_	_	TMA 11	TAM 11	XAM 11	XAMI 11	XAMD 11	LXY
1100	с	_	_	_	_	TAW2		_	_		_	_	TMA 12	TAM 12	XAM 12	XAMI 12	XAMD 12	LXY
1101	D	_	_	TPU0A	ТЗНАВ	TAW3	_	_	_	_	_	-	TMA 13	TAM 13	XAM 13	XAMI 13	XAMD 13	LXY
1110	Е	TW1A	_	_	_	_	_	_	_	_	_	_	TMA 14	TAM 14	XAM 14	XAMI 14	XAMD 14	LXY
1111	F	TW2A	_	_	_	TAW5	_	_	_	_	_	_	TMA 15	TAM 15	XAM 15	XAMI 15	XAMD 15	LXY
							-	-	-	-		-				-		-

## **INSTRUCTION CODE TABLE (CONTINUED)**

The above table shows the relationship between machine language codes and machine language instructions. D  $_{3}$ -D₀ show the low-order 4 bits of the machine language code, and D  $_{9}$ -D₄ show the high-order 6 bits of the machine language code. The hexadecimal representation of the code is also provided. There are one-word instructions and two-word instructions, but only the first word of each instruction is shown. Do not use code marked "-."

The codes for the second word of a two-word instruction are described below.

	Tł	ne second	l word
BL	1 p	рааа	aaaa
BML	1 p	paaa	aaaa
BLA	1 p	p	рррр
BMLA	1 p	p	рррр
SEA	0 0	0111	nnnn
SZD	0 0	0010	1011

#### SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

## MACHINE INSTRUCTIONS

Parameter						In	struc	ction	coc	de					er of ds	er of es	
Type of instructions	Mnemonic	D9	D8	D7	D6	D5	D4	Dз	D2	D1	D ₀		adec		Number of words	Number of cycles	Function
	ТАВ	0	0	0	0	0	1	1	1	1	0	0	1	E	1	1	(A) ← (B)
	ТВА	0	0	0	0	0	0	1	1	1	0	0	0	Е	1	1	$(B) \leftarrow (A)$
	TAY	0	0	0	0	0	1	1	1	1	1	0	1	F	1	1	$(A) \gets (Y)$
	TYA	0	0	0	0	0	0	1	1	0	0	0	0	С	1	1	$(Y) \gets (A)$
sfer	TEAB	0	0	0	0	0	1	1	0	1	0	0	1	A	1	1	$(E_7-E_4) \leftarrow (B)$ $(E_3-E_0) \leftarrow (A)$
Register to register transfer	TABE	0	0	0	0	1	0	1	0	1	0	0	2	A	1	1	(B) ← (E7–E4) (A) ← (E3–E0)
er to re	TDA	0	0	0	0	1	0	1	0	0	1	0	2	9	1	1	$(DR_2-DR_0) \leftarrow (A_2-A_0)$
Regist	TAD	0	0	0	1	0	1	0	0	0	1	0	5	1	1	1	$(A_2-A_0) \leftarrow (DR_2-DR_0)$ $(A_3) \leftarrow 0$
	TAZ	0	0	0	1	0	1	0	0	1	1	0	5	3	1	1	$\begin{array}{l} (A_1,A_0) \leftarrow (Z_1,Z_0) \\ (A_3,A_2) \leftarrow 0 \end{array}$
	ТАХ	0	0	0	1	0	1	0	0	1	0	0	5	2	1	1	$(A) \gets (X)$
	TASP	0	0	0	1	0	1	0	0	0	0	0	5	0	1	1	$(A_2-A_0) \leftarrow (SP_2-SP_0)$ $(A_3) \leftarrow 0$
	LXY x, y	1	1	Х3	<b>X</b> 2	<b>X</b> 1	<b>X</b> 0	уз	<b>y</b> 2	<b>у</b> 1	yo	3	x	У	1	1	$\begin{array}{l} (X) \leftarrow x,  x = 0 \text{ to } 15 \\ (Y) \leftarrow y,  y = 0 \text{ to } 15 \end{array}$
RAM addresses	LZ z	0	0	0	1	0	0	1	0	<b>Z</b> 1	<b>Z</b> 0	0	4	8 +z	1	1	$(Z) \leftarrow z, z = 0 \text{ to } 3$
RAM	INY	0	0	0	0	0	1	0	0	1	1	0	1	3	1	1	$(Y) \leftarrow (Y) + 1$
	DEY	0	0	0	0	0	1	0	1	1	1	0	1	7	1	1	(Y) ← (Y) − 1

Skip condition	Carry flag CY	
-	_	Transfers the contents of register B to re
-	-	Transfers the contents of register A to re
-	_	Transfers the contents of register Y to re
-	-	Transfers the contents of register A to re
-	_	Transfers the contents of registers A an
_	_	Transfers the contents of register E to re
-	_	Transfers the contents of register A to re
_	_	Transfers the contents of register D to re
-	-	Transfers the contents of register Z to re
-	-	Transfers the contents of register X to re
_	_	Transfers the contents of stack pointer (
Continuous description	_	Loads the value x in the immediate field Y. When the LXY instructions are continuous and other LXY instructions coded contin
-	-	Loads the value z in the immediate field
(Y) = 0	_	Adds 1 to the contents of register Y. As next instruction is skipped.
(Y) = 15	_	Subtracts 1 from the contents of registe is 15, the next instruction is skipped.

# MITSUBISHI MICROCOMPUTERS 4570 Group

## SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

# Detailed description

register A.

register B.

register A.

register Y.

nd B to register E.

registers A and B.

register D.

register A.

register A.

register A.

(SP) to register A.

d to register X, and the value y in the immediate field to register

usly coded and executed, only the first LXY instruction is executed inuously are skipped.

ld to register Z.

as a result of addition, when the contents of register Y is 0, the

er Y. As a result of subtraction, when the contents of register Y

#### SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

#### MACHINE INSTRUCTIONS (CONTINUED)

Parameter		Instruction code											er of ds	er of es	କ୍ଷ୍ Eunction		
Type of instructions	Mnemonic	D9	D8	D7	D6	D5	D4	Dз	D2	D1	D ₀		kadeo otati		Number of words	Number of cycles	Function
	ТАМ ј	1	0	1	1	0	0	j	j	j	j	2	С	j	1	1	$\begin{array}{l} (A) \leftarrow (M(DP)) \\ (X) \leftarrow (X)EXOR(j) \\ j = 0 \text{ to } 15 \end{array}$
	XAM j	1	0	1	1	0	1	j	j	j	j	2	D	j	1	1	$\begin{array}{l} (A) \leftarrow \rightarrow (M(DP)) \\ (X) \leftarrow (X)EXOR(j) \\ j = 0 \text{ to } 15 \end{array}$
RAM to register transfer	XAMD j	1	0	1	1	1	1	j	j	j	j	2	F	j	1	1	$\begin{array}{l} (A) \leftarrow \rightarrow (M(DP)) \\ (X) \leftarrow (X)EXOR(j) \\ j = 0 \text{ to } 15 \\ (Y) \leftarrow (Y) - 1 \end{array}$
RAM to	XAMI j	1	0	1	1	1	0	j	j	j	j	2	E	j	1	1	$\begin{array}{l} (A) \leftarrow \rightarrow (M(DP)) \\ (X) \leftarrow (X)EXOR(j) \\ j = 0 \text{ to } 15 \\ (Y) \leftarrow (Y) + 1 \end{array}$
	TMA j	1	0	1	0	1	1	j	j	j	j	2	В	j	1	1	(M(DP)) ← (A) (X) ← (X)EXOR(j) j = 0 to 15
	LA n	0	0	0	1	1	1	n	n	n	n	0	7	n	1	1	(A) ← n n = 0 to 15
Arithmetic operation	TABP p	0	0	1	0	ръ	р4	рз	p2	p1	po	0	8 +p	р	1	3	$\begin{split} (SP) &\leftarrow (SP) + 1 \\ (SK(SP)) &\leftarrow (PC) \\ (PCH) &\leftarrow p \\ (PCL) &\leftarrow (DR2 - DR0, A3 - A0) \\ (W5) &\leftarrow (ROM(PC))9 \text{ to } 8 \\ (B) &\leftarrow (ROM(PC))7 \text{ to } 4 \\ (A) &\leftarrow (ROM(PC))3 \text{ to } 0 \\ (PC) &\leftarrow (SK(SP)) \\ (SP) &\leftarrow (SP) - 1 \\ (Note) \end{split}$

	Carry flag CY	Skip condition
After transferring the contents of M(DP) to register X and the value j in the immediat	_	-
After exchanging the contents of M(DP) v performed between register X and the val	_	_
After exchanging the contents of M(DP) of performed between register X and the val Subtracts 1 from the contents of register is 15, the next instruction is skipped.	_	(Y) = 15
After exchanging the contents of M(DP) of performed between register X and the val Adds 1 to the contents of register Y. As a next instruction is skipped.	_	(Y) = 0
After transferring the contents of register / register X and the value j in the immediat	_	_
Loads the value n in the immediate field When the LA instructions are continuously and other LA instructions coded continuo	_	Continuous description
Transfers bits 9 and 8 to register W5, bits to 0 are the ROM pattern in address (DF page p. When this instruction is executed, 1 stag When this instruction is executed after ex When this instruction is executed after ex When this instruction is executed after sy pages 0 to 63 are specified.	_	_

Note: p is 0 to 31 for M34570M4 and p is 0 to 63 for M34570E8 and M34570M8.

p is 0 to 127 for M34570ED and M34570MD, and p6 is specified with the SBK and RBK instructions.

# MITSUBISHI MICROCOMPUTERS 4570 Group

#### SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

Detailed description

to register A, an exclusive OR operation is performed between iate field, and stores the result in register X.

v) with the contents of register A, an exclusive OR operation is value j in the immediate field, and stores the result in register X.

) with the contents of register A, an exclusive OR operation is value j in the immediate field, and stores the result in register X. For Y. As a result of subtraction, when the contents of register Y

) with the contents of register A, an exclusive OR operation is value j in the immediate field, and stores the result in register X. s a result of addition, when the contents of register Y is 0, the

er A to M(DP), an exclusive OR operation is performed between iate field, and stores the result in register X.

d to register A.

sly coded and executed, only the first LA instruction is executed uously are skipped.

its 7 to 4 to register B and bits 3 to 0 to register A. These bits 9 DR2 DR1 DR0 A3 A2 A1 A0)2 specified by registers A and D in

age of stack register is used.

executing the SBK instruction, pages 64 to 127 are specified. executing the RBK instruction, pages 0 to 63 are specified. system is released from reset or returned from RAM back-up,

#### SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

## MACHINE INSTRUCTIONS (CONTINUED)

Parameter	r		Instruction code												oer of rds oer of		S Euroction	
Type of instructions	Mnemonic	D9	D8	D7	D6	D5	D4	Dз	D2	D1	Do			cimal ion	Number of words	Number of cycles	Function	
	AM	0	0	0	0	0	0	1	0	1	0	0	0	A	1	1	$(A) \leftarrow (A) + (M(DP))$	
	AMC	0	0	0	0	0	0	1	0	1	1	0	0	В	1	1	$(A) \leftarrow (A) + (M(DP))+ (CY)$ $(CY) \leftarrow Carry$	
	A n	0	0	0	1	1	0	n	n	n	n	0	6	n	1	1	(A) ← (A) + n n = 0 to 15	
Arithmetic operation	AND	0	0	0	0	0	1	1	0	0	0	0	1	8	1	1	$(A) \leftarrow (A)AND(M(DP))$	
Arithmetic	OR	0	0	0	0	0	1	1	0	0	1	0	1	9	1	1	$(A) \leftarrow (A)OR(M(DP))$	
	SC	0	0	0	0	0	0	0	1	1	1	0	0	7	1	1	(CY) ← 1	
	RC	0	0	0	0	0	0	0	1	1	0	0	0	6	1	1	$(CY) \leftarrow 0$	
	SZC	0	0	0	0	1	0	1	1	1	1	0	2	F	1	1	(CY) = 0 ?	
	СМА	0	0	0	0	0	1	1	1	0	0	0	1	С	1	1	$(A) \leftarrow (A)$	
	RAR	0	0	0	0	0	1	1	1	0	1	0	1	D	1	1		
	SB j	0	0	0	1	0	1	1	1	j	j	0	5	C +j	1	1	$(Mj(DP)) \leftarrow 1$ j = 0 to 3	
Bit operation	RB j	0	0	0	1	0	0	1	1	j	j	0	4	C +j	1	1	$(Mj(DP)) \leftarrow 0$ j = 0  to  3	
Bi.	SZB j	0	0	0	0	1	0	0	0	j	j	0	2	j	1	1	(Mj(DP)) = 0 ? j = 0 to 3	
u u	SEAM	0	0	0	0	1	0	0	1	1	0	0	2	6	1	1	(A) = (M(DP)) ?	
Comparison operation	SEA n	0	0	0	0	1	0	0	1	0	1	0	2	5	2	2	(A) = n ? n = 0 to 15	
o C C		0	0	0	1	1	1	n	n	n	n	0	7	n				

Skip condition	Carry flag CY	
-	-	Adds the contents of M(DP) to register A remains unchanged.
_	0/1	Adds the contents of M(DP) and carry flag
Overflow = 0	_	Adds the value n in the immediate field to The contents of carry flag CY remains un Skips the next instruction when there is n
-	-	Performs the AND operation between the the result in register A.
-	-	Performs the OR operation between the the result in register A.
_	1	Sets carry flag CY to "1."
_	0	Clears carry flag CY to "0."
(CY) = 0	-	Skips the next instruction when the conte
-	-	Stores the one's complement for register
-	0/1	Rotates the contents of register A includ
_	-	Sets the contents of bit j (bit specified by
-	-	Clears the contents of bit j (bit specified
(Mj(DP)) = 0 j = 0 to 3	_	Skips the next instruction when the conte of M(DP) is "0."
(A) = (M(DP))	-	Skips the next instruction when the conte
(A) = n	-	Skips the next instruction when the conte

# MITSUBISHI MICROCOMPUTERS 4570 Group

#### SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

#### Detailed description

A. Stores the result in register A. The contents of carry flag CY

lag CY to register A. Stores the result in register A and carry flag

I to register A. unchanged. s no overflow as the result of operation.

he contents of register A and the contents of M(DP), and stores

he contents of register A and the contents of M(DP), and stores

ntents of carry flag CY is "0."

ter A's contents in register A.

uding the contents of carry flag CY to the right by 1 bit.

by the value j in the immediate field) of M(DP) to "1."

d by the value j in the immediate field) of M(DP) to "0."

ntents of bit j (bit specified by the value j in the immediate field)

ntents of register A is equal to the contents of M(DP).

ntents of register A is equal to the value n in the immediate field.

#### SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

## MACHINE INSTRUCTIONS (CONTINUED)

		<u>10</u>				0				_							
Parameter						In	struc	ction	cod	le					er of Is	er of es	
Type of instructions	Mnemonic	D9	D8	D7	D6	D5	D4	Dз	D2	D1	Do	Hex no	adec otati		Number of words	Number of cycles	Function
	Ва	0	1	1	<b>a</b> 6	<b>a</b> 5	<b>a</b> 4	<b>a</b> 3	<b>a</b> 2	aı	<b>a</b> 0	1	8 +a	а	1	1	(PC∟) ← a6–a0
ation	BL p, a	0	0	1	1	1	p4	рз	p2	рı	<b>p</b> o	0	Е +р	р	2	2	(РСн) ← р (РС∟) ← а6-а0
Branch operation		1	0	<b>p</b> 5	<b>a</b> 6	<b>a</b> 5	<b>a</b> 4	<b>a</b> 3	<b>a</b> 2	aı	<b>a</b> 0	2	р +а	а			(Note)
Ъ	BLA p	0	0	0	0	0	1	0	0	0	0	0	1	0	2	2	(РСн) ← р
		1	0	p5	<b>p</b> 4	0	0	рз	p2	рı	<b>p</b> 0	2	р	р			$(PCL) \leftarrow (DR_2 - DR_0, A_3 - A_0)$ (Note)
	BM a	0	1	0	<b>a</b> 6	<b>a</b> 5	a4	аз	<b>a</b> 2	<b>a</b> 1	<b>a</b> 0	1	а	а	1	1	$(SP) \leftarrow (SP) + 1$ $(SK(SP)) \leftarrow (PC)$ $(PCH) \leftarrow 2$ $(PCL) \leftarrow a_{6}-a_{0}$
beration	BML p, a	0	0	1	1	0	p4	рз	p2	p1	<b>p</b> o	0	С +р	р	2	2	$(SP) \leftarrow (SP) + 1$ $(SK(SP)) \leftarrow (PC)$ $(PCH) \leftarrow p$
Subroutine operation		1	0	p₅	<b>a</b> 6	<b>a</b> 5	a4	<b>a</b> 3	<b>a</b> 2	aı	<b>a</b> 0	2	p +a	а			$(PCL) \leftarrow a_{6}-a_{0}$ (Note)
Sut	BMLA p	0	0	0	0	1	1	0	0	0	0	0	3	0	2	2	(SP) ← (SP) + 1 (SK(SP)) ← (PC)
		1	0	p5	<b>p</b> 4	0	0	рз	p2	рı	<b>p</b> o	2	р	р			$(PCH) \leftarrow p$ $(PCL) \leftarrow (DR_2-DR_0, A_3-A_0)$ (Note)
tion	RTI	0	0	0	1	0	0	0	1	1	0	0	4	6	1	1	$(PC) \leftarrow (SK(SP))$ $(SP) \leftarrow (SP) - 1$
Return operation	RT	0	0	0	1	0	0	0	1	0	0	0	4	4	1	2	$(PC) \leftarrow (SK(SP))$ $(SP) \leftarrow (SP) - 1$
2	RTS	0	0	0	1	0	0	0	1	0	1	0	4	5	1	2	$(PC) \leftarrow (SK(SP))$

Skip condition	Carry flag CY	
_	_	Branch within a page : Branches to add
-	-	Branch out of a page : Branches to add
_	_	Branch out of a page : Branches to add A in page p.
_	_	Call the subroutine in page 2 : Calls the
_	_	Call the subroutine : Calls the subroutin
_	_	Call the subroutine : Calls the subroutine D and A in page p.
		Returns from interrupt service routine to Returns each value of data pointer (X, Y description of the LA/LXY instruction, re
_	_	Returns from subroutine to the routine of
Skip unconditionally	_	Returns from subroutine to the routine cal

Note: p is 0 to 31 for M34570M4 and p is 0 to 63 for M34570E8 and M34570M8.

p is 0 to 127 for M34570ED and M34570MD, and p6 is specified with the SBK and RBK instructions.

# MITSUBISHI MICROCOMPUTERS 4570 Group

#### SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

Detailed description

dress a in the identical page.

dress a in page p.

dress (DR2 DR1 DR0 A3 A2 A1 A0)2 specified by registers D and

ne subroutine at address a in page 2.

ine at address a in page p.

ne at address (DR2 DR1 DR0 A3 A2 A1 A0)2 specified by registers o.

to main routine.

Y, Z), carry flag, skip status, NOP mode status by the continuous register A and register B to the states just before interrupt.

called the subroutine.

alled the subroutine, and skips the next instruction unconditionally.

#### SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

## **MACHINE INSTRUCTIONS (CONTINUED)**

Parameter									coc	-					er of Is	r of s	
Type of instructions	Mnemonic	D9	D8	D7	D6	D5	D4	Dз	D2	D1	Do		kade otat	cimal ion	Number of words	Number of cycles	Function
	DI	0	0	0	0	0	0	0	1	0	0	0	0	4	1	1	(INTE) ← 0
	EI	0	0	0	0	0	0	0	1	0	1	0	0	5	1	1	(INTE) ← 1
	SNZ0	0	0	0	0	1	1	1	0	0	0	0	3	8	1	1	(EXF0) = 1 ? After skipping the next instruction, $(EXF0) \leftarrow 0$
ation	SNZI0	0	0	0	0	1	1	1	0	1	0	0	3	A	1	1	l12 = 1 : (INT) = "H" ?
Interrupt operation																	l12 = 0 : (INT) = "L" ?
-	TAV1	0	0	0	1	0	1	0	1	0	0	0	5	4	1	1	(A) ← (V1)
	TV1A	0	0	0	0	1	1	1	1	1	1	0	3	F	1	1	(V1) ← (A)
	TAV2	0	0	0	1	0	1	0	1	0	1	0	5	5	1	1	(A) ← (V2)
	TV2A	0	0	0	0	1	1	1	1	1	0	0	3	Е	1	1	(V2) ← (A)
	TAI1	1	0	0	1	0	1	0	0	1	1	2	5	3	1	1	$(A) \leftarrow (I1)$
	TI1A	1	0	0	0	0	1	0	1	1	1	2	1	7	1	1	(I1) ← (A)
	TAW1	1	0	0	1	0	0	1	0	1	1	2	4	В	1	1	$(A) \leftarrow (W1)$
	TW1A	1	0	0	0	0	0	1	1	1	0	2	0	Е	1	1	(W1) ← (A)
	TAW2	1	0	0	1	0	0	1	1	0	0	2	4	С	1	1	$(A) \leftarrow (W2)$
ç	TW2A	1	0	0	0	0	0	1	1	1	1	2	0	F	1	1	(W2) ← (A)
peratio	TAW3	1	0	0	1	0	0	1	1	0	1	2	4	D	1	1	$(A) \leftarrow (W3)$
Timer operation	ТѠЗА	1	0	0	0	0	1	0	0	0	0	2	1	0	1	1	(W3) ← (A)
	TAW5	1	0	0	1	0	0	1	1	1	1	2	4	F	1	1	(A) ← (0, 0, W51, W50)
	TW5A	1	0	0	0	0	1	0	0	1	0	2	1	2	1	1	(W51, W50) ← (A1, A0)

Skip condition	Carry flag CY	
-	_	Clears the interrupt enable flag INTE to "
-	-	Sets the interrupt enable flag INTE to "1,"
(EXF0) = 1	-	Skips the next instruction when the conte After skipping, clears the EXF0 flag to "0.
(INT) = "H" However, I1 ₂ = 1	-	When bit 2 (I12) of register I1 is "1" : Skip
(INT) = "L" However, I1 ₂ = 0	-	When bit 2 (I12) of register I1 is "0" : Skip
-	-	Transfers the contents of interrupt contro
-	-	Transfers the contents of register A to int
-	-	Transfers the contents of interrupt contro
-	-	Transfers the contents of register A to int
-	-	Transfers the contents of interrupt contro
-	-	Transfers the contents of register A to int
-	-	Transfers the contents of timer control re-
-	-	Transfers the contents of register A to tim
-	-	Transfers the contents of timer control re
-	-	Transfers the contents of register A to tim
-	-	Transfers the contents of timer control re
-	-	Transfers the contents of register A to tim
-	_	Transfers the contents of timer count valu contents of the high-order 2 bits of registe
_	_	Transfers the contents of the low-order 2

# MITSUBISHI MICROCOMPUTERS 4570 Group

### SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

# Detailed description "0," and disables the interrupt. ," and enables the interrupt. tents of EXF0 flag is "1." 0." ips the next instruction when the level of INT pin is "H." ips the next instruction when the level of INT pin is "L." ol register V1 to register A. nterrupt control register V1. ol register V2 to register A. nterrupt control register V2. ol register I1 to register A. nterrupt control register I1. register W1 to register A. imer control register W1. egister W2 to register A. imer control register W2. register W3 to register A. imer control register W3. alue store register W5 to the low-order 2 bits of register A. The ster A is set to "0." 2 bits of register A to timer count value store register W5.

#### SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

### MACHINE INSTRUCTIONS (CONTINUED)

Parameter								ction						er of ds	er of es			
Type of instructions	Mnemonic	D9	D8	D7	D6	D5	D4	Dз	D2	D1	D ₀	Hexad nota		Number of words	Number of cycles	Function		
	TAB1	1	0	0	1	1	1	0	0	0	0	2 7	0	1	1	$(W5) \leftarrow (T19, T18)$ (B) $\leftarrow (T17-T14)$ (A) $\leftarrow (T13-T10)$		
	T1AB	1	0	0	0	1	1	0	0	0	0	23	0	1		At timer 1 stop (W10=0), (R19, R18) $\leftarrow$ (W5) (T19, T18) $\leftarrow$ (W5) (R17-R14) $\leftarrow$ (B) (T17-T14) $\leftarrow$ (B) (R13-R10) $\leftarrow$ (A) (T13-T10) $\leftarrow$ (A) At timer 1 operating (W10=1), (R19, R18) $\leftarrow$ (W5) (R17-R14) $\leftarrow$ (B) (R13-R10) $\leftarrow$ (A)		
	TAB2	1	0	0	1	1	1	0	0	0	1	2 7	1	1	1	(B) ← (T27–T24) (A) ← (T23–T20)		
Timer operation	T2AB	1	0	0	0	1	1	0	0	0	1	23	1	1	1	$(R27-R24) \leftarrow (B)$ $(T27-T24) \leftarrow (B)$ $(R23-R20) \leftarrow (A)$ $(T23-T20) \leftarrow (A)$		
Time	TR2AB	1	0	0	0	1	1	1	0	1	0	23	A	1	1	(R27–R24) ← (B) (R23–R20) ← (A)		
	TAB3	1	0	0	1	1	1	0	0	1	0	2 7	2	1	1	(B) ← (T37–T34) (A) ← (T33–T30)		
	ТЗАВ	1	0	0	0	1	1	0	0	1	0	2 3	2	1		$(R3L7-R3L4) \leftarrow (B)$ $(T37-T34) \leftarrow (B)$ $(R3L3-R3L0) \leftarrow (A)$ $(T33-T30) \leftarrow (A)$		
	ТЗНАВ	1	0	0	0	1	1	1	1	0	1	2 3	D	1	1	(R3H7–R3H4) ← (B) (R3H3–R3H0) ← (A)		

ndition Carry flag CY	Skip condition
<ul> <li>Transfers the contents of the high-order 2 the low-order 8 bits of timer 1 to registers</li> </ul>	-
<ul> <li>When stopping (W10=0), transfers the cont timer 1 and of the timer 1 reload register, a of the low-order 8 bits of timer 1 and of the When operating (W10=1), transfers the con of the timer 1 reload register, and transfers order 8 bits of the timer 1 reload register.</li> </ul>	_
<ul> <li>Transfers the contents of timer 2 to register</li> </ul>	-
<ul> <li>Transfers the contents of registers A and I</li> </ul>	-
<ul> <li>Transfers the contents of registers A and I</li> </ul>	_
<ul> <li>Transfers the contents of timer 3 to registe</li> </ul>	-
<ul> <li>Transfers the contents of registers A and I</li> </ul>	-
<ul> <li>Transfers the contents of registers A and I</li> </ul>	-

# MITSUBISHI MICROCOMPUTERS 4570 Group

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

Detailed description

er 2 bits of timer 1 to register W5, and transfers the contents of ers A and B.

contents of register W5 to the contents of the high-order 2 bits of r, and transfers the contents of registers A and B to the contents the timer 1 reload register.

contents of register W5 to the contents of the high-order 2 bits ers the contents of registers A and B to the contents of the lower.

isters A and B.

nd B to timer 2 and timer 2 reload register.

nd B to timer 2 reload register.

isters A and B.

nd B to timer 3 and timer 3 reload register R3L.

nd B to timer 3 reload register R3H.

### SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

## MACHINE INSTRUCTIONS (CONTINUED)

Parameter					-		struc								er of s	r of s	
Type of instructions	Mnemonic	D9	D8	D7	D6	D5	D4	D3	D2	D1	Do		ade otat	cimal ion	Number of words	Number of cycles	Function
	SNZT1	1	0	1	0	0	0	0	0	0	0	2	8	0	1	1	(T1F) = 1 ? After skipping the next instruction (T1F) $\leftarrow$ 0
Timer operation	SNZT2	1	0	1	0	0	0	0	0	0	1	2	8	1	1	1	(T2F) = 1 ? After skipping the next instruction (T2F) $\leftarrow$ 0
T	SNZT3	1	0	1	0	0	0	0	0	1	0	2	8	2	1	1	(T3F) = 1 ? After skipping the next instruction (T3F) $\leftarrow$ 0
	IAP0	1	0	0	1	1	0	0	0	0	0	2	6	0	1	1	$(A) \gets (P0)$
	OP0A	1	0	0	0	1	0	0	0	0	0	2	2	0	1	1	(P0) ← (A)
	IAP1	1	0	0	1	1	0	0	0	0	1	2	6	1	1	1	$(A) \gets (P1)$
	OP1A	1	0	0	0	1	0	0	0	0	1	2	2	1	1	1	$(P1) \leftarrow (A)$
	IAP2	1	0	0	1	1	0	0	0	1	0	2	6	2	1	1	$(A_1, A_0) \leftarrow (P2_1, P2_0)$ $(A_3, A_2) \leftarrow 0$
	IAP3	1	0	0	1	1	0	0	0	1	1	2	6	3	1	1	$(A) \gets (P3)$
ation	ОРЗА	1	0	0	0	1	0	0	0	1	1	2	2	3	1	1	$(P3) \leftarrow (A)$
ut oper	IAP4	1	0	0	1	1	0	0	1	0	0	2	6	4	1	1	(A) ← (P4)
Input/Output operation	CLD	0	0	0	0	0	1	0	0	0	1	0	1	1	1	1	(D) ← 1
ndul	RD	0	0	0	0	0	1	0	1	0	0	0	1	4	1	1	(D(Y)) ← 0 (Y) = 0 to 9
	SD	0	0	0	0	0	1	0	1	0	1	0	1	5	1	1	(D(Y)) ← 1 (Y) = 0 to 9
	TK0A	1	0	0	0	0	1	1	0	1	1	2	1	В	1	1	$(K0) \gets (A)$
	ТАК0	1	0	0	1	0	1	0	1	1	0	2	5	6	1	1	(A) ← (K0)
	TPU0A	1	0	0	0	1	0	1	1	0	1	2	2	D	1	1	$(PU0) \leftarrow (A)$
	TAPU0	1	0	0	1	0	1	0	1	1	1	2	5	7	1	1	(A) ← (PU0)

Skip conditi	uo Carry flag CY	
(T1F) = 1	-	Skips the next instruction when the cont After skipping, clears T1F flag.
(T2F) = 1	-	Skips the next instruction when the cont After skipping, clears T2F flag.
(T3F) = 1	-	Skips the next instruction when the contend of the skipping, clears T3F flag.
_	-	Transfers the input of port P0 to register
_	-	Outputs the contents of register A to por
_	-	Transfers the input of port P1 to register
-	-	Outputs the contents of register A to por
-	-	Transfers the input of port P2 to register
_	-	Transfers the input of port P3 to register
_	-	Outputs the contents of register A to por
-	-	Transfers the input of port P4 to register
-	-	Sets port D to "1."
-	-	Clears a bit of port D specified by registe
-	-	Sets a bit of port D specified by register
_	-	Transfers the contents of register A to ke
-	-	Transfers the contents of key-on wakeup
-	-	Transfers the contents of register A to p
-	-	Transfers the contents of pull-up control

# MITSUBISHI MICROCOMPUTERS 4570 Group

## SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

Detailed description
ntents of T1F flag is "1."
ntents of T2F flag is "1."
ntents of T3F flag is "1."
er A.
ort P0.
er A.
ort P1.
er A.
er A.
ort P3.
er A.
ster Y to "0."
er Y to "1."
key-on wakeup control register K0.
up control register K0 to register A.
pull-up control register PU0.
ol register PU0 to register A.

#### SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

Parameter			Instruction code													er of es	
Type of instructions	Mnemonic	D9	D8	D7	D6	D5 D4 D3 D2 D1 [		Do		ade otati	cimal ion	Number of words	Number of cycles	Function			
Carrier generating circuit operation	TC2A	1	0	1	0	1	0	1	0	0	1	2	A	9	1	1	(C21, C20) ← (A1, A0)
	NOP	0	0	0	0	0	0	0	0	0	0	0	0	0	1	1	(PC) ← (PC) + 1
	POF	0	0	0	0	0	0	0	0	1	0	0	0	2	1	1	Transition to RAM back-up mode
	EPOF	0	0	0	1	0	1	1	0	1	1	0	5	В	1	1	POF instruction valid
Other operation	SNZP	0	0	0	0	0	0	0	0	1	1	0	0	3	1	1	(P) = 1 ?
Other	WRST	1	0	1	0	1	0	0	0	0	0	2	A	0	1	1	$(WDF1) \leftarrow 0, (WEF) \leftarrow 1$
	TABSI	1	0	0	1	1	1	1	0	0	0	2	7	8	1	1	
	TSIAB	1	0	0	0	1	1	1	0	0	0	2	3	8	1	1	$(SI_7-SI_4) \leftarrow (B)$ $(SI_3-SI_0) \leftarrow (A)$
	TAMR	1	0	0	1	0	1	0	0	1	0	2	5	2	1	1	$(A) \leftarrow (MR_3 – MR_0)$
	TMRA	1	0	0	0	0	1	0	1	1	0	2	1	6	1	1	(MR3–MR₀) ← (A)
	SBK	0	0	0	1	0	0	0	0	0	1	0	4	1	1	1	When executing the TABP p instruction, $p_6 \leftarrow 1$
	RBK	0	0	0	1	0	0	0	0	0	0	0	4	0	1	1	When executing the TABP p instruction, $p_{6} \leftarrow 0$

Skip condition	Carry flag CY	
-	-	Transfers the contents of register A to ca
-	-	No operation
_	-	Puts the system in RAM back-up mode EPOF instruction.
_	-	Validates the POF instruction which is e instruction.
(P) = 1	-	Skips the next instruction when P flag is After skipping, P flag remains unchange
-	-	Operates the watchdog timer and initialized
-	-	Transfers the contents of general-purpos
-	-	Transfers the contents of registers A and
-	-	Transfers the contents of clock control re
-	-	Transfers the contents of register A to cl
_	-	Data area which is referred when execut This setting is valid only for the TABP p
_	-	Data area which is referred when execut This setting is valid only for the TABP p If the SBK instruction is not executed, pe

# MITSUBISHI MICROCOMPUTERS 4570 Group

### SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

#### Detailed description

carrier wave output control register C2.

de state by executing the POF instruction after executing the

s executed after the EPOF instruction by executing the EPOF

is "1." jed.

alizes the watchdog timer flag (WDF1).

ose register SI to registers A and B.

nd B to general-purpose register SI.

register MR to register A.

clock control register MR.

uting the TABP p instruction is set to pages 64 to 127.

cuting the TABP p instruction is set to pages 0 to 63. p instruction. p6 when executing the TABP p instruction is "0."

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

## **CONTROL REGISTERS**

	Interrupt control register V1		reset : 00002	RAM back-up : 00002	R/W
V13	Timer 2 interrupt enable bit	0	Interrupt disabled (	SNZT2 instruction is valid)	
V15		1	Interrupt enabled (	SNZT2 instruction is invalid)	
\/10	V12 Timer 1 interrupt enable bit	0	Interrupt disabled (SNZT1 instruction is valid)		
V 12		1	Interrupt enabled (SNZT1 instruction is invalid)		
V11	Not used	0	This bit has as function, but as allowith is such lad		
VII		1	This bit has no function, but read/write is enabled.		
V10	External 0 interrupt enable bit	0	Interrupt disabled (	SNZ0 instruction is valid)	
VIO		1	Interrupt enabled (	SNZ0 instruction is invalid)	

Interrupt control register V2		at reset : 00002		at RAM back-up : 00002	R/W	
V23	V23 Not used	0	This bit has no function, but read/write is enabled.			
-		1	This bit has no function, but read/while is enabled.			
V22	V22 Not used	0	This bit has no function, but read/write is enabled.			
V 22		1				
V21	Not used	0	This hit has no function, but read/units is suchlad			
VZI		1	This bit has no function, but read/write is enabled.			
V20	Timer 3 interrupt enable bit	0	Interrupt disabled (	SNZT3 instruction is valid)		
v 20	Timer 3 interrupt enable bit	1	Interrupt enabled (SNZT3 instruction is invalid)			

	Interrupt control register I1		reset : 00002	at RAM back-up : state retained	R/W
113	Not used	0	This bit has no function, but read/write is enabled.		
113		1			
		0	Falling waveform ("L" level of INT pin is recognized with the SN		he SNZI0
14-	Interrupt valid waveform for INT pin /return		instruction)/"L" level		
I12	level selection bit (Note 2)	1	Rising waveform ("H" level of INT pin is recognized with the SNZIO		
			instruction)/"H" leve	el	
111	Not used	0	This hit has no fun	ation, but road/write is enabled	
111	Not used	1	This bit has no function, but read/write is enabled.		
14.	Netwood	0			
<b>11</b> 0	Not used	1	This bit has no function, but read/write is enabled.		

Notes 1: "R" represents read enabled, and "W" represents write enabled.

2: Depending on the input state of P21/INT pin, the external interrupt request flag EXF0 may be set to "1" when the contents of I12 is changed. Accordingly, set a value to bit 2 of register I1 and execute the SNZ0 instruction to clear the EXF0 flag after executing at least one instruction.

#### SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

## **CONTROL REGISTERS (CONTINUED)**

	Timer control register W1		reset : 00002	at RAM back-up : 00002	R/W		
\\/1e	Prescaler control bit	0	Stop (prescaler stat	Stop (prescaler state initialized)			
VV 13	W13 Prescaler control bit	1	Operating				
10/4 -	M/4 December dividing ratio coloritien bit	0	Instruction clock divided by 4				
W12	Prescaler dividing ratio selection bit	1	Instruction clock divided by 8				
	Timer 1 count source selection bit	0	Prescaler output (ORCLK)				
W11	Timer T count source selection bit	1	Carrier output (CARRY)				
10/4 -	Timor 1 control bit	0	Stop (state retained)				
W1o	Timer 1 control bit	1	Operating				

	Timer control register W2		at reset : 00002		at RAM back-up : state retained R	/W
W23	Timer 2 control bit	(	)	Stop (state retained	)	
VVZ3		-	1	Operating		
W22	Port D ₉ /Tout pin function selection bit	0		Port D ₉		
V V Z Z			1	Τουτ pin		
	Timer 2 count source selection bits	W21	W20		Count source	
W21		0	0	Prescaler output (C	DRCLK)	
		0	1	Timer 1 underflow signal		
W20		1	0	Instruction clock		
		1	1	16-bit timer underfle	ow signal	

	Timer control register W3		at reset : 00002		at RAM back-up : state retained R/W	
W33	Timer 3 control bit	(	)	Stop (state retained)		
VV 33		1		Operating		
W32	Not used	0 This bit ha		This bit has no func	his bit has no function, but read/write is enabled.	
					Courteouros	
		W31	W30		Count source	
W31		0	0	Timer 2 underflow s	signal	
	Timer 3 count source selection bits	0	1	Prescaler output (ORCLK)		
W30		1	0	f(XIN) or f(XIN)/2		
		1	1	Not available		

at reset : 002	at RAM back-up : state retained	R/W					
2-bit register. The contents of the high-order 2 bits (bits 9 and 8) of the 10-bit ROM pattern at address (D2D1D0A3A2A1A0) in page							
ster W5 with the TABP p instr	uction.						
In addition, data can be transferred between the low-order 2 bits of register A and this register W5 with the TW5A or TAW5							
instruction. Data can be read/written to/from the high-order 2 bits of timer 1 with the T1AB or TAB1 instruction.							
	its 9 and 8) of the 10-bit ROM ster W5 with the TABP p instr ow-order 2 bits of register A a	its 9 and 8) of the 10-bit ROM pattern at address (D2D1D0A3A2A1A ster W5 with the TABP p instruction. w-order 2 bits of register A and this register W5 with the TW5A					

Note: "R" represents read enabled, and "W" represents write enabled.

#### SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

## **CONTROL REGISTERS (CONTINUED)**

Ca	Carrier wave output control register C2		at reset : 002	at RAM back-up : 002	W
C21	C21 Port CARR output control bit	0	Port CARR "L" level output		
021		1	Port CARR "H" level output		
C20	Carrier wave output auto-control bit	0	Auto-control output by timer 1 is invalid		
020	Carrier wave output auto-control bit	1	Auto-control output	by timer 1 is valid	

	Key-on wakeup control register K0		reset : 00002	at RAM back-up : state retained	R/W
K03	Port P43 key-on wakeup	0	Key-on wakeup not	used	
KU3	control bit	1 Key-on wakeup used		d	
K02	Port P42 key-on wakeup	0	Key-on wakeup not used		
K02	control bit		Key-on wakeup used		
K01	Port P41 key-on wakeup	0	Key-on wakeup not used		
KU1	control bit	1	Key-on wakeup used		
KO	Port P40 key-on wakeup		Key-on wakeup not	used	
K00	control bit	1	Key-on wakeup use	d	

Pull-up control register PU0		at reset : 00002		at RAM back-up : state retained	R/W	
PU03	Port P43 pull-up transistor	0	Pull-up transistor O	FF		
P003	control bit	1	Pull-up transistor O	N		
DUO	Port P42 pull-up transistor	0	0 Pull-up transistor OFF			
PU02	control bit	1	Pull-up transistor ON			
PU01	Port P41 pull-up transistor	0	Pull-up transistor OFF			
P001	control bit	1	Pull-up transistor ON			
PU00	Port P40 and P01 pull-up transistor	0	Pull-up transistor O	FF		
P000	control bit	1	Pull-up transistor ON			

	Clock control register MR		reset : 10002	at RAM back-up : state retained	R/W	
MR3			f(Xin)			
IVIR3	System clock selection bit	1	f(XIN)/4			
MR ₂	MR2 Not used	0	This hit has no function, but read/write is anabled			
IVIT\2		1	This bit has no function, but read/write is enabled.			
MR1	Not used	0	This bit has no function, but read/write is enabled.			
		1				
MRo	Netwood	0	This bit has no function, but read/write is enabled.			
IVIR0	Not used	1				

8-bit general purpose register PU0	at reset : 0016	at RAM back-up : state retained	R/W
8-bit general purpose register.			

8-bit data can be transferred between this register PU0 and registers A and B with the TSIAB instruction and TABSI instruction.

Note: "R" represents read enabled, and "W" represents write enabled.

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

## **ABSOLUTE MAXIMUM RATINGS**

Symbol	Parameter	Conditions	Ratings	Unit
Vdd	Supply voltage		-0.3 to 7.0	V
	Input voltage		0.2 to 1/22 : 0.2	V
Vi	P0, P1, P2, P3, P4, RESET, XIN, VDCE		-0.3 to VDD+0.3	V
Vo	Output voltage P0, P1, P3, D	Output transistors in cut-off state	-0.3 to VDD+0.3	V
Vo	Output voltage CARR, Xout		-0.3 to VDD+0.3	V
Pd	Power dissipation		300	mW
Topr	Operating temperature range		-20 to 70	°C
Tstg	Storage temperature range		-40 to 125	°C

#### **RECOMMENDED OPERATING CONDITIONS1**

(Mask ROM version: Ta = -20 °C to 70 °C, VDD = 2.0 V to 5.5 V, unless otherwise noted)

(One Time PROM version:Ta = -20 °C to 70 °C, VDD = 2.5 V to 5.5 V, unless otherwise noted)

Symbol	Parameter		Conditions	Limits				
Symbol			Conditions	Min.	Тур.	Max.	Unit	
		Mask ROM version System clock =f(XIN)/4	$f(X_{IN}) \le 4.2 \text{ MHz}$ Ceramic resonator	2.0		5.5		
		Mask ROM version System clock	$f(X_{IN}) \le 2.0 \text{ MHz}$ Ceramic resonator	4.5		5.5		
Vdd	Supply voltage	=f(XiN)	$f(X_{IN}) \le 1.0 \text{ MHz}$ Ceramic resonator	2.0		5.5	v	
		One Time PROM version System clock =f(XIN)/4	$f(X_{IN}) \le 4.2 \text{ MHz}$ Ceramic resonator	2.5		5.5		
		One Time PROM version System clock	$f(X_{IN}) \le 2.0 \text{ MHz}$ Ceramic resonator	4.5		5.5		
		=f(XIN)	f(Xıℕ) ≤ 1.0 MHz Ceramic resonator	2.5		5.5		
Vram	RAM back-up	Mask ROM version	PAM book up	1.8		5.5	V	
	voltage	One Time PROM version	RAM back-up	2.0		5.5	V	
Vss	Supply voltage				0		V	
		Mask ROM version System clock =f(XIN)/4	VDD=2.0 V to 5.5V			4.2		
	Oscillation	Mask ROM version	VDD=4.5 V to 5.5V			2.0	_	
f(Xın)	frequency	=f(X _{IN} )	Vdd=2.0 V to 5.5V			1.0	MHz	
	(at ceramic resonance)	One Time PROM version System clock =f(XIN)/4	VDD=2.5 V to 5.5V			4.2	_	
		One Time PROM version System clock	VDD=4.5 V to 5.5V			2.0	_	
		=f(XIN)	VDD=2.5 V to 5.5V			1.0	1	

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

### **RECOMMENDED OPERATING CONDITIONS 2**

(Mask ROM version:Ta = -20 °C to 70 °C, V_{DD} = 2.0 V to 5.5 V, unless otherwise noted) (One Time PROM version:Ta = -20 °C to 70 °C, V_{DD} = 2.5 V to 5.5 V, unless otherwise noted)

Cumple of		Conditions	Limits				
Symbol	Parameter	Conditions	Min.	Тур.	Max.	Unit	
Viн	"H" level input voltage P0, P1, P2, P3, P4, VDCE		0.8Vdd		Vdd	V	
Vін	"H" level input voltage XIN		0.7Vdd		Vdd	V	
Vін	"H" level input voltage RESET		0.85Vdd		Vdd	V	
Vін	"H" level input voltage INT		0.8Vdd		Vdd	V	
VIL	"L" level input voltage P0, P1, P2, P3, P4, VDCE		0		0.3Vdd	V	
Vil	"L" level input voltage XIN		0		0.3Vdd	V	
Vil	"L" level input voltage RESET		0		0.3Vdd	V	
Vil	"L" level input voltage INT		0		0.2Vdd	V	
loL(peak)	"L" level peak output current	VDD=5.0 V			10	mA	
	P0, P1, D0–D9, CARR	VDD=3.0 V			4	- mA	
loL(peak)	"L" level peak output current P3	VDD=5.0 V			30	mA	
		VDD=3.0 V			24	- mA	
lo∟(avg)	"L" level average output current	VDD=5.0 V			5	- mA	
	P0, P1, D0–D9, CARR (Note)	VDD=3.0 V			2	- mA	
lo∟(avg)	"L" level average output current P3	VDD=5.0 V			15	mA	
	(Note)	VDD=3.0 V			12	1	
loн <b>(peak)</b>	"H" level peak output current	VDD=5.0 V			-30	mA	
	CARR	VDD=3.0 V			-15		
loн(avg)	"H" level average output current	VDD=5.0 V			-15	mA	
	CARR (Note)	VDD=3.0 V			-7	IIIA	
Σ Ιοι	"L" total current P0, P1, P3				30	mA	
Σ Ιοι	"L" total current D				20	mA	
Γρον	Power reset circuit valid power rising	Mask ROM version					
	time	VDD = 0 to 2.0 V			100		
		One Time PROM version			100	μs	
		VDD = 0 to 2.5 V					

Note: The average output current is the average current value at the 100 ms interval.

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

## **ELECTRICAL CHARACTERISTICS**

(Mask ROM version:Ta = -20 °C to 70 °C, V_{DD} = 2.0 V to 5.5 V, unless otherwise noted) (One Time PROM version:Ta = -20 °C to 70 °C, V_{DD} = 2.5 V to 5.5 V, unless otherwise noted)

Symbol	Parameter		Test conditions			Unit		
-		Parameter			Min.	Тур.	Max.	
Vol	"L" level output voltage		lo∟ = 5 mA	Vdd = 5.0 V			0.9	v
VOL	P0, P1, D0–D9,	CARR, RESET	IoL = 2 mA	VDD = 3.0 V			0.9	1
Vol	"I " lovel output	voltage D2	lo∟ = 15 mA	Vdd = 5.0 V			1.5	v
VOL	"L" level output	vollage P3	IoL = 12 mA	VDD = 3.0 V			1.5	1
Vон	"II" lovel output		Іон = 15 mA	Vdd = 5.0 V	2.4			v
VOH	"H" level output	Vollage CARR	Іон = –7 mA	Vdd = 3.0 V	1.0			] V
Ін	"H" level input RESET, VDCE	current P0, P1, P2, P3, P4,	VI = VDD (Note)				1	μA
lı∟	"L" level input current P2, P3, P4, RESET, VDCE		VI = 0 V (Note)		-1			μA
loz	Output current a	at off-state D0–D9	Vo = Vdd				1	μA
	Supply current		Vdd = 5.0 V, f(Xin) = 4.2 MHz			1.3	2.6	
			System clock = $f(XIN)/4$			1.5	2.0	
			VDD = 5.0 V	f(XIN) = 2 MHz		1.9	3.8	
		at CBLI operating mode	System clock = f(XIN)	f(XIN) = 1 MHz		1.3	2.6	mA
DD			VDD = 3.0 V, f(XIN) = 4.2	2 MHz		0.0	4.0	
			System clock = $f(X_{IN})/4$			0.6	1.2	
			VDD = 3.0 V	f(XIN) = 1 MHz		0.5	1.0	
			System clock = f(XIN)	f(XIN) = 500 kHz		0.4	0.8	
		at RAM back-up mode	f(XIN) = stop, typical val	ue at Ta = 25 °C		0.1	10	μA
		P0, P1, P4	VDD = 5.0 V, VI = 0 V		20	50	125	1.0
Devi	Pull-up resistor	FU, F1, F4	Vdd = 3.0 V, VI = 0 V		40	100	250	kΩ
Rрн	value	RESET	VDD = 5.0 V, VI = 0 V		12	30	70	1.0
		RESEI	$V_{DD} = 3.0 V, V_I = 0 V$		25	60	130	kΩ
		INT	Vdd = 5.0 V			0.5		- v
Vt+ – Vt	Hysteresis		VDD = 3.0 V			0.4		V
v i t v I -		RESET	Vdd = 5.0 V			1.5		
	KESEI		Vdd = 3.0 V			0.6		V

Note: In this case, the pull-up transistor of port P4 is turned off by software.

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

BASIC TIMING DIA	Machine cycle	Mi	Mi+1				
Parameter	State Pin name	 Тз	T1	T ₂	Тз		
Clock	Xın (System clock=f(Xın))						
	XIN (System clock=f(XIN)/4)	חחח					
Port D output	D0-D9	X			X		
Port P0, P1, P3 output	P00-P03 P10-P13 P30-P33	 X			X		
Port P0, P1, P2, P3, P4 input	P00–P03 P10–P13 P20, P21 P30–P33 P40–P43		X		X		
Interrupt input	INT			X			

## **BASIC TIMING DIAGRAM**

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

## **BUILT-IN PROM VERSION**

In addition to the mask ROM version, the 4570 Group has the programmable ROM version software compatible with mask ROM. The One Time PROM version has PROM which can only be written to and not be erased.

The built-in PROM version has functions similar to those of the mask ROM version, but it has a PROM mode that enables writing to built-in PROM.

Table 16 shows the product of built-in PROM version. Figure 35 shows the pin configurations of built-in PROM version. The One Time PROM version has pin-compatibility with the mask ROM version.

## Table 16 Product of built-in PROM version

Product	PROM size (X 10 bits)	RAM size (X 4 bits)	Package	ROM type
M34570E8FP	8192 words	128 words	36P2R-A	One Time PROM
M34570EDFP	16384 words	128 words	36P2R-A	

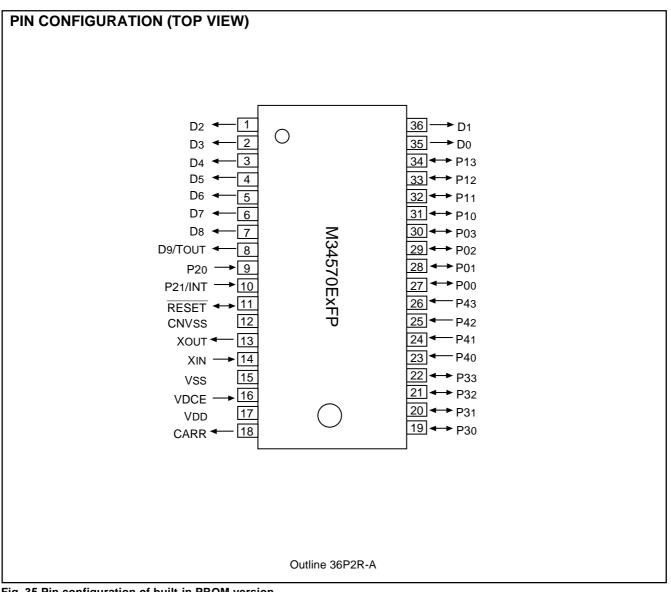


Fig. 35 Pin configuration of built-in PROM version

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

#### (1) PROM mode

The built-in PROM version has a PROM mode in addition to a normal operation mode. The PROM mode is used to write to and read from the built-in PROM.

In the PROM mode, the programming adapter can be used with a general-purpose PROM programmer to write to or read from the built-in PROM as if it were M5M27C256K. Programming adapter is listed in Table 17. Contact addresses at the end of this book for the appropriate PROM programmer.

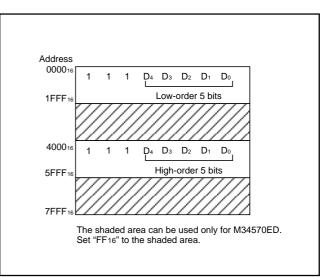
 Writing and reading of built-in PROM Programming voltage is 12.5 V. Write the program in the PROM of the built-in PROM version as shown in Figure 36.

#### (2) Notes on handling

- ① A high-voltage is used for writing. Take care that overvoltage is not applied. Take care especially at turning on the power.
- ② For the One Time PROM version Mitsubishi Electric corp. does not perform PROM writing test and screening in the assembly process and following processes. In order to improve reliability after writing, performing writing and test according to the flow shown in Figure 37 before using is recommended.

#### **Table 17 Programming adapter**

Microcomputer	Programming adapter				
M34570E8FP, M34570EDFP	PCA7425				



#### Fig. 36 PROM memory map

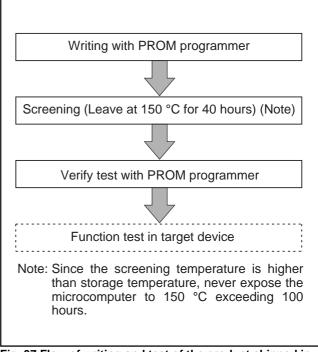


Fig. 37 Flow of writing and test of the product shipped in blank

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

GZ	GZZ-SH55-08B <91A0>					OM number	
	4500 SERIES MASK ROM ORDER CONFIRMATION FORM SINGLE-CHIP MICROCOMPUTER M34570M4-XXXFP MITSUBISHI ELECTRIC					Date: Section head signature	Supervisor signature
	Please fil	ll in all ite	ms marked *.	,	Receipt		
		Company					
		name				Responsible officer	Supervisor
*	Customer		TEL ( )		iture		
		Date issued	Date:	0100	signature		

* 1. Confirmation

Three sets of EPROMs are required for each pattern if this order is performed by EPROMs. One floppy disk is required for each pattern if this order is performed by floppy disk.

Ordering by the EPROMs

Specify the type of EPROMs submitted (check in the approximate box).

If at least two of the three sets of EPROMs submitted contain the identical data, we will produce masks based on this data. We shall assume the responsibility for errors only if the mask ROM data on the products we produce differ from this data. Thus, the customer must be especially careful in verifying the data contained in the EPROMs submitted.

Checksum code for entire EPROM area

(hexadecimal notation)

EPROM	Type:
-------	-------

27C256	27C512			
Low-order 5-bit data High-order 5-bit data 4.00K 4.00K 4.00K 4.00K 4.00K 4.00K 4.00K 4.00K 4.00K 4.00K 4.00K	Low-order 5-bit data High-order 5-bit data 4.00K 4.00K 4.00K 4.00K 4.00K 4.00K 4.00K 4.00K 4.00K 4.00K			

## Set "FF16" in the shaded area.

Set "1112" in the area of low-order and high-order 5-bit data.

**MITSUBISHI MICROCOMPUTERS** 

## 4570 Group

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

GZZ-SH55-08B <91A0>		
022-01133-00B <91702		Mask ROM number
4500 SERIES MASK R SINGLE-CHIP MICR MITSI		
sume the responsibility for er mask file. Thus, extreme car	rors only if the mask ROM data e must be taken to verify the ma ust be-3.5 inch 2HD type and	by the mask file generating utility. We shall as- a on the products we produce differs from this ask file in the submitted floppy disk. DOS/V format. And the number of the mask
File code		(hexadecimal notation)
Mask file name		.MSK (equal or less than eight characters)

#### * 2. Mark Specification

Mark specification must be submitted using the correct form for the type of package being ordered. Fill out the approximate Mark Specification Form (36P2R-A for M34570M4-XXXFP) and attach to the Mask ROM Order Confirmation Form.

* 3. Comments

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

GZ	GZZ-SH55-09B <91A0>					Mask ROM number			
	4500 SERIES MASK ROM ORDER CONFIRMATION FORM SINGLE-CHIP MICROCOMPUTER M34570M8-XXXFP MITSUBISHI ELECTRIC				Receipt	Date: Section head signature	Supervisor signature		
	Please fi	ll in all ite	ms marked *.		Rec				
		Company							
	0	name	/			Responsible officer	Supervisor		
*	Customer		TEL ( )		iture				
		Date issued	Date:	oucraal	signature				

* 1. Confirmation

Three sets of EPROMs are required for each pattern if this order is performed by EPROMs. One floppy disk is required for each pattern if this order is performed by floppy disk.

Ordering by the EPROMs

Specify the type of EPROMs submitted (check in the approximate box).

If at least two of the three sets of EPROMs submitted contain the identical data, we will produce masks based on this data. We shall assume the responsibility for errors only if the mask ROM data on the products we produce differ from this data. Thus, the customer must be especially careful in verifying the data contained in the EPROMs submitted.

Checksum code for entire EPROM area

(hexadecimal notation)

27C256	27C512
Low-order 5-bit data High-order 5-bit data High-order 5-bit data 000016 1FFF16 8.00K 5FFF16 8.00K 5FFF16 7FFF16	Low-order 5-bit data High-order 5-bit data High-order 5-bit data 5FFF ₁₆ 8.00K 8.00K 5FFF ₁₆

## Set "FF16" in the shaded area.

Set "1112" in the area of low-order and high-order 5-bit data.

**MITSUBISHI MICROCOMPUTERS** 

## 4570 Group

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

GZZ-SH55-09B <91A0>		Mask ROM number
SINGLE-CHIP MICR	OM ORDER CONFIRMATIO OCOMPUTER M34570M8-2 UBISHI ELECTRIC	
sume the responsibility for e mask file. Thus, extreme car	rrors only if the mask ROM da e must be taken to verify the r ust be-3.5 inch 2HD type and	by the mask file generating utility. We shall as- ta on the products we produce differs from this nask file in the submitted floppy disk. d DOS/V format. And the number of the mask
File code		(hexadecimal notation)
Mask file name		.MSK (equal or less than eight characters)

#### * 2. Mark Specification

Mark specification must be submitted using the correct form for the type of package being ordered. Fill out the approximate Mark Specification Form (36P2R-A for M34570M8-XXXFP) and attach to the Mask ROM Order Confirmation Form.

#### * 3. Comments

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

GZZ-SH55-10B <91A0>				M	lask R		
4500 SERIES MASK ROM ORDER CONFIRMATION FORM SINGLE-CHIP MICROCOMPUTER M34570MD-XXXFP MITSUBISHI ELECTRIC						Date: Section head signature	Supervisor signature
	Please fil	II in all ite	ms marked *.	,	Rec		
		Company					
٧-		name			υØ	Responsible officer	Supervisor
*	Customer		TEL ( )				
		Date issued	Date:	ouciaa	signature		

* 1. Confirmation

Three sets of EPROMs are required for each pattern if this order is performed by EPROMs. One floppy disk is required for each pattern if this order is performed by floppy disk.

Ordering by the EPROMs

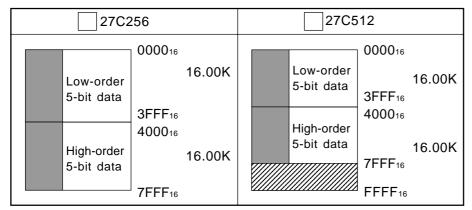
Specify the type of EPROMs submitted (check in the approximate box).

If at least two of the three sets of EPROMs submitted contain the identical data, we will produce masks based on this data. We shall assume the responsibility for errors only if the mask ROM data on the products we produce differ from this data. Thus, the customer must be especially careful in verifying the data contained in the EPROMs submitted.

Checksum code for entire EPROM area

(hexadecimal notation)

EPROM Type:



Set "FF16" in the shaded area.

Set "1112" in the area of low-order and high-order 5-bit data.

**MITSUBISHI MICROCOMPUTERS** 

## 4570 Group

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

C77 CHEE 10D -0100		
GZZ-SH55-10B <91A0>		Mask ROM number
SINGLE-CHIP MICR	OM ORDER CONFIRMATIO OCOMPUTER M34570MD-X JBISHI ELECTRIC	
sume the responsibility for er mask file. Thus, extreme care	rors only if the mask ROM data e must be taken to verify the ma ust be-3.5 inch 2HD type and	by the mask file generating utility. We shall as- a on the products we produce differs from this ask file in the submitted floppy disk. DOS/V format. And the number of the mask
File code		(hexadecimal notation)
Mask file name		.MSK (equal or less than eight characters)

#### * 2. Mark Specification

Mark specification must be submitted using the correct form for the type of package being ordered. Fill out the approximate Mark Specification Form (36P2R-A for M34570MD-XXXFP) and attach to the Mask ROM Order Confirmation Form.

* 3. Comments

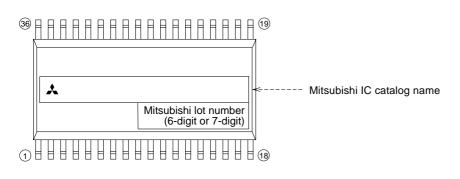
SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

## MARK SPECIFICATION FORM 36P2R-A (36-PIN SHRINK SOP) MARK SPECIFICATION FORM

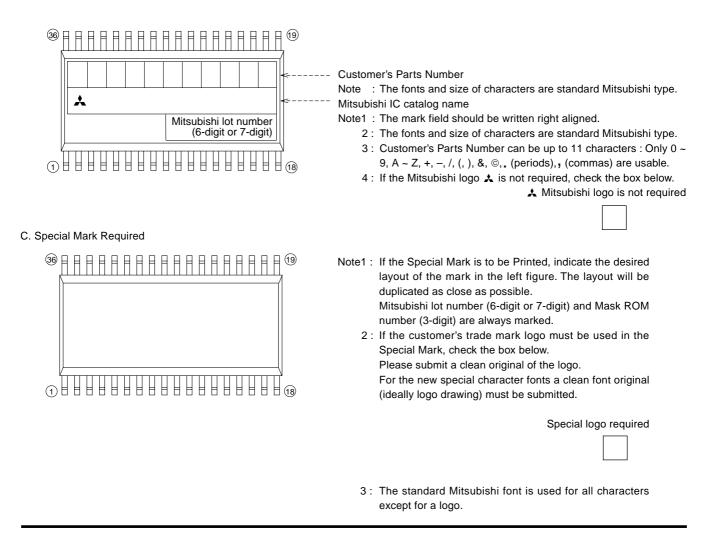
Mitsubishi IC catalog name

Please choose one of the marking types below (A, B, C), and enter the Mitsubishi catalog name and the special mark (if needed).

#### A. Standard Mitsubishi Mark



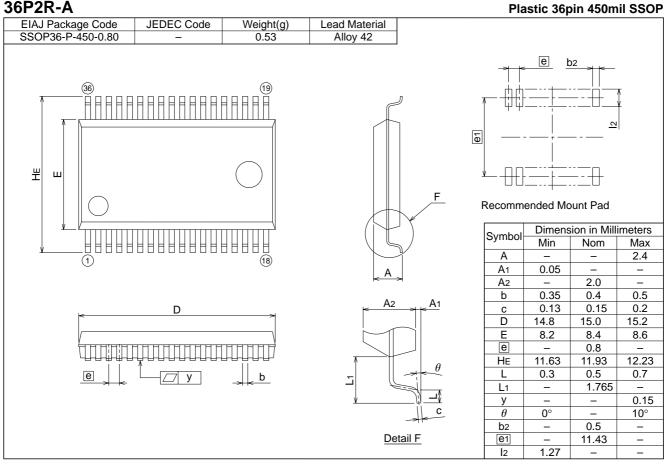
B. Customer's Parts Number + Mitsubishi catalog name



SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

## **PACKAGE OUTLINE**

#### 36P2R-A



**MITSUBISHI MICROCOMPUTERS** 

## 4570 Group

SINGLE-CHIP 4-BIT CMOS MICROCOMPUTER

#### Keep safety first in your circuit designs!

Misubishi Electric Corporation puts the maximum effort into making semiconductor products better and more reliable, but there is always the possibility that trouble may occur with them. Trouble with semiconductors may lead to personal injury, fire or property damage. Remember to give due consideration to safety when making your circuit designs, with appropriate measures such as (i) placement of substitutive, auxiliary circuits, (ii) use of non-flammable material or (iii) prevention against any malfunction or mishap.

#### Notes regarding these materials

- These materials are intended as a reference to assist our customers in the selection of the Mitsubishi semiconductor product best suited to the customer's application; they do not convey any license under any intellectual property rights, or any other rights, belonging to Mitsubishi Electric Corporation or a third party. Mitsubishi Electric Corporation assumes no responsibility for any damage, or infringement of any third-party's rights, originating in the use of any product data, diagrams, charts or circuit application examples contained in these materials.
- All information contained in these materials, including product data, diagrams and charts, represent information on products at the time of publication of these materials, and are subject to change by Mitsubishi
- All miorination contained in these materials, including product data, diagrams and charts, represent miorimation on products at the time of publication or these materials, and are subject to charge by Missobish Electric Corporation without notice due to product improvements or other reasons. It is therefore recommended that customers contact Mitsubishi Electric Corporation or an authorized Mitsubishi Semiconductor product distributor for the latest product information before purchasing a product listed herein. Mitsubishi Electric Corporation semiconductors are not designed or manufactured for use in a device or system that is used under circumstances in which human life is potentially at stake. Please contact Mitsubishi Electric Corporation or an authorized Mitsubishi Semiconductor product distributor when considering the use of a product contained herein for any specific purposes, such as apparatus or systems for transportation, vehicular, medical, aerospace, nuclear, or undersea repeater use. The prior written approval of Mitsubishi Electric Corporation is necessary to reprint or reprint or reprint or in part these materials. If these products or technologies are subject to the Japanese export control restrictions, they must be exported under a license from the Japanese government and cannot be imported into a country other than the construction testeristic.
- Approved destination. Any diversion or reexport contrary to the export control laws and regulations of Japan and/or the country of destination is prohibited. Please contact Mitsubishi Electric Corporation or an authorized Mitsubishi Semiconductor product distributor for further details on these materials or the products contained therein.

© 1999 MITSUBISHI ELECTRIC CORP. New publication, effective April. 1999. Specifications subject to change without notice.

## **REVISION DESCRIPTION LIST**

## 4570 GROUP DATA SHEET

Revision Description	Rev. date
First Edition	971022
Main revision points are described below.	990331
•M34570MD-XXXFP and M34570EDFP (ROM expansion products [size: 16K 5 10 bits] ) added.	
<ul> <li>SBK and RBK instructions added and TABP p instruction function is expanded.</li> </ul>	
(TABP p instruction: When this instruction is executed after executing the SBK instruction, pages 64 to 127 are specified. When this instruction is executed after executing the RBK instruction, pages 0 to 63 are specified. When this instruction is executed after system is released from reset and returned from the RAM back-up mode, pages 0 to 63 are specified.)	
<ul> <li>BL, BML, BLA and BMLA instructions revised. Referred pages are expanded to pages 0 to 127 (p6 can be used for page specification.)</li> </ul>	
	<ul> <li>First Edition</li> <li>Main revision points are described below.</li> <li>•M34570MD-XXXFP and M34570EDFP (ROM expansion products [size: 16K 5 10 bits] ) added.</li> <li>•SBK and RBK instructions added and TABP p instruction function is expanded.</li> <li>(TABP p instruction: When this instruction is executed after executing the SBK instruction, pages 64 to 127 are specified. When this instruction is executed after executing the RBK instruction, pages 0 to 63 are specified. When this instruction is executed after system is released from reset and returned from the RAM back-up mode, pages 0 to 63 are specified.)</li> <li>•BL, BML, BLA and BMLA instructions revised. Referred pages are expanded to pages 0 to 127</li> </ul>