

# LH28F320S5NS-L90

# Flash Memory 32M (4MB × 8 / 2M × 16)

(Model No.: LHF32K10)

Spec No.: EL108051A Issue Date: December 22, 1998

<u>To;</u>	SPEC No.         E L 1 0 8 0 5 1 A           I S S U E :         Dec. 22, 1998
SPECIFICA	TIONS
Product Type <u>32Mbit</u> Flas	sh Memory
LH28F320	S 5 N S – L 9 0
Model No. (LHF32	K10)
**This specifications contains 55 pages i If you have any objections, please conta CUSTONERS ACCEPTANCE	
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#### LHF32K10

# LH28F320S5NS-L90 32-MBIT (4MBx8/2MBx16) Smart 5 Flash MEMORY

- Smart 5 Technology — 5V V<sub>CC</sub>
  - 5V V<sub>PP</sub>
- Common Flash Interface (CFI)
   Universal & Upgradable Interface
- Scalable Command Set (SCS)
- High Speed Write Performance
   32 Bytes x 2 plane Page Buffer
   2µs/Byte Write Transfer Rate
- High Speed Read Performance
   90ns(5V±0.25V), 100ns(5V±0.5V)
- Operating Temperature
   0°C to +70°C
- High-Density Symmetrically-Blocked Architecture
   — Sixty-four 64-Kbyte Erasable Blocks
- Extended Cycling Capability
   100,000 Block Erase Cycles
  - 6.4 Million Block Erase Cycles/Chip
- Automated Write and Erase
   Command User Interface
  - Status Register

- Enhanced Automated Suspend Options
  - Write Suspend to Read
  - Block Erase Suspend to Write
  - Block Erase Suspend to Read
- Enhanced Data Protection Features
  - Absolute Protection with V<sub>PP</sub>=GND
  - Flexible Block Locking
  - Erase/Write Lockout during Power Transitions
- Low Power Management
  - Deep Power-Down Mode
  - Automatic Power Savings Mode
     Decreases I<sub>CC</sub> in Static Mode
- User-Configurable x8 or x16 Operation
- SRAM-Compatible Write Interface
- Industry-Standard Packaging — 56-Lead SSOP
- ETOX<sup>TM\*</sup> V Nonvolatile Flash Technology
- CMOS Process (P-type silicon substrate)
- Not designed or rated as radiation hardened

SHARP's LH28F320S5NS-L90 Flash memory with Smart 5 technology is a high-density, low-cost, nonvolatile, read/write storage solution for a wide range of applications. Its symmetrically-blocked architecture, flexible voltage and extended cycling provide for highly flexible component suitable for resident flash arrays, SIMMs and memory cards. Its enhanced suspend capabilities provide for an ideal solution for code + data storage applications. For secure code storage applications, such as networking, where code is either directly executed out of flash or downloaded to DRAM, the LH28F320S5NS-L90 offers three levels of protection: absolute protection with  $V_{PP}$  at GND, selective hardware block locking, or flexible software block locking. These alternatives give designers ultimate control of their code security needs.

The LH28F320S5NS-L90 is conformed to the flash Scalable Command Set (SCS) and the Common Flash Interface (CFI) specification which enable universal and upgradable interface, enable the highest system/device data transfer rates and minimize device and system-level implementation costs.

The LH28F320S5NS-L90 is manufactured on SHARP's 0.4µm ETOX<sup>TM\*</sup> V process technology. It come in industrystandard package: the 56-Lead SSOP, ideal for board constrained applications.

\*ETOX is a trademark of Intel Corporation.

#### **1 INTRODUCTION**

This datasheet contains LH28F320S5NS-L90 specifications. Section 1 provides a flash memory overview. Sections 2, 3, 4, and 5 describe the memory organization and functionality. Section 6 covers electrical specifications.

#### 1.1 Product Overview

The LH28F320S5NS-L90 is a high-performance 32-Mbit Smart 5 Flash memory organized as 4MBx8/2MBx16. The 4MB of data is arranged in sixty-four 64-Kbyte blocks which are individually erasable, lockable, and unlockable in-system. The memory map is shown in Figure 3.

Smart 5 technology provides a choice of V<sub>CC</sub> and V<sub>PP</sub> combinations, as shown in Table 1, to meet system performance and power expectations. 5V V<sub>CC</sub> provides the highest read performance. V<sub>PP</sub> at 5V eliminates the need for a separate 12V converter, while V<sub>PP</sub>=5V maximizes erase and write performance. In addition to flexible erase and program voltages, the dedicated V<sub>PP</sub> pin gives complete data protection when V<sub>PP</sub>  $\leq$  V<sub>PPI K</sub>.

Table 1. V<sub>CC</sub> and V<sub>PP</sub> Voltage CombinationsOffered by Smart 5 Technology

V <sub>CC</sub> Voltage	V <sub>PP</sub> Voltage
5V	5V

Internal  $V_{CC}$  and  $V_{PP}$  detection Circuitry automatically configures the device for optimized read and write operations.

A Command User Interface (CUI) serves as the interface between the system processor and internal operation of the device. A valid command sequence written to the CUI initiates device automation. An internal Write State Machine (WSM) automatically executes the algorithms and timings necessary for block erase, full chip erase, (multi) word/byte write and block lock-bit configuration operations.

A block erase operation erases one of the device's 64-Kbyte blocks typically within 0.34s (5V  $V_{CC}$ , 5V  $V_{PP}$ ) independent of other blocks. Each block can be independently erased 100,000 times (6.4 million block erases per device). Block erase suspend mode allows system software to suspend block erase to read or write data from any other block.

A word/byte write is performed in byte increments typically within 9.24 $\mu$ s (5V V<sub>CC</sub>, 5V V<sub>PP</sub>). A multi word/byte write has high speed write performance of 2 $\mu$ s/byte (5V V<sub>CC</sub>, 5V V<sub>PP</sub>). (Multi) Word/byte write suspend mode enables the system to read data or

execute code from any other flash memory array location.

Individual block locking uses a combination of bits and WP#, Sixty-four block lock-bits, to lock and unlock blocks. Block lock-bits gate block erase, full chip erase and (multi) word/byte write operations. Block lock-bit configuration operations (Set Block Lock-Bit and Clear Block Lock-Bits commands) set and cleared block lock-bits.

The status register indicates when the WSM's block erase, full chip erase, (multi) word/byte write or block lock-bit configuration operation is finished.

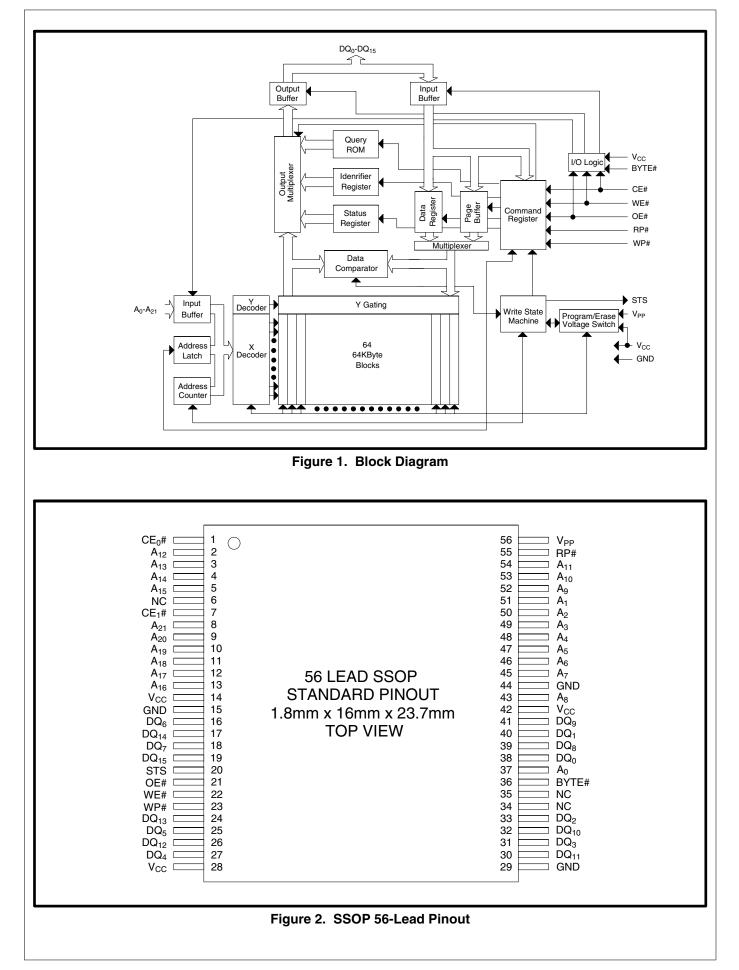
The STS output gives an additional indicator of WSM activity by providing both a hardware signal of status (versus software polling) and status masking (interrupt masking for background block erase, for example). Status polling using STS minimizes both CPU overhead and system power consumption. STS pin can be configured to different states using the Configuration command. The STS pin defaults to RY/BY# operation. When low, STS indicates that the WSM is performing a block erase, full chip erase, (multi) word/byte write or block lock-bit configuration. STS-High Z indicates that the WSM is ready for a new command, block erase is suspended and (multi) word/byte write are inactive, (multi) word/byte write are suspended, or the device is in deep power-down mode. The other 3 alternate configurations are all pulse mode for use as a system interrupt.

The access time is 90ns ( $t_{AVQV}$ ) over the commercial temperature range (0°C to +70°C) and V<sub>CC</sub> supply voltage range of 4.75V-5.25V. At lower V<sub>CC</sub> voltage, the access time is 100ns (4.5V-5.5V).

The Automatic Power Savings (APS) feature substantially reduces active current when the device is in static mode (addresses not switching). In APS mode, the typical  $I_{CCR}$  current is 1 mA at 5V  $V_{CC}$ .

When either  $CE_0^{\#}$  or  $CE_1^{\#}$ , and RP# pins are at  $V_{CC}$ , the  $I_{CC}$  CMOS standby mode is enabled. When the RP# pin is at GND, deep power-down mode is enabled which minimizes power consumption and provides write protection during reset. A reset time  $(t_{PHQV})$  is required from RP# switching high until outputs are valid. Likewise, the device has a wake time  $(t_{PHEL})$  from RP#-high until writes to the CUI are recognized. With RP# at GND, the WSM is reset and the status register is cleared.

The device is available in 56-Lead SSOP (Shrink Small Outline Package). Pinout is shown in Figure 2.



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Symbol	Tunn	Table 2. Pin Descriptions Name and Function
Symbol	Туре	
		ADDRESS INPUTS: Inputs for addresses during read and write operations. Addresses are
		internally latched during a write cycle.
A <sub>0</sub> -A <sub>21</sub>	INPUT	Ao: Byte Select Address. Not used in x16 mode(can be floated).
0 21		A1-A4: Column Address. Selects 1 of 16 bit lines.
		A5-A15: Row Address. Selects 1 of 2048 word lines.
		A16-A21 : Block Address.
		DATA INPUT/OUTPUTS:
		DQ0-DQ7: Inputs data and commands during CUI write cycles; outputs data during memory
		array, status register, query, and identifier code read cycles. Data pins float to high-
	INPUT/	impedance when the chip is deselected or outputs are disabled. Data is internally latched
DQ <sub>0</sub> -DQ <sub>15</sub>	OUTPUT	during a write cycle.
		DQ <sub>8</sub> -DQ <sub>15</sub> :Inputs data during CUI write cycles in x16 mode; outputs data during memory
		array read cycles in x16 mode; not used for status register, query and identifier code read
		mode. Data pins float to high-impedance when the chip is deselected, outputs are
		disabled, or in x8 mode(Byte#=V <sub>IL</sub> ). Data is internally latched during a write cycle.
CE <sub>0</sub> #,		CHIP ENABLE: Activates the device's control logic, input buffers decoders, and sense
CE <sub>1</sub> #	INPUT	amplifiers. Either $CE_0$ # or $CE_1$ # $V_{IH}$ deselects the device and reduces power consumption
1		to standby levels. Both $CE_0$ # and $CE_1$ # must be $V_{IL}$ to select the devices.
		<b>RESET/DEEP POWER-DOWN:</b> Puts the device in deep power-down mode and resets
RP#	INPUT	internal automation. RP# $V_{IH}$ enables normal operation. When driven $V_{IL}$ , RP# inhibits
	-	write operations which provides data protection during power transitions. Exit from deep
		power-down sets the device to read array mode.
OE#	INPUT	OUTPUT ENABLE: Gates the device's outputs during a read cycle.
WE#	INPUT	WRITE ENABLE: Controls writes to the CUI and array blocks. Addresses and data are
		latched on the rising edge of the WE# pulse.
		STS (RY/BY#): Indicates the status of the internal WSM. When configured in level mode
		(default mode), it acts as a RY/BY# pin. When low, the WSM is performing an internal
	OPEN	operation (block erase, full chip erase, (multi) word/byte write or block lock-bit
STS	DRAIN	configuration). STS High Z indicates that the WSM is ready for new commands, block
	OUTPUT	erase is suspended, and (multi) word/byte write is inactive, (multi) word/byte write is
		suspended or the device is in deep power-down mode. For alternate configurations of the
		STATUS pin, see the Configuration command.
WP#	INPUT	<b>WRITE PROTECT:</b> Master control for block locking. When $V_{IL}$ , Locked blocks can not be
	_	erased and programmed, and block lock-bits can not be set and reset.
		BYTE ENABLE: BYTE# V <sub>IL</sub> places device in x8 mode. All data is then input or output on
BYTE#	INPUT	$DQ_{0-7}$ , and $DQ_{8-15}$ float. BYTE# V <sub>IH</sub> places the device in x16 mode , and turns off the A <sub>0</sub>
		input buffer.
		BLOCK ERASE, FULL CHIP ERASE, (MULTI) WORD/BYTE WRITE, BLOCK LOCK-
		BIT CONFIGURATION POWER SUPPLY: For erasing array blocks, writing bytes or
V <sub>PP</sub>	SUPPLY	configuring block lock-bits. With V <sub>PP</sub> ≤V <sub>PPLK</sub> , memory contents cannot be altered. Block
		erase, full chip erase, (multi) word/byte write and block lock-bit configuration with an invalid
		V <sub>PP</sub> (see DC Characteristics) produce spurious results and should not be attempted.
		<b>DEVICE POWER SUPPLY:</b> Internal detection configures the device for 5V operation. Do
V	SUPPLY	not float any power pins. With $V_{CC} \leq V_{LKO}$ , all write attempts to the flash memory are
V <sub>CC</sub>	SOLILI	inhibited. Device operations at invalid V <sub>CC</sub> voltage (see DC Characteristics) produce
		spurious results and should not be attempted.
GND	SUPPLY	GROUND: Do not float any ground pins.
NC		<b>NO CONNECT:</b> Lead is not internal connected; it may be driven or floated.

#### 2 PRINCIPLES OF OPERATION

The LH28F320S5NS-L90 Flash memory includes an on-chip WSM to manage block erase, full chip erase, (multi) word/byte write and block lock-bit configuration functions. It allows for: 100% TTL-level control inputs, fixed power supplies during block erase, full chip erase, (multi) word/byte write and block lock-bit configuration, and minimal processor overhead with RAM-Like interface timings.

After initial device power-up or return from deep power-down mode (see Bus Operations), the device defaults to read array mode. Manipulation of external memory control pins allow array read, standby, and output disable operations.

Status register, query structure and identifier codes can be accessed through the CUI independent of the  $V_{PP}$  voltage. High voltage on  $V_{PP}$  enables successful block erase, full chip erase, (multi) word/byte write and block lock-bit configuration. All functions associated with altering memory contents—block erase, full chip erase, (multi) word/byte write and block lock-bit configuration, status, query and identifier codes—are accessed via the CUI and verified through the status register.

Commands are written using standard microprocessor write timings. The CUI contents serve as input to the WSM, which controls the block erase, full chip erase, (multi) word/byte write and block lock-bit configuration. The internal algorithms are regulated by the WSM, including pulse repetition, internal verification, and margining of data. Addresses and data are internally latch during write cycles. Writing the appropriate command outputs array data, accesses the identifier codes, outputs query structure or outputs status register data.

Interface software that initiates and polls progress of block erase, full chip erase, (multi) word/byte write and block lock-bit configuration can be stored in any block. This code is copied to and executed from system RAM during flash memory updates. After successful completion, reads are again possible via the Read Array command. Block erase suspend allows system software to suspend a block erase to read or write data from any other block. Write suspend allows system software to suspend a (multi) word/byte write to read data from any other flash memory array location.

#### 2.1 Data Protection

Depending on the application, the system designer may choose to make the  $V_{PP}$  power supply switchable (available only when block erase, full chip erase, (multi) word/byte write and block lock-bit configuration are required) or hardwired to  $V_{PPH1}$ . The device accommodates either design practice and encourages optimization of the processor-memory interface.

When  $V_{PP} \le V_{PPLK}$ , memory contents cannot be altered. The CUI, with multi-step block erase, full chip erase, (multi) word/byte write and block lock-bit configuration command sequences, provides protection from unwanted operations even when high voltage is applied to  $V_{PP}$ . All write functions are disabled when  $V_{CC}$  is below the write lockout voltage  $V_{LKO}$  or when RP# is at  $V_{IL}$ . The device's block locking capability provides additional protection from inadvertent code or data alteration by gating block erase, full chip erase and (multi) word/byte write operations.

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1FFFFF 1F0000	64-Kbyte Block	31
1EFFFF 1E0000	64-Kbyte Block	30
1DFFFF 1D0000	64-Kbyte Block	29
1CFFFF 1C0000	64-Kbyte Block	28
1BFFFF 1B0000	64-Kbyte Block	27
1AFFFF 1A0000	64-Kbyte Block	26
19FFFF	64-Kbyte Block	25
190000 18FFFF	64-Kbyte Block	24
180000 17FFFF	64-Kbyte Block	23
170000 16FFFF	64-Kbyte Block	22
160000 15FFFF	64-Kbyte Block	21
150000 14FFFF	64-Kbyte Block	20
140000 13FFFF	64-Kbyte Block	19
130000 12FFFF	64-Kbyte Block	18
120000 11FFFF	64-Kbyte Block	17
110000 10FFFF	64-Kbyte Block	16
100000 0FFFFF	-	15
0F0000 DEFFFF	64-Kbyte Block	
	64-Kbyte Block	14
	64-Kbyte Block	13
	64-Kbyte Block	12
0В0000	64-Kbyte Block	11
	64-Kbyte Block	10
09FFFF 090000	64-Kbyte Block	ç
08FFFF 080000	64-Kbyte Block	8
07FFFF 070000	64-Kbyte Block	7
06FFFF 060000	64-Kbyte Block	6
05FFFF 050000	64-Kbyte Block	5
04FFFF 040000	64-Kbyte Block	4
03FFFF 030000	64-Kbyte Block	З
02FFFF 020000	64-Kbyte Block	2
01FFFF 010000	64-Kbyte Block	1

3FFFFF	,	
3F0000	64-Kbyte Block	63
3EFFFF 3E0000	64-Kbyte Block	62
3DFFFF 3D0000	64-Kbyte Block	61
3CFFFF 3C0000	64-Kbyte Block	60
3BFFFF 3B0000	64-Kbyte Block	59
3AFFFF 3A0000	64-Kbyte Block	58
39FFFF	64-Kbyte Block	57
390000 38FFFF	64-Kbyte Block	56
380000 37FFFF	64-Kbyte Block	55
370000 36FFFF	64-Kbyte Block	54
360000 35FFFF	64-Kbyte Block	53
350000 34FFFF	64-Kbyte Block	52
340000 33FFFF	64-Kbyte Block	51
330000 32FFFF	64-Kbyte Block	50
320000 31FFFF	64-Kbyte Block	49
310000 30FFFF	64-Kbyte Block	48
300000 2FFFFF	64-Kbyte Block	47
2F0000 2EFFFF	,	47
2E0000 2DFFFF	64-Kbyte Block	
2D0000 2CFFFF	64-Kbyte Block	45
2C0000 2BFFFF	64-Kbyte Block	44
2B0000 2AFFFF	64-Kbyte Block	43
2A0000	64-Kbyte Block	42
29FFFF 290000	64-Kbyte Block	41
28FFFF 280000	64-Kbyte Block	40
27FFFF 270000	64-Kbyte Block	39
26FFFF 260000	64-Kbyte Block	38
25FFFF 250000	64-Kbyte Block	37
24FFFF 240000	64-Kbyte Block	36
23FFFF 230000	64-Kbyte Block	35
22FFFF 220000	64-Kbyte Block	34
21FFFF 210000	64-Kbyte Block	33
20FFFF	64-Kbyte Block	32
200000	L	

Figure 3. Memory Map

#### **3 BUS OPERATION**

The local CPU reads and writes flash memory insystem. All bus cycles to or from the flash memory conform to standard microprocessor bus cycles.

#### 3.1 Read

Information can be read from any block, identifier codes, query structure, or status register independent of the  $V_{PP}$  voltage. RP# must be at  $V_{IH}$ .

The first task is to write the appropriate read mode command (Read Array, Read Identifier Codes, Query or Read Status Register) to the CUI. Upon initial device power-up or after exit from deep power-down mode, the device automatically resets to read array mode. Five control pins dictate the data flow in and out of the component: CE# (CE<sub>0</sub>#, CE<sub>1</sub>#), OE#, WE#, RP# and WP#. CE<sub>0</sub>#, CE<sub>1</sub># and OE# must be driven active to obtain data at the outputs. CE<sub>0</sub>#, CE<sub>1</sub># is the device selection control, and when active enables the selected memory device. OE# is the data output (DQ<sub>0</sub>-DQ<sub>15</sub>) control and when active drives the selected memory data onto the I/O bus. WE# and RP# must be at V<sub>IH</sub>. Figure 17, 18 illustrates a read cycle.

#### 3.2 Output Disable

With OE# at a logic-high level ( $V_{IH}$ ), the device outputs are disabled. Output pins  $DQ_0$ - $DQ_{15}$  are placed in a high-impedance state.

#### 3.3 Standby

Either  $CE_0$ # or  $CE_1$ # at a logic-high level (V<sub>IH</sub>) places the device in standby mode which substantially reduces device power consumption.  $DQ_0$ - $DQ_{15}$ outputs are placed in a high-impedance state independent of OE#. If deselected during block erase, full chip erase, (multi) word/byte write and block lock-bit configuration, the device continues functioning, and consuming active power until the operation completes.

#### 3.4 Deep Power-Down

RP# at  $V_{II}$  initiates the deep power-down mode.

In read modes, RP#-low deselects the memory, places output drivers in a high-impedance state and turns off all internal circuits. RP# must be held low for a minimum of 100 ns. Time  $t_{PHQV}$  is required after return from power-down until initial memory access outputs are valid. After this wake-up interval, normal operation is restored. The CUI is reset to read array mode and status register is set to 80H.

During block erase, full chip erase, (multi) word/byte write or block lock-bit configuration modes, RP#-low will abort the operation. STS remains low until the reset operation is complete. Memory contents being altered are no longer valid; the data may be partially erased or written. Time  $t_{PHWL}$  is required after RP# goes to logic-high (V<sub>IH</sub>) before another command can be written.

As with any automated device, it is important to assert RP# during system reset. When the system comes out of reset, it expects to read from the flash memory. Automated flash memories provide status information when accessed during block erase, full chip erase, (multi) word/byte write and block lock-bit configuration. If a CPU reset occurs with no flash memory reset, proper CPU initialization may not occur because the flash memory may be providing status information instead of array data. SHARP's flash memories allow proper CPU initialization following a system reset through the use of the RP# input. In this application, RP# is controlled by the same RESET# signal that resets the system CPU.

#### 3.5 Read Identifier Codes Operation

The read identifier codes operation outputs the manufacturer code, device code, block status codes for each block (see Figure 4). Using the manufacturer and device codes, the system CPU can automatically match the device with its proper algorithms. The block status codes identify locked or unlocked block setting and erase completed or erase uncompleted condition.

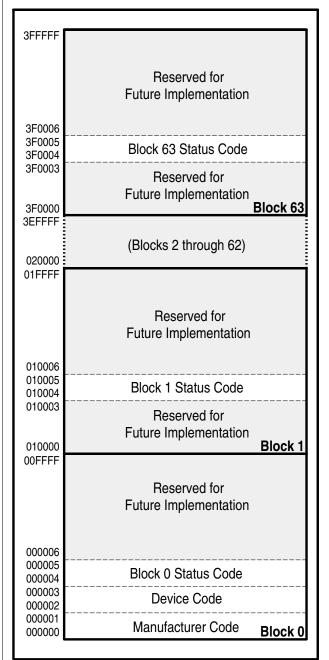


Figure 4. Device Identifier Code Memory Map

#### 3.6 Query Operation

The query operation outputs the query structure. Query database is stored in the 48Byte ROM. Query structure allows system software to gain critical information for controlling the flash component. Query structure are always presented on the lowestorder data output ( $DQ_0$ - $DQ_7$ ) only.

#### 3.7 Write

Writing commands to the CUI enable reading of device data and identifier codes. They also control inspection and clearing of the status register. When  $V_{CC}=V_{CC1/2}$  and  $V_{PP}=V_{PPH1}$ , the CUI additionally controls block erase, full chip erase, (multi) word/byte write and block lock-bit configuration.

The Block Erase command requires appropriate command data and an address within the block to be erased. The Word/byte Write command requires the command and address of the location to be written. Set Block Lock-Bit command requires the command and block address within the device (Block Lock) to be locked. The Clear Block Lock-Bits command requires the command and address within the device.

The CUI does not occupy an addressable memory location. It is written when WE# and CE# are active. The address and data needed to execute a command are latched on the rising edge of WE# or CE# (whichever goes high first). Standard microprocessor write timings are used. Figures 19 and 20 illustrate WE# and CE#-controlled write operations.

#### **4 COMMAND DEFINITIONS**

When the V<sub>PP</sub> voltage  $\leq$  V<sub>PPLK</sub>, Read operations from the status register, identifier codes, query, or blocks are enabled. Placing V<sub>PPH1</sub> on V<sub>PP</sub> enables successful block erase, full chip erase, (multi) word/byte write and block lock-bit configuration operations.

Device operations are selected by writing specific commands into the CUI. Table 4 defines these commands.

Table 3. Bus Operations(BYTE#=V <sub>IH</sub> )										
Mode	Notes	RP#	CE <sub>0</sub> #	CE <sub>1</sub> #	OE#	WE#	Address	V <sub>PP</sub>	DQ <sub>0-15</sub>	STS
Read	1,2,3,9	VIH	V	VII	VII	V <sub>IH</sub>	Х	X	D <sub>OUT</sub>	Х
Output Disable	3	V <sub>IH</sub>		V	V <sub>IH</sub>	V <sub>IH</sub>	X	Х	High Z	Х
Standby	3	V <sub>IH</sub>	V <sub>IH</sub> V <sub>IH</sub> V <sub>II</sub>	V <sub>IH</sub> V <sub>IL</sub> V <sub>IH</sub>	x	x	x	х	High Z	х
Deep Power-Down	4	VII	X	X	Х	Х	Х	Х	High Z	High Z
Read Identifier Codes	9	V <sub>IH</sub>	V <sub>IL</sub>	V <sub>IL</sub>	V <sub>IL</sub>	V <sub>IH</sub>	See Figure 4	Х	Note 5	High Z
Query	9	V <sub>IH</sub>	V <sub>IL</sub>	V <sub>IL</sub>	V <sub>IL</sub>	V <sub>IH</sub>	See Table 7~11	Х	Note 6	High Z
Write	3,7,8,9	VIH	VIL	VII	VIH	VIL	Х	Х	D <sub>IN</sub>	Х

Table 3.1. Bus Operations(BYTE#=V.,)

Mode	Notes	RP#	CE <sub>0</sub> #	CE <sub>1</sub> #	OE#	WE#	Address	V <sub>PP</sub>	DQ <sub>0-7</sub>	STS
Read	1,2,3,9	VIH	VII	VII	VII	VIH	X	X	D <sub>OUT</sub>	Х
Output Disable	3	V <sub>IH</sub>		V	V <sub>IH</sub>	V <sub>IH</sub>	X	Х	High Z	Х
Standby	3	V <sub>IH</sub>	V <sub>IH</sub> V <sub>IH</sub> V <sub>II</sub>	V <sub>IH</sub> V <sub>IL</sub> V <sub>IH</sub>	X	x	x	х	High Z	х
Deep Power-Down	4	VII	X	X	Х	Х	X	Х	High Z	High Z
Read Identifier Codes	9	V <sub>IH</sub>	V <sub>IL</sub>	V <sub>IL</sub>	V <sub>IL</sub>	V <sub>IH</sub>	See Figure 4	х	Note 5	High Z
Query	9	V <sub>IH</sub>	V <sub>IL</sub>	V <sub>IL</sub>	V <sub>IL</sub>	V <sub>IH</sub>	See Table 7~11	Х	Note 6	High Z
Write	3,7,8,9	VIH	V <sub>IL</sub>	V <sub>IL</sub>	VIH	V <sub>IL</sub>	X	Х	D <sub>IN</sub>	Х

#### NOTES:

Refer to DC Characteristics. When V<sub>PP</sub>≤V<sub>PPLK</sub>, memory contents can be read, but not altered.
 X can be V<sub>IL</sub> or V<sub>IH</sub> for control pins and addresses, and V<sub>PPLK</sub> or V<sub>PPH1</sub> for V<sub>PP</sub>. See DC Characteristics for V<sub>PPLK</sub> and V<sub>PPH1</sub> voltages.

3. STS is V<sub>OI</sub> (if configured to RY/BY# mode) when the WSM is executing internal block erase, full chip erase, (multi) word/byte write or block lock-bit configuration algorithms. It is floated during when the WSM is not busy, in block erase suspend mode with (multi) word/byte write inactive, (multi) word/byte write suspend mode, or deep power-down mode.

- 4. RP# at GND±0.2V ensures the lowest deep power-down current.
- 5. See Section 4.2 for read identifier code data.
- 6. See Section 4.5 for query data.
- 7. Command writes involving block erase, full chip erase, (multi) word/byte write or block lock-bit configuration are reliably executed when  $V_{PP}=V_{PPH1}$  and  $V_{CC}=V_{CC1/2}$ . 8. Refer to Table 4 for valid  $D_{IN}$  during a write operation. 9. Don't use the timing both OE# and WE# are  $V_{IL}$ .

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	Table 4.	Comma	and Defin	itions <sup>(10)</sup>					
	Bus Cycles	Notes		First Bus Cycle			Second Bus Cycle		
Command	Req'd		Oper <sup>(1)</sup>	Addr <sup>(2)</sup>	Data <sup>(3)</sup>	Oper <sup>(1)</sup>	Addr <sup>(2)</sup>	Data <sup>(3)</sup>	
Read Array/Reset	1		Write	Х	FFH	_			
Read Identifier Codes	≥2	4	Write	Х	90H	Read	IA	ID	
Query	≥2		Write	Х	98H	Read	QA	QD	
Read Status Register	2		Write	Х	70H	Read	Х	SRD	
Clear Status Register	1		Write	Х	50H				
Block Erase Setup/Confirm	2	5	Write	BA	20H	Write	BA	D0H	
Full Chip Erase Setup/Confirm	2		Write	Х	30H	Write	Х	D0H	
Word/Byte Write Setup/Write	2	5,6	Write	WA	40H	Write	WA	WD	
Alternate Word/Byte Write Setup/Write	2	5,6	Write	WA	10H	Write	WA	WD	
Multi Word/Byte Write Setup/Confirm	≥4	9	Write	WA	E8H	Write	WA	N-1	
Block Erase and (Multi) Word/byte Write Suspend	1	5	Write	Х	B0H				
Confirm and Block Erase and (Multi) Word/byte Write Resume	1	5	Write	Х	D0H				
Block Lock-Bit Set Setup/Confirm	2	7	Write	BA	60H	Write	BA	01H	
Block Lock-Bit Reset Setup/Confirm	2	8	Write	Х	60H	Write	Х	D0H	
STS Configuration Level-Mode for Erase and Write (RY/BY# Mode)	2		Write	х	B8H	Write	х	00H	
STS Configuration Pulse-Mode for Erase	2		Write	х	B8H	Write	Х	01H	
STS Configuration Pulse-Mode for Write	2		Write	Х	B8H	Write	Х	02H	
STS Configuration Pulse-Mode for Erase and Write	2		Write	Х	B8H	Write	Х	03H	

#### NOTES:

1. BUS operations are defined in Table 3 and Table 3.1.

2. X=Any valid address within the device.

IA=Identifier Code Address: see Figure 4.

QA=Query Offset Address.

BA=Address within the block being erased or locked.

WA=Address of memory location to be written.

 SRD=Data read from status register. See Table 14 for a description of the status register bits. WD=Data to be written at location WA. Data is latched on the rising edge of WE# or CE# (whichever goes high first).

ID=Data read from identifier codes.

QD=Data read from query database.

- 4. Following the Read Identifier Codes command, read operations access manufacturer, device and block status codes. See Section 4.2 for read identifier code data.
- If the block is locked, WP# must be at V<sub>IH</sub> to enable block erase or (multi) word/byte write operations. Attempts
  to issue a block erase or (multi) word/byte write to a locked block while RP# is V<sub>IH</sub>.
- 6. Either 40H or 10H are recognized by the WSM as the byte write setup.
- 7. A block lock-bit can be set while WP# is  $V_{IH}$ .
- 8. WP# must be at V<sub>IH</sub> to clear block lock-bits. The clear block lock-bits operation simultaneously clears all block lock-bits.
- 9. Following the Third Bus Cycle, inputs the write address and write data of 'N' times. Finally, input the confirm command 'D0H'.
- 10. Commands other than those shown above are reserved by SHARP for future device implementations and should not be used.

#### 4.1 Read Array Command

Upon initial device power-up and after exit from deep power-down mode, the device defaults to read array mode. This operation is also initiated by writing the Read Array command. The device remains enabled for reads until another command is written. Once the internal WSM has started a block erase, full chip erase, (multi) word/byte write or block lock-bit configuration, the device will not recognize the Read Array command until the WSM completes its operation unless the WSM is suspended via an Erase Suspend and (Multi) Word/byte Write Suspend command. The Read Array command functions independently of the V<sub>PP</sub> voltage and RP# must be V<sub>IH</sub>.

#### 4.2 Read Identifier Codes Command

The identifier code operation is initiated by writing the Read Identifier Codes command. Following the command write, read cycles from addresses shown in Figure 4 retrieve the manufacturer, device, block lock configuration and block erase status (see Table 5 for identifier code values). To terminate the operation, write another valid command. Like the Read Array command, the Read Identifier Codes command functions independently of the V<sub>PP</sub> voltage and RP# must be V<sub>IH</sub>. Following the Read Identifier Codes command, the following information can be read:

Code	Address A21-A0	Data
Manufacture Code	000000 000001	B0
Device Code	000002 000003	D4
Block Status Code	X0004 <sup>(1)</sup> X0005 <sup>(1)</sup>	
Block is Unlocked		DQ <sub>0</sub> =0
Block is Locked		DQ <sub>0</sub> =1
•Last erase operation completed successfully		DQ <sub>1</sub> =0
<ul> <li>Last erase operation did not completed successfully</li> </ul>		DQ <sub>1</sub> =1
Reserved for Future Use		DQ <sub>2-7</sub>

 Table 5. Identifier Codes

#### NOTE:

 X selects the specific block status code to be read. See Figure 4 for the device identifier code memory map.

#### 4.3 Read Status Register Command

The status register may be read to determine when a block erase, full chip erase, (multi) word/byte write or block lock-bit configuration is complete and whether the operation completed successfully(see Table 14). It may be read at any time by writing the Read Status Register command. After writing this command, all subsequent read operations output data from the status register until another valid command is written. The status register contents are latched on the falling edge of OE# or CE#(Either CE<sub>0</sub># or CE<sub>1</sub>#), whichever occurs. OE# or CE#(Either CE<sub>0</sub># or CE<sub>1</sub>#) must toggle to V<sub>IH</sub> before further reads to update the status register latch. The Read Status Register command functions independently of the V<sub>PP</sub> voltage. RP# must be V<sub>IH</sub>.

The extended status register may be read to determine multi word/byte write availability(see Table 14.1). The extended status register may be read at any time by writing the Multi Word/Byte Write command. After writing this command, all subsequent read operations output data from the extended status register, until another valid command is written. Multi Word/Byte Write command must be re-issued to update the extended status register latch.

#### 4.4 Clear Status Register Command

Status register bits SR.5, SR.4, SR.3 and SR.1 are set to "1"s by the WSM and can only be reset by the Clear Status Register command. These bits indicate various failure conditions (see Table 14). By allowing system software to reset these bits, several operations (such as cumulatively erasing or locking multiple blocks or writing several bytes in sequence) may be performed. The status register may be polled to determine if an error occurs during the sequence.

To clear the status register, the Clear Status Register command (50H) is written. It functions independently of the applied  $V_{PP}$  Voltage. RP# must be  $V_{IH}$ . This command is not functional during block erase, full chip erase, (multi) word/byte write block lock-bit configuration, block erase suspend or (multi) word/byte write suspend modes.

#### 4.5 Query Command

Query database can be read by writing Query command (98H). Following the command write, read cycle from address shown in Table 7~11 retrieve the critical information to write, erase and otherwise control the flash component.  $A_0$  of query offset address is ignored when X8 mode (BYTE#=V<sub>IL</sub>).

Query data are always presented on the low-byte data output ( $DQ_0$ - $DQ_7$ ). In x16 mode, high-byte ( $DQ_8$ - $DQ_{15}$ ) outputs 00H. The bytes not assigned to any information or reserved for future use are set to "0". This command functions independently of the V<sub>PP</sub> voltage. RP# must be V<sub>IH</sub>.

Table 6. Example of Query Structure Output								
Mode	Offset Address	Offset Address Output						
		DQ <sub>15~8</sub>	DQ <sub>7~0</sub>					
	A <sub>5</sub> , A <sub>4</sub> , A <sub>3</sub> , A <sub>2</sub> , A <sub>1</sub> , A <sub>0</sub>							
	1 , 0 , 0 , 0 , 0 , 0 (20H)	High-Z	"Q"					
X8 mode	1,0,0,0,0,1(21H)	High-Z	"Q"					
	1, 0,0,0,1,0 (22H)	High-Z	"R"					
	1,0,0,0,1,1(23H)	High-Z	"R"					
	A <sub>5</sub> , A <sub>4</sub> , A <sub>3</sub> , A <sub>2</sub> , A <sub>1</sub>							
X16 mode		00H	"Q"					
	1,0,0,0,1 (11H)	00H	"R"					

#### 4.5.1 Block Status Register

This field provides lock configuration and erase status for the specified block. These informations are only available when device is ready (SR.7=1). If block erase or full chip erase operation is finished irregulary, block erase status bit will be set to "1". If bit 1 is "1", this block is invalid.

	Table 7. Query Block Status Register					
Offset (Word Address)	Length	Description				
(BA+2)H	01H	Block Status Register bit0 Block Lock Configuration 0=Block is unlocked 1=Block is Locked bit1 Block Erase Status 0=Last erase operation completed successfully 1=Last erase operation not completed successfully bit2-7 reserved for future use				

#### Note:

1. BA=The beginning of a Block Address.

#### 4.5.2 CFI Query Identification String

The Identification String provides verification that the component supports the Common Flash Interface specification. Additionally, it indicates which version of the spec and which Vendor-specified command set(s) is(are) supported.

Offset (Word Address)	Length	Description
10H,11H,12H	03H	Query Unique ASCII string "QRY" 51H,52H,59H
13H,14H	02H	Primary Vendor Command Set and Control Interface ID Code 01H,00H (SCS ID Code)
15H,16H	02H	Address for Primary Algorithm Extended Query Table 31H,00H (SCS Extended Query Table Offset)
17H,18H	02H	Alternate Vendor Command Set and Control Interface ID Code 0000H (0000H means that no alternate exists)
19H,1AH	02H	Address for Alternate Algorithm Extended Query Table 0000H (0000H means that no alternate exists)

#### Table 8. CFI Query Identification String

#### 4.5.3 System Interface Information

The following device information can be useful in optimizing system interface software.

Offset (Word Address)	Length	Description	
1BH	01H	V <sub>CC</sub> Logic Supply Minimum Write/Erase voltage 45H (4.5V)	
1CH	01H	V <sub>CC</sub> Logic Supply Maximum Write/Erase voltage 55H (5.5V)	
1DH	01H	V <sub>PP</sub> Programming Supply Minimum Write/Erase voltage 45H (4.5V)	
1EH	01H	V <sub>PP</sub> Programming Supply Maximum Write/Erase voltage 55H (5.5V)	
1FH	01H	Typical Timeout per Single Byte/Word Write 04H (2 <sup>4</sup> =16µs)	
20H	01H	Typical Timeout for Maximum Size Buffer Write (32 Bytes) 06H (2 <sup>6</sup> =64µs)	
21H	01H	Typical Timeout per Individual Block Erase 09H (09H=9, 2 <sup>9</sup> =512ms)	
22H	01H	Typical Timeout for Full Chip Erase 0FH (0FH=15, 2 <sup>15</sup> =32768ms)	
23H	01H	Maximum Timeout per Single Byte/Word Write, 2 <sup>N</sup> times of typical. 04H (2 <sup>4</sup> =16, 16µsx16=256µs)	
24H	01H	Maximum Timeout Maximum Size Buffer Write, $2^{N}$ times of typical. 04H ( $2^{4}$ =16, 64µsx16=1024µs)	
25H	01H	Maximum Timeout per Individual Block Erase, 2 <sup>N</sup> times of typical. 04H (2 <sup>4</sup> =16, 1024msx16=16384ms)	
26H	01H	Maximum Timeout for Full Chip Erase, $2^{N}$ times of typical. 04H ( $2^{4}$ =16, 32768msx16=524288ms)	

#### Table 9. System Information String

#### 4.5.4 Device Geometry Definition

This field provides critical details of the flash device geometry.

Table 10. Device Geometry Definition				
Offset (Word Address)	Length	Description		
27H	01H	Device Size 16H (16H=22, 2 <sup>22</sup> =4194304=4M Bytes)		
28H,29H	02H	Flash Device Interface description 02H,00H (x8/x16 supports x8 and x16 via BYTE#)		
2AH,2BH	02H	Maximum Number of Bytes in Multi-byte 05H,00H (2 <sup>5</sup> =32 Bytes )		
2CH	01H	lumber of Erase Block Regions within device 1H (symmetrically blocked)		
2DH,2EH	02H	The Number of Erase Blocks 3FH,00H (3FH=63 ==> 63+1=64 Blocks)		
2FH,30H	02H	The Number of "256 Bytes" cluster in a Érase block 00H,01H (0100H=256 ==>256 Bytes x 256= 64K Bytes in a Erase Block)		

#### 4.5.5 SCS OEM Specific Extended Query Table

Certain flash features and commands may be optional in a vendor-specific algorithm specification. The optional vendor-specific Query table(s) may be used to specify this and other types of information. These structures are defined solely by the flash vendor(s).

Offset (Word Address)	Length	Description
31H,32H,33H	03H	PRI
		50H,52H,49H
34H	01H	31H (1) Major Version Number , ASCII
35H	01H	30H (0) Minor Version Number, ASCII
36H,37H,	04H	0FH,00H,00H
38H,39H		Optional Command Support
		bit0=1 : Chip Erase Supported
		bit1=1 : Suspend Erase Supported
		bit2=1 : Suspend Write Supported
		bit3=1 : Lock/Unlock Supported
		bit4=0 : Queued Erase Not Supported
		bit5-31=0 : reserved for future use
3AH	01H	01H
		Supported Functions after Suspend
		bit0=1 : Write Supported after Erase Suspend
		bit1-7=0 : reserved for future use
3BH,3CH	02H	03H,00H
		Block Status Register Mask
		bit0=1 : Block Status Register Lock Bit [BSR.0] active
		bit1=1 : Block Status Register Valid Bit [BSR.1] active
		bit2-15=0 : reserved for future use
3DH	01H	V <sub>CC</sub> Logic Supply Optimum Write/Erase voltage(highest performance) 50H(5.0V)
3EH	01H	V <sub>PP</sub> Programming Supply Optimum Write/Erase voltage(highest performance 50H(5.0V)
3FH	reserved	Reserved for future versions of the SCS Specification

#### Table 11. SCS OEM Specific Extended Query Table

#### 4.6 Block Erase Command

Block erase is executed one block at a time and initiated by a two-cycle command. A block erase setup is first written, followed by an block erase confirm. This command sequence requires appropriate sequencing and an address within the block to be erased (erase changes all block data to FFH). Block preconditioning, erase and verify are handled internally by the WSM (invisible to the system). After the two-cycle block erase sequence is written, the device automatically outputs status register data when read (see Figure 5). The CPU can detect block erase completion by analyzing the output data of the STS pin or status register bit SR.7.

When the block erase is complete, status register bit SR.5 should be checked. If a block erase error is detected, the status register should be cleared before system software attempts corrective actions. The CUI remains in read status register mode until a new command is issued.

This two-step command sequence of set-up followed by execution ensures that block contents are not accidentally erased. An invalid Block Erase command sequence will result in both status register bits SR.4 and SR.5 being set to "1". Also, reliable block erasure can only occur when  $V_{CC}=V_{CC1/2}$  and  $V_{PP}=V_{PPH1}$ . In the absence of this high voltage, block contents are protected against erasure. If block erase is attempted while  $V_{PP}\leq V_{PPLK}$ , SR.3 and SR.5 will be set to "1". Successful block erase requires that the corresponding block lock-bit be cleared or if set, that  $WP\#=V_{IH}$ . If block erase is attempted when the corresponding block lock-bit is set and  $WP\#=V_{IL}$ , SR.1 and SR.5 will be set to "1".

#### 4.7 Full Chip Erase Command

This command followed by a confirm command (D0H) erases all of the unlocked blocks. A full chip

erase setup is first written, followed by a full chip erase confirm. After a confirm command is written, device erases the all unlocked blocks from block 0 to Block 63 block by block. This command sequence requires appropriate sequencing. Block preconditioning, erase and verify are handled internally by the WSM (invisible to the system). After the two-cycle full chip erase sequence is written, the device automatically outputs status register data when read (see Figure 6). The CPU can detect full chip erase completion by analyzing the output data of the STS pin or status register bit SR.7.

When the full chip erase is complete, status register bit SR.5 should be checked. If erase error is detected, the status register should be cleared before system software attempts corrective actions. The CUI remains in read status register mode until a new command is issued. If error is detected on a block during full chip erase operation, WSM stops erasing the block and begin to erase the next block. Reading the block valid status by issuing Read ID Codes command or Query command informs which blocks failed to its erase.

This two-step command sequence of set-up followed by execution ensures that block contents are not accidentally erased. An invalid Full Chip Erase command sequence will result in both status register bits SR.4 and SR.5 being set to "1". Also, reliable full chip erasure can only occur when  $V_{CC}=V_{CC1/2}$  and  $V_{PP}=V_{PPH1}$ . In the absence of this high voltage, block contents are protected against erasure. If full chip erase is attempted while  $V_{PP}\leq V_{PPLK}$ , SR.3 and SR.5 will be set to "1". When WP#=V<sub>IH</sub>, all blocks are erased independent of block lock-bits status. When WP#=V<sub>IL</sub>, only unlocked blocks are erased. Full chip erase can not be suspended.

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#### 4.8 Word/Byte Write Command

Word/byte write is executed by a two-cycle command sequence. Word/Byte Write setup (standard 40H or alternate 10H) is written, followed by a second write that specifies the address and data (latched on the rising edge of WE#). The WSM then takes over, controlling the word/byte write and write verify algorithms internally. After the word/byte write sequence is written, the device automatically outputs status register data when read (see Figure 7). The CPU can detect the completion of the word/byte write event by analyzing the STS pin or status register bit SR.7.

When word/byte write is complete, status register bit SR.4 should be checked. If word/byte write error is detected, the status register should be cleared. The internal WSM verify only detects errors for "1"s that do not successfully write to "0"s. The CUI remains in read status register mode until it receives another command.

Reliable word/byte writes can only occur when  $V_{CC}=V_{CC1/2}$  and  $V_{PP}=V_{PPH1}$ . In the absence of this high voltage, memory contents are protected against word/byte writes. If word/byte write is attempted while  $V_{PP}\leq V_{PPLK}$ , status register bits SR.3 and SR.4 will be set to "1". Successful word/byte write requires that the corresponding block lock-bit be cleared or, if set, that WP#=V<sub>IH</sub>. If word/byte write is attempted when the corresponding block lock-bit is set and WP#=V<sub>IL</sub>, SR.1 and SR.4 will be set to "1". Word/byte write operations with  $V_{IL}<WP#<V_{IH}$  produce spurious results and should not be attempted.

#### 4.9 Multi Word/Byte Write Command

Multi word/byte write is executed by at least fourcycle or up to 35-cycle command sequence. Up to 32 bytes in x8 mode (16 words in x16 mode) can be loaded into the buffer and written to the Flash Array. First, multi word/byte write setup (E8H) is written with the write address. At this point, the device automatically outputs extended status register data (XSR) when read (see Figure 8, 9). If extended status register bit XSR.7 is 0, no Multi Word/Byte Write command is available and multi word/byte write setup which just has been written is ignored. To retry,

continue monitoring XSR.7 by writing multi word/byte write setup with write address until XSR.7 transitions to 1. When XSR.7 transitions to 1, the device is ready for loading the data to the buffer. A word/byte count (N)-1 is written with write address. After writing a word/byte count(N)-1, the device automatically turns back to output status register data. The word/byte count (N)-1 must be less than or equal to 1FH in x8 mode (0FH in x16 mode). On the next write, device start address is written with buffer data. Subsequent writes provide additional device address and data, depending on the count. All subsequent address must lie within the start address plus the count. After the final buffer data is written, write confirm (D0H) must be written. This initiates WSM to begin copying the buffer data to the Flash Array. An invalid Multi Word/Byte Write command sequence will result in both status register bits SR.4 and SR.5 being set to "1". For additional multi word/byte write, write another multi word/byte write setup and check XSR.7. The Multi Word/Byte Write command can be gueued while WSM is busy as long as XSR.7 indicates "1", because LH28F320S5NS-L90 has two buffers. If an error occurs while writing, the device will stop writing and flush next multi word/byte write command loaded in multi word/byte write command. Status register bit SR.4 will be set to "1". No multi word/byte write command is available if either SR.4 or SR.5 are set to "1". SR.4 and SR.5 should be cleared before issuing multi word/byte write command. If a multi word/byte write command is attempted past an erase block boundary, the device will write the data to Flash Array up to an erase block boundary and then stop writing. Status register bits SR.4 and SR.5 will be set to "1".

Reliable multi byte writes can only occur when  $V_{CC}=V_{CC1/2}$  and  $V_{PP}=V_{PPH1}$ . In the absence of this high voltage, memory contents are protected against multi word/byte writes. If multi word/byte write is attempted while  $V_{PP}\leq V_{PPLK}$ , status register bits SR.3 and SR.4 will be set to "1". Successful multi word/byte write requires that the corresponding block lock-bit be cleared or, if set, that WP#= $V_{IH}$ . If multi byte write is attempted when the corresponding block lock-bit is set and WP#= $V_{IL}$ , SR.1 and SR.4 will be set to "1".

#### 4.10 Block Erase Suspend Command

The Block Erase Suspend command allows blockerase interruption to read or (multi) word/byte-write data in another block of memory. Once the blockerase process starts, writing the Block Erase Suspend command requests that the WSM suspend the block erase sequence at a predetermined point in the algorithm. The device outputs status register data when read after the Block Erase Suspend command is written. Polling status register bits SR.7 and SR.6 can determine when the block erase operation has been suspended (both will be set to "1"). STS will also transition to High Z. Specification t<sub>WHRH2</sub> defines the block erase suspend latency.

At this point, a Read Array command can be written to read data from blocks other than that which is suspended. A (Multi) Word/Byte Write command sequence can also be issued during erase suspend to program data in other blocks. Using the (Multi) Word/Byte Write Suspend command (see Section 4.11), a (multi) word/byte write operation can also be suspended. During a (multi) word/byte write operation with block erase suspended, status register bit SR.7 will return to "0" and the STS (if set to RY/BY#) output will transition to  $V_{OL}$ . However, SR.6 will remain "1" to indicate block erase suspend status.

The only other valid commands while block erase is suspended are Read Status Register and Block Erase Resume. After a Block Erase Resume command is written to the flash memory, the WSM will continue the block erase process. Status register bits SR.6 and SR.7 will automatically clear and STS will return to  $V_{OL}$ . After the Erase Resume command is written, the device automatically outputs status register data when read (see Figure 10).  $V_{PP}$  must remain at  $V_{PPH1}$  (the same  $V_{PP}$  level used for block erase) while block erase is suspended. RP# must also remain at  $V_{IH}$ . Block erase cannot resume until

(multi) word/byte write operations initiated during block erase suspend have completed.

#### 4.11 (Multi) Word/Byte Write Suspend Command

The (Multi) Word/Byte Write Suspend command allows (multi) word/byte write interruption to read data in other flash memory locations. Once the (multi) word/byte write process starts, writing the (Multi) Word/Byte Write Suspend command requests that the WSM suspend the (multi) word/byte write sequence at a predetermined point in the algorithm. The device continues to output status register data when read after the (Multi) Word/Byte Write Suspend command is written. Polling status register bits SR.7 and SR.2 can determine when the (multi) word/byte write operation has been suspended (both will be set to "1"). STS will also transition to High Z. Specification  $t_{WHRH1}$  defines the (multi) word/byte write suspend latency.

At this point, a Read Array command can be written to read data from locations other than that which is suspended. The only other valid commands while (multi) word/byte write is suspended are Read Status Register and (Multi) Word/Byte Write Resume. After (Multi) Word/Byte Write Resume command is written to the flash memory, the WSM will continue the (multi) word/byte write process. Status register bits SR.2 and SR.7 will automatically clear and STS will return to VOL. After the (Multi) Word/Byte Write command is written, the device automatically outputs status register data when read (see Figure 11). V<sub>PP</sub> must remain at V<sub>PPH1</sub> (the same V<sub>PP</sub> level used for (multi) word/byte write) while in (multi) word/byte write suspend mode. WP# must also remain at VIH or V<sub>IL</sub>.

#### 4.12 Set Block Lock-Bit Command

A flexible block locking and unlocking scheme is enabled via block lock-bits. The block lock-bits gate program and erase operations With  $WP\#=V_{IH}$ , individual block lock-bits can be set using the Set Block Lock-Bit command. See Table 13 for a summary of hardware and software write protection options.

Set block lock-bit is executed by a two-cycle command sequence. The set block lock-bit setup along with appropriate block or device address is written followed by either the set block lock-bit confirm (and an address within the block to be locked). The WSM then controls the set block lock-bit algorithm. After the sequence is written, the device automatically outputs status register data when read (see Figure 12). The CPU can detect the completion of the set block lock-bit event by analyzing the STS pin output or status register bit SR.7.

When the set block lock-bit operation is complete, status register bit SR.4 should be checked. If an error is detected, the status register should be cleared. The CUI will remain in read status register mode until a new command is issued.

This two-step sequence of set-up followed by execution ensures that block lock-bits are not accidentally set. An invalid Set Block Lock-Bit command will result in status register bits SR.4 and SR.5 being set to "1". Also, reliable operations occur only when  $V_{CC}=V_{CC1/2}$  and  $V_{PP}=V_{PPH1}$ . In the absence of this high voltage, block lock-bit contents are protected against alteration.

A successful set block lock-bit operation requires WP#=V<sub>IH</sub>. If it is attempted with WP#=V<sub>IL</sub>, SR.1 and SR.4 will be set to "1" and the operation will fail. Set block lock-bit operations with WP#<V<sub>IH</sub> produce spurious results and should not be attempted.

#### 4.13 Clear Block Lock-Bits Command

All set block lock-bits are cleared in parallel via the Clear Block Lock-Bits command. With WP#= $V_{IH}$ ,

block lock-bits can be cleared using only the Clear Block Lock-Bits command. See Table 13 for a summary of hardware and software write protection options.

Clear block lock-bits operation is executed by a twocycle command sequence. A clear block lock-bits setup is first written. After the command is written, the device automatically outputs status register data when read (see Figure 13). The CPU can detect completion of the clear block lock-bits event by analyzing the STS Pin output or status register bit SR.7.

When the operation is complete, status register bit SR.5 should be checked. If a clear block lock-bit error is detected, the status register should be cleared. The CUI will remain in read status register mode until another command is issued.

This two-step sequence of set-up followed by execution ensures that block lock-bits are not accidentally cleared. An invalid Clear Block Lock-Bits command sequence will result in status register bits SR.4 and SR.5 being set to "1". Also, a reliable clear block lock-bits operation can only occur when  $V_{CC}=V_{CC1/2}$  and  $V_{PP}=V_{PPH1}$ . If a clear block lock-bits operation is attempted while V<sub>PP</sub>≤V<sub>PPLK</sub>, SR.3 and SR.5 will be set to "1". In the absence of this high voltage, the block lock-bits content are protected against alteration. A successful clear block lock-bits operation requires WP#=VIH. If it is attempted with  $WP\#=V_{II}$ , SR.1 and SR.5 will be set to "1" and the operation will fail. Clear block lock-bits operations with V<sub>IH</sub><RP# produce spurious results and should not be attempted.

If a clear block lock-bits operation is aborted due to  $V_{PP}$  or  $V_{CC}$  transitioning out of valid range or RP# active transition, block lock-bit values are left in an undetermined state. A repeat of clear block lock-bits is required to initialize block lock-bit contents to known values.

#### 4.14 STS Configuration Command

The Status (STS) pin can be configured to different states using the STS Configuration command. Once the STS pin has been configured, it remains in that configuration until another configuration command is issued, the device is powered down or RP# is set to  $V_{II}$ . Upon initial device power-up and after exit from deep power-down mode, the STS pin defaults to RY/BY# operation where STS low indicates that the WSM is busy. STS High Z indicates that the WSM is ready for a new operation.

To reconfigure the STS pin to other modes, the STS Configuration is issued followed by the appropriate configuration code. The three alternate configurations are all pulse mode for use as a system interrupt. The STS Configuration command functions independently of the  $V_{PP}$  voltage and RP# must be  $V_{IH}$ .

Table 12. 313	Configuration Coding Description
Configuration Bits	Effects
00H	Set STS pin to default level mode (RY/BY#). RY/BY# in the default level-mode of operation will indicate WSM status condition.
01H	Set STS pin to pulsed output signal for specific erase operation. In this mode, STS provides low pulse at the completion of BLock Erase, Full Chip Erase and Clear Block Lock-bits operations.
02H	Set STS pin to pulsed output signal for a specific write operation. In this mode, STS provides low pulse at the completion of (Multi) Byte Write and Set Block Lock-bit operation.
03H	Set STS pin to pulsed output signal for specific write and erase operation. STS provides low pulse at the completion of Block Erase, Full Chip Erase, (Multi) Word/Byte Write and Block Lock-bit Configuration operations.

Operation Block Lock-Bit WP#		WP#	Effect
Block Erase,	0	V <sub>II</sub> or V <sub>IH</sub>	Block Erase and (Multi) Word/Byte Write Enabled
(Multi) Word/Byte Write	1	V <sub>IL</sub>	Block is Locked. Block Erase and (Multi) Word/Byte Write Disabled
		V <sub>IH</sub>	Block Lock-Bit Override. Block Erase and (Multi) Word/Byte Write Enabled
Full Chip Erase	0,1	V <sub>II</sub>	All unlocked blocks are erased, locked blocks are not erased
	Х	VIH	All blocks are erased
Set Block Lock-Bit	Х	V <sub>II</sub>	Set Block Lock-Bit Disabled
		V <sub>IH</sub>	Set Block Lock-Bit Enabled
Clear Block Lock-Bits	Х	VII	Clear Block Lock-Bits Disabled
		V <sub>IH</sub>	Clear Block Lock-Bits Enabled

#### **Table 13. Write Protection Alternatives**

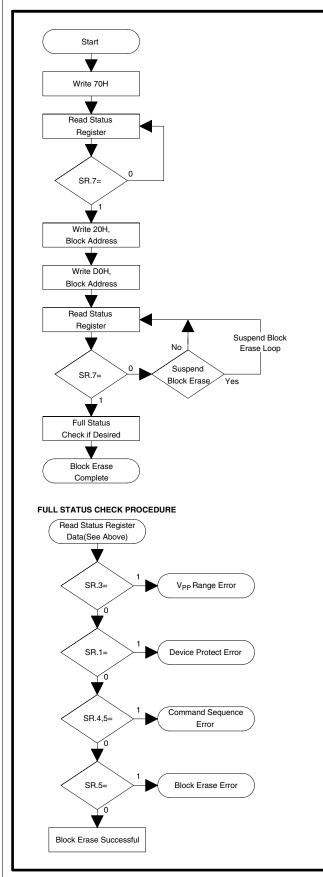
Table 12. STS	<b>Configuration Coding Description</b>	
<b>.</b>		

#### LHF32K10

7     6     5     4     3     2     1     0       IR.7 = WRITE STATE MACHINE STATUS     1     Ready     0     NOTES:       1 = Beady     0     Block Erase Suspended     0     Block Erase in Progress/Completed     NOTES:       IR.6 = BLOCK ERASE SUSPEND STATUS     1     Block Erase in Progress/Completed     SR.6 -0 are invalid wind/byte write or block lock-bit configuration completion.       IR.5 = ERASE AND CLEAR BLOCK LOCK-BITS STATUS     1     Form in Erase or Clear Block Lock-Bits       I = Error in Erase or Clear Block Lock-Bits     SR.3 does not provide a continuous indication of Vpp level. The WSM interrogates and indicates the Vpp level.       I = Virte AND SET ELOCK LOCK-BIT     SR.3 does not provide a continuous indication of Vpp level.       I = Virte INP STATUS     SR.1 does not provide a continuous indication of block lock-bit.       I = Virte SUSPEND STATUS     SR.1 does not provide a continuous indication of block lock-bit.       I = Virte SUSPEND STATUS     SR.1 does not provide a continuous indication of block lock-bit.       I = DEVICE PROTECT STATUS     SR.1 does not provide a continuous indication of block lock-bit configuration commant sequences. It informs the system, depending on the attempted operation, if the block lock-bit status.       I = DEVICE PROTECT STATUS     SR.0 is reserved for future use and should be masked out when polling the status register.       I = Malti H1. Extended Status Register Definition     SR.7 = STATE MACHINE STATUS	WSMS	BESS	ECBLBS	e 14. Status WSBLBS	VPPS	WSS	DPS	R
R.7 = WRITE STATE MACHINE STATUS1 = ReadyCheck STS or SR.7 to determine block erase, full chip erase, (multi) word/byte write or block lock-bit configuration completion.R.6 = BLOCK ERASE SUSPEND STATUSSR.6-0 are invalid while SR.7-"0".1 = Biock Erase SuspendedSR.6-0 are invalid write or block lock-bit configuration or STS configuration attempt, an impropriod 	7	6	5	4	3	2	1	0
IR.7 = WRITE STATE MACHINE STATUS       1 = Ready         1 = Ready       Check STS or SR.7 to determine block erase, full chip erase, (multi) word/byte write or block lock-bit configuration completion.         IR.6 = BLOCK ERASE SUSPEND STATUS       SR.6-0 are invalid while SR.7="0".         1 = Biock Erase Suspended       SR.6-0 are invalid while SR.7="0".         I = Biock Erase Suspended       SR.6-0 are invalid while SR.7="0".         I = Error in Erase or Clear Block Lock-Bits       SR.3 does not provide a continuous indicates the V <sub>PP</sub> level. The WSM interrogates and indicates the V <sub>PP</sub> level. The WSM interrogates and indicates the V <sub>PP</sub> level. The WSM interrogates and indicates the V <sub>PP</sub> level. The WSM interrogates and indicates the V <sub>PP</sub> level. The WSM interrogates block lock-bit configuration or stres, full chip erase, (multi) word/byte write or block lock-bit configuration or stres equences. SR.3 is not guaranteed to reports accurate feedback only when V <sub>PP</sub> =V <sub>PPH1</sub> .         IR.3 = V <sub>PP</sub> STATUS       SR.1 does not provide a continuous indication of block lock-bit is set and/or WP# only after block lock-bit configuration or omman sequences. It informs the system, depending on the attempted operation, if the block lock-bit is set and/or WP# is not V <sub>H+</sub> . Reading the block lock-bit configuration comman sequences. It informs the system, depending on the attempted operation, if the block lock-bit is set and/or WP# is not V <sub>H+</sub> . Reading the block lock-bit is set and/or WP# is not V <sub>H+</sub> . Reading the block lock-bit set at and/or WP# is not V <sub>H+</sub> . Reading the block lock-bit is set and/or WP# is not V <sub>H+</sub> . Reading the block lock-bit is set and/or WP# is not V <sub>H+</sub> . Reading the block lock-bit is set and/or WP# is not V <sub>H+</sub> . Reading the block lock-bit is set and/or WP					NOTES:			
0 = Busy       erase, (multi) word/byte write or block lock-bit configuration completion.         R.6 = BLOCK ERASE SUSPEND STATUS       SR.6 - are invalid winie SR.7="0".         1 = Block Erase Suspended       SR.6 - are invalid winie SR.7="0".         R.5 = ERASE AND CLEAR BLOCK LOCK-BITS STATUS       SR.6 - are invalid winie SR.7="0".         1 = Error in Erase or Clear Block Lock-Bits       Iboth SR.5 and SR.4 are "1"s after a block lock-bit configuration or STS configuration attempt, an improprior of St Block Lock-Bits         R.4 = WRITE AND SET BLOCK LOCK-BIT STATUS       SR.3 does not provide a continuous indicates the V <sub>pp</sub> level. The WSM interrogates and indicates the V <sub>pp</sub> level. The WSM interrogates block lock-bit configuration command sequences. St.3 is not guaranteed to reports accurate feedback only when V <sub>pp=7</sub> V <sub>pp+1</sub> .         R.3 = V <sub>pp</sub> STATUS       SR.1 does not provide a continuous indication of block lock-bit configuration command sequences. It informs the system, depending on the attempted operation, if the block lock-bit configuration comman sequences. It informs the system, depending on the attempted operation, if the block lock-bit status.         R.1 = DEVICE PROTECT STATUS       SR.0 is reserved for future use and should be masked out when polling the status register.         B.1 = DEVICE PROTECT STATUS       SR.0 is reserved for future use and should be masked out when polling the status register.         0 = Unlock       SR.1 a R       R       R       R       R       R       R       R       R       R       R       R       R <td>R.7 = WRITE</td> <td>STATE MA</td> <td>CHINE STATUS</td> <td>5</td> <td></td> <td></td> <td></td> <td></td>	R.7 = WRITE	STATE MA	CHINE STATUS	5				
R.6 = BLOCK ERASE SUSPEND STATUS       a Biock Erase Suspended         1 = Biock Erase in Progress/Completed       SR.6-0 are invalid while SR.7="0".         R.5 = ERASE AND CLEAR BLOCK LOCK-BITS       SR.6-0 are invalid while SR.7="0".         I botk Erase or Clear Bloch Lock-Bits       I both SR.5 and SR.4 are "1"s after a block erase, full thip erase, (multi) word/byte write, block lock-bit configuration of STS configuration attempt, an improprommand sequence was entered.         I = Error in Write or Set Block Lock-Bits       SR.3 does not provide a continuous indication of V <sub>PP</sub> level. The WSM interrogates and indicates the V <sub>PP</sub> erase, (multi) word/byte write or block lock-bit configuration command sequences. SR.3 is not guaranteed to reports accurate feedback only when V <sub>PP</sub> =V <sub>PPH1</sub> .         R.3 = V <sub>PP</sub> STATUS       SR.1 does not provide a continuous indication of block lock-bit values. The WSM interrogates block lock-bit, and WP# only after block corese, full chip erase, (multi) word/byte write or block lock-bit, and VPP oN         R.1 = DEVICE PROTECT STATUS       SR.1 does not provide a continuous indication of block lock-bit and/or WP# is not V <sub>H</sub> . Reading the block lock-bit is set and/or WP# is not V <sub>H</sub> . Reading the block lock-bit is and/or WP# is not V <sub>H</sub> . Reading the block lock-bit status.         R.1 = DEVICE PROTECT STATUS       SR.0 is reserved for future use and should be masked out when polling the status register.         R.1 = DEVICE PROTECT STATUS       SR.0 is reserved for future use and should be masked out when polling the status register.         SR.2 = RESERVED FOR FUTURE ENHANCEMENTS       A fter issue a Multi Word/Byte Write command					Check STS or	SR.7 to determ	nine block era	se, full chip
R.6 = BLOCK ERASE SUSPEND STATUS       SR.6-0 are invalid while SR.7="0".         1 = Block Erase Suspended       Iboth SR.5 and SR.4 are "1"s after a block lock-bit configuration or STS configuration attempt, an improprocommend sequence was entered.         1 = Error in Erase or Clear BlocL Lock-Bits       Iboth SR.5 and SR.4 are "1"s after a block lock-bit configuration or STS configuration attempt, an improprocommend sequence was entered.         1 = Error in Erase or Clear BlocL Lock-Bits       SR.3 does not provide a continuous indicates the V <sub>pp</sub> lev only after block erase, full chip erase, (multi) word/byte write or block lock-bit configuration command sequences. SR.3 is not guaranteed to reports accurate feedback only when V <sub>pp=</sub> Zv <sub>ppH1</sub> .         R.3 = V <sub>pp</sub> STATUS       SR.1 does not provide a continuous indication of block lock-bit values. The WSM interrogates block lock-bit is alues, (multi) word/byte write or block lock-bit colck-bit configuration command sequences. It informs the system, depending on the otwork lock-bit satus.         R.1 = DEVICE PROTECT STATUS       SR.1 enserved for future use and should be masked out when polling the status register.         R.1 = DEVICE PROTECT STATUS       SR.0 is reserved for future use and should be masked out when polling the status register.         R.0 = RESERVED FOR FUTURE ENHANCEMENTS       SR.0 is reserved for future use and should be masked out when polling the status register.         SR.7 = STATE MACHINE STATUS       1 = Multi Word/Byte Write not available       After issue a Multi Word/Byte Write command i available.         SR.6-0 = RESERVED FOR FUTURE ENHANCEMENTS       SR.6 -0 is reserve	0 = Busy						or block lock-	-bit
1 = Block Erase Suspended       1         0 = Block Erase in Progress/Completed       If both SR.5 and SR.4 are "1"s after a block erase, full chip erase, (multi) word/byte write, block lock-bit configuration or STS configuration attempt, an improprommand sequence was entered.         1 = Error in Erase or Clear Block Lock-Bits       SR.4 are "1"s after a block tock-bit configuration or STS configuration attempt, an improprommand sequence was entered.         1 = Error in Write or Set Block Lock-Bits       SR.3 does not provide a continuous indicates the V <sub>PP</sub> level. The WSM interrogates and indicates the V <sub>PP</sub> level. The WSM interrogates and indicates the V <sub>PP</sub> level.         1 = Vpp STATUS       SR.1 does not provide a continuous indication of V <sub>PP</sub> level.         1 = Vpp Low Detect, Operation Abort       SR.1 does not provide a continuous indication of block lock-bit configuration command sequences. SR.3 is not guaranteed to reports accurate feedback only when V <sub>PP</sub> ≠V <sub>PPH1</sub> .         SR.1 does not provide a continuous indication of block lock-bit values. The WSM interrogates block lock-bit configuration command sequences. SR.1 does not provide a continuous indication of block lock-bit values. The WSM interrogates block lock-bit configuration command sequences. SR.1 does not provide a continuous indication of block lock-bit values. The WSM interrogates block lock-bit is and/or WP# took Detected, Operation Abort         0 = Unlock       SR.1 erase R       R       R       R       R       R       R       R       R       R       R       R       R       R       R       R       R       R				10				
0 = Block Erase in Progress/Completed       If both SR.5 and SR.4 are "1"s after a block erase, full chip erase, (multi) word/byte write, block lock-bit chip erase, (multi) word/byte write, block lock-bit comfiguration or STS configuration attempt, an impropriod mand sequence was entered.         1 = Error in Erase or Clear Block Lock-Bits       SR.3 does not provide a continuous indicates the V <sub>pp</sub> level. The WSM interrogates and indicates the V <sub>pp</sub> level. The WSM interrogates and indicates the V <sub>pp</sub> level. The WSM interrogates and indicates the V <sub>pp</sub> level. The WSM interrogates and indicates the V <sub>pp</sub> level. The WSM interrogates and indicates the V <sub>pp</sub> level. The WSM interrogates and indicates the V <sub>pp</sub> level. The WSM interrogates and indicates the V <sub>pp</sub> level. The WSM interrogates and indicates the V <sub>pp</sub> level. The WSM interrogates and indicates the V <sub>pp</sub> level. The WSM interrogates and indicates the V <sub>pp</sub> level. The WSM interrogates and indicates the V <sub>pp</sub> level. The WSM interrogates and indicates the V <sub>pp</sub> level. The WSM interrogates and indicates the V <sub>pp</sub> level. The WSM interrogates and indicates the V <sub>pp</sub> level. The WSM interrogates and indicates the V <sub>pp</sub> level. The WSM interrogates and indicates the V <sub>pp</sub> level. The WSM interrogates and indicates the V <sub>pp</sub> level. The WSM interrogates and indicates the V <sub>pp</sub> level. The WSM interrogates and indicates the V <sub>pp</sub> level. The WSM interrogates block lock-bit and very orby atter block lock-bit and WP# only atter block lock-bit configuration command sequences. SR.3 is not graving the block lock-bit and WP# only atter block lock-bit and WP# only atter block lock-bit is set and/or WP# is not V <sub>im</sub> . Reading the block lock corting unton a tempted operation, if the block lock-bit is set and/or WP# is not V <sub>im</sub> . Reading the block lock cot configuration command sequences. It informs the system, depending on the anticates that therot V <sub>im</sub> . Reading thenolock lock-bit is set and/or WP# is not V <sub>i</sub>				15	SR.6-0 are inv	alid while SR./	="0".	
chip erase, (multi) word/byte write, block lock-bit configuration or STS configuration attempt, an impropri command sequence was entered.1 = Error in Erase or Clear Block Lock-Bits 0 = Successful Erase or Clear Block Lock-Bit 1 = Error in Write or Set Block Lock-Bit 0 = Successful Write or Set Block Lock-BitSR.3 does not provide a continuous indication of $V_{PP}$ level. The WSM interrogates and indicates the $V_{Pp}$ lev only after block lock-bit configuration command sequences. SR.3 is not guaranteed to reports accurate feedback only when $V_{PP} \neq V_{PPH1}$ .IR.3 = $V_{PP}$ STATUS 1 = $V_{PP}$ Low Detect, Operation Abort 0 = $V_{PP}$ OKSR.1 does not provide a continuous indication of block lock-bit values. The WSM interrogates block lock-bit, and WP# only after block lock-bit configuration command sequences. It informs the system, depending on the attempted operation, Abort 0 = UnlockI = DEVICE PROTECT STATUS 1 = Block Lock-Bit and/or WP# Lock Detected, Operation Abort 0 = UnlockSR.1 does not provide a continuous indication of block lock-bit configuration command sequences. It informs the system, depending on the attempted operation, if the block lock-bit configuration codes after writing the Read Identifier Codes commann indicates block lock-bit status. SR.0 is reserved for future use and should be masked out when polling the status register.SMSRRRRR76543211 = Multi Word/Byte Write cort available 0 = Multi Word/Byte Write not availableNOTES: After issue a Multi Word/Byte Write command i available.SR.6-0 is reserved for future use and should be				4	If both SB 5 a	nd SR 4 are "1'	's after a bloc	k oraso full
IR-5 = ERASE AND CLEAR BLOCK LOCK-BITS STATUSI = Error in Erase or Clear Block Lock-Bitsconfiguration or STS configuration attempt, an improprison command sequences was entered.I = Error in Erase or Clear Block Lock-BitsSR.3 does not provide a continuous indication of $V_{PP}$ level. The WSM interrogates and indicates the $V_{PP}$ level.IR.4 = WRITE AND SET BLOCK LOCK-BIT STATUSSR.3 does not provide a continuous indication of block only after block lock-bit configuration command sequences. SR.3 is not guaranteed to reports accurate feedback only when $V_{PP} \neq V_{PPH1}$ .IR.3 = V <sub>PP</sub> STATUSSR.1 does not provide a continuous indication of block lock-bit values. The WSM interrogates block lock-bit only after block lock-bit configuration command sequences. It informs the system, depending on the attempted operation, if the block lock-bit set and/or WP# is not V <sub>IH</sub> . Reading the block lock-bit set and/or WP# is not V <sub>IH</sub> . Reading the block lock configuration codes after writing the Read Identifier Codes commani indicates block lock-bit status.I = DEVICE PROTECT STATUS 1 = Block Lock-Bit and/or WP# Lock Detected, Operation AbortSR.0 is reserved for future use and should be masked out when polling the status register.IR.0 = RESERVED FOR FUTURE ENHANCEMENTSSR.0 is reserved for future use and should be masked out when yold/byte Write availableSR.7 = STATE MACHINE STATUS 1 = Multi Word/Byte Write availableNOTES: After issue a Multi Word/Byte Write command i available.SR.6-0 = RESERVED FOR FUTURE ENHANCEMENTSAfter issue a Multi Word/Byte Write command i available.			grees, completed					
STATUS1 = Error in Erase or Clear Block Lock-Bits0 = Successful Erase or Clear Block Lock-Bits1 = Error in Write or Set Block Lock-Bits1 = Error in Write or Set Block Lock-Bits0 = Successful Write or Set Block Lock-Bits1 = Vrite or Dock lock-Dit configuration command1 = Vrite Suspended0 = Write in Progress/Completed1R.1 = DEVICE PROTECT STATUS1 = Block Lock-Bit and/or WP# Lock Detected, Operation Abort0 = Unlock10 = Unlock12 = RESERVED FOR FUTURE ENHANCEMENTS13 = Mick Sit and/or WP# Lock Detected, Operation Abort0 = Unlock14 = 14.1. Extended Status Register Definition15 = SMS1 = Multi Word/Byte Write available1 = Multi Word/Byte Write available1 = Multi Word/Byte Write available1 = Multi Word/Byte Write not available1 = Multi Word/Byte Write not available1 = SIGK - 0=RESERVED FOR FUTURE ENHANCEMENTS1 = Multi Word/Byte Write not available1 = SR.6-0=RESERVED FOR FUTURE ENHANCEMENTS	R.5 = ERASE	AND CLEA	AR BLOCK LOC	K-BITS				
0 = Successful Erase or Clear Block Lock-BitSR.3 does not provide a continuous indicates the $V_{PP}$ IR.4 = WRITE AND SET BLOCK LOCK-BIT STATUSI = Error in Write or Set Block Lock-BitI error in Write or Set Block Lock-Bit0 = Successful Write or Set Block Lock-Bitwrite or block lock-bit configuration commandsequences. SR.3 is not guaranteed to reports accurateIR.3 = $V_{PP}$ STATUSI = $V_{PP}$ Low Detect, Operation AbortSR.1 does not provide a continuous indication of block1 = $V_{PP}$ Low Detect, Operation AbortI = Write SuspendedSR.1 does not provide a continuous indication of blockIR.2 = WRITE SUSPEND STATUSI = Mvite SuspendedSR.1 does not provide a continuous indication of block lock-bit values. The WSM interrogates block lock-bitIR.1 = DEVICE PROTECT STATUSI = Block Lock-Bit and/or WP# Lock Detected, Operation AbortSR.0 is reserved for future use and should be masked out when polling the status register.IR.0 = RESERVED FOR FUTURE ENHANCEMENTSSR.0 is reserved for future use and should be masked out when polling the status register.SMSRRRR7654321 = Multi Word/Byte Write availableNOTES:1 = Multi Word/Byte Write not availableAfter issue a Multi Word/Byte Write command: XSR.71 = Multi Word/Byte Write not availableSR.6-0 is reserved for future use and should be	STATU	IS						/ I I
Ievel. The WSM interrogates and indicates the $V_{PP}$ few only after block crase, full chip erase, (multi) word/byte write or block lock-bit configuration command sequences. SR.3 is not guaranteed to reports accurate feedback only when $V_{PP} \neq V_{PPH1}$ .IR.3 = $V_{PP}$ Low Detect, Operation Abort 0 = V_{PP} OKSR.1 does not provide a continuous indication of block lock-bit values. The WSM interrogates block lock-bit, and WP# only after block erase, full chip erase, (multi) word/byte write or block lock-bit configuration comman sequences. It informs the system, depending on the attempted operation, if the block lock-bit is set and/or WP# is not $V_{IH}$ . Reading the block lock configuration codes after writing the Read Identifier Codes comman indicates block lock-bit status. RR.0 = RESERVED FOR FUTURE ENHANCEMENTSTable 14.1. Extended Status Register Definition SR.7 = STATE MACHINE STATUS 1 = Multi Word/Byte Write not available 0 = Multi Word/Byte Write not available SR.6-0 = RESERVED FOR FUTURE ENHANCEMENTSNOTES: After issue a Multi Word/Byte Write command: XSR.7 indicates that a next Multi Word/Byte Write command i available.								
R.4 = WRITE AND SET BLOCK LOCK-BIT STATUS       only after block case, full chip erase, (multi) word/byte         1 = Error in Write or Set Block Lock-Bit       only after block erase, full chip erase, (multi) word/byte         R.3 = V <sub>PP</sub> STATUS       sequences. SR.3 is not guaranteed to reports accurate feedback only when $V_{PP} \neq V_{PPH1}$ .         R.3 = V <sub>PP</sub> Dow Detect, Operation Abort       SR.1 does not provide a continuous indication of block lock-bit values. The WSM interrogates block lock-bit, and WP# only after block erase, full chip erase, (multi) word/byte write or block lock-bit configuration command sequences. It informs the system, depending on the attempted operation, if the block lock-bit is set and/or         R.1 = DEVICE PROTECT STATUS       SR.1 does not provide a continuous indication of block lock-bit values. The WSM interrogates block lock-bit, and WP# only after block lock-bit is set and/or         R.1 = DEVICE PROTECT STATUS       attempted operation, if the block lock-bit is set and/or         WP# is not V <sub>H</sub> . Reading the block lock-bit is set and/or       WP# is not V <sub>H</sub> . Reading the block lock configuration comman indicates block lock-bit status.         0 = Unlock       SR.0 is reserved for future use and should be masked out when polling the status register.         R.0 = RESERVED FOR FUTURE ENHANCEMENTS       SR.7 = STATE MACHINE STATUS         1 = Multi Word/Byte Write available       After issue a Multi Word/Byte Write command: XSR.7 indicates that a next Multi Word/Byte Write command: XSR.7 indicates that a next Multi Word/Byte Write command: XSR.7 indicates that a next Multi Word/Byte Write command: available.	0 = Succes	sful Erase c	or Clear Block Lo	ock-Bits				
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feedback only when $V_{PP} \neq V_{PPH1}$ . $R.3 = V_{PP}$ Low Detect, Operation Abort $SR.1$ does not provide a continuous indication of block $1 = V_{PP}$ OK $SR.1$ does not provide a continuous indication of block $R.2 = WRITE$ SUSPEND STATUS $SR.1$ does not provide a continuous indication of block $1 = Write$ Suspended $OCK$ -bit values. The WSM interrogates block lock-bit, $0 = Write in Progress/CompletedV_{IH}-Reading the block lock-bit set and/orR.1 = DEVICE PROTECT STATUSSR.0 is reserved for future use and should be masked1 = Block Lock-Bit and/or WP# Lock Detected,Operation AbortOFR FUTURE ENHANCEMENTSR.0 = RESERVED FOR FUTURE ENHANCEMENTSSR.0 is reserved for future use and should be maskedout when polling the status register.SR.7 = STATE MACHINE STATUS1 = Multi Word/Byte Write available0 = Multi Word/Byte Write not availableSR.6-0=RESERVED FOR FUTURE ENHANCEMENTSAfter issue a Multi Word/Byte Write command: XSR.7indicates that a next Multi Word/Byte Write command: available.SR.6-0 = RESERVED FOR FUTURE ENHANCEMENTSXSR.6-0 is reserved for future use and should be$								
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0 = V <sub>PP</sub> OK         iR.2 = WRITE SUSPEND STATUS         1 = Write Suspended         0 = Write in Progress/Completed         iR.1 = DEVICE PROTECT STATUS         1 = Block Lock-Bit and/or WP# Lock Detected, Operation Abort         0 = Unlock         iR.0 = RESERVED FOR FUTURE ENHANCEMENTS         Table 14.1. Extended Status Register Definition         SR.7 = STATE MACHINE STATUS         1 = Multi Word/Byte Write available         0 = Multi Word/Byte Write not available         SR.6-0=RESERVED FOR FUTURE ENHANCEMENTS					,			
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0 = Write in Progress/Completed       attempted operation, if the block lock-bit is set and/or         R.1 = DEVICE PROTECT STATUS       attempted operation, if the block lock-bit is set and/or         1 = Block Lock-Bit and/or WP# Lock Detected, Operation Abort       WP# is not V <sub>IH</sub> . Reading the block lock configuration codes after writing the Read Identifier Codes comman indicates block lock-bit status.         0 = Unlock       SR.0 = RESERVED FOR FUTURE ENHANCEMENTS         Table 14.1. Extended Status Register Definition         SR.0 = RESERVED FOR FUTURE ENHANCEMENTS         Motion WP# Lock Detected, Operation Abort         SR.0 = RESERVED FOR FUTURE ENHANCEMENTS         SR.7 = STATE MACHINE STATUS         1 = Multi Word/Byte Write available       NOTES:         After issue a Multi Word/Byte Write command: XSR.7 indicates that a next Multi Word/Byte Write command available.         SR.6-0 = RESERVED FOR FUTURE ENHANCEMENTS			51A105					
Image: Second			Completed					
Image: R.1 = DEVICE PROTECT STATUS       codes after writing the Read Identifier Codes commany indicates block lock-bit status.         1 = Block Lock-Bit and/or WP# Lock Detected, Operation Abort       codes after writing the Read Identifier Codes commany indicates block lock-bit status.         0 = Unlock       SR.0 is reserved for future use and should be masked out when polling the status register.         Table 14.1. Extended Status Register Definition         SMS         R       R         7       6       5         7       6       5         4       3       2         2       1         0 = Multi Word/Byte Write available       After issue a Multi Word/Byte Write command: XSR.7         0 = Multi Word/Byte Write not available       After issue a Multi Word/Byte Write command: XSR.7         3       2       1         0 = Multi Word/Byte Write not available       XSR.6-0 is reserved for future use and should be								
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0 = Unlock       SR.0 is reserved for future use and should be masked out when polling the status register.         SR.0 = RESERVED FOR FUTURE ENHANCEMENTS       SR.0 is reserved for future use and should be masked out when polling the status register.         Table 14.1. Extended Status Register Definition         SMS       R       R       R       R         7       6       5       4       3       2       1       0         SR.7 = STATE MACHINE STATUS       1 = Multi Word/Byte Write available       NOTES:       After issue a Multi Word/Byte Write command: XSR.7 indicates that a next Multi Word/Byte Write command: XSR.7 indicates that a next Multi Word/Byte Write command i available.         SR.6-0=RESERVED FOR FUTURE ENHANCEMENTS       XSR.6-0 is reserved for future use and should be			Or WP# LOCK De	etected,	indicates block	< lock-bit status	i.	
On the total of the one one of the	•				SR 0 is record	od for futuro us	o and should	ho maskod
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SR.6-0=RESERVED FOR FUTURE ENHANCEMENTS XSR.6-0 is reserved for future use and should be	0 = Multi W	/ord/Byte W	rite not available	•		a next Multi Wo	ord/Byte Write	command
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	3n.0-0=hES					erved for futur	a lise and eho	uld he

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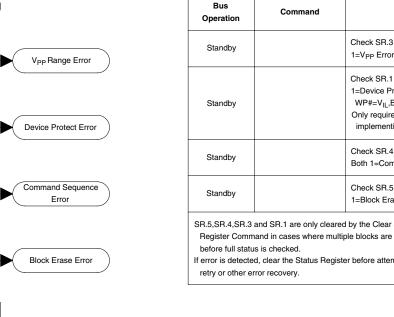


Bus Operation	Command	Comments
Write	Read Status Register	Data=70H Addr=X
Read		Status Register Data
Standby		Check SR.7 1=WSM Ready 0=WSM Busy
Write	Erase Setup	Data=20H Addr=Within Block to be Erased
Write	Erase Confirm	Data=D0H Addr=Within Block to be Erased
Read		Status Register Data
Standby		Check SR.7 1=WSM Ready 0=WSM Busy

Repeat for subsequent block erasures. Full status check can be done after each block erase or after a sequence of

block erasures.

Write FFH after the last operation to place device in read array mode.



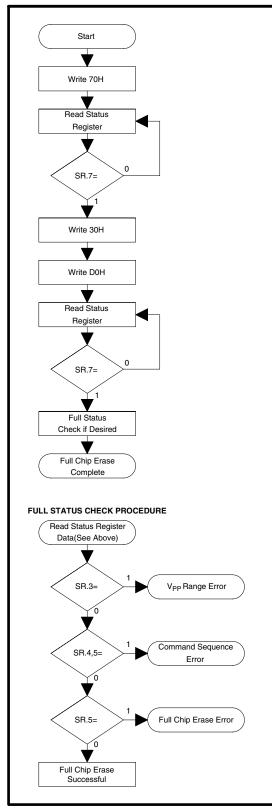
#### 1=V<sub>PP</sub> Error Detect Check SR.1 1=Device Protect Detect WP#=VIL,Block Lock-Bit is Set Only required for systems implementing lock-bit configuration Check SR.4,5 Both 1=Command Sequence Error Check SR.5 1=Block Erase Error SR.5, SR.4, SR.3 and SR.1 are only cleared by the Clear Status Register Command in cases where multiple blocks are erased If error is detected, clear the Status Register before attempting

Comments

#### Figure 5. Automated Block Erase Flowchart

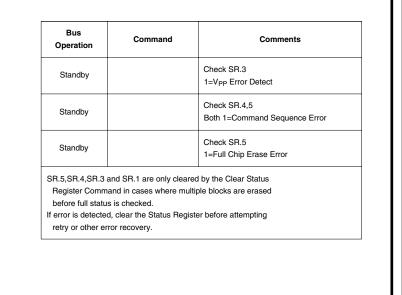
22

#### LHF32K10



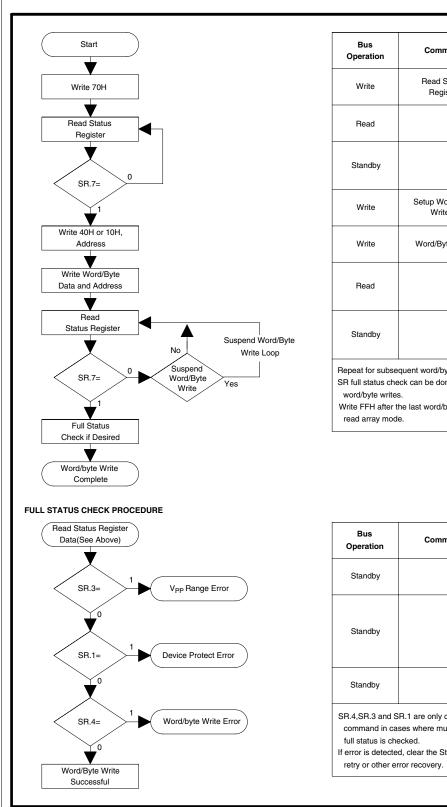
Bus Operation	Command	Comments
Write	Read Status Register	Data=70H Addr=X
Read		Status Register Data
Standby		Check SR.7 1=WSM Ready 0=WSM Busy
Write	Full Chip Erase Setup	Data=30H Addr=X
Write	Full Chip Erase Confirm	Data=D0H Addr=X
Read		Status Register Data
Standby		Check SR.7 1=WSM Ready 0=WSM Busy

Write FFH after the last operation to place device in read array mode.





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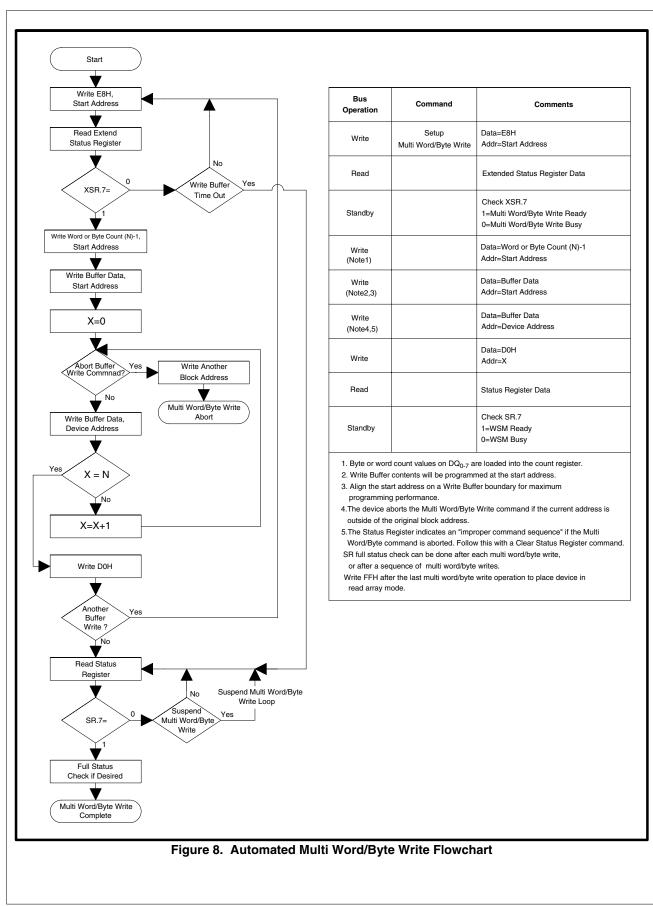
Bus Operation	Command	Comments
Write	Read Status Register	Data=70H Addr=X
Read		Status Register Data
Standby		Check SR.7 1=WSM Ready 0=WSM Busy
Write	Setup Word/Byte Write	Data=40H or 10H Addr=Location to Be Written
Write	Word/Byte Write	Data=Data to Be Written Addr=Location to Be Written
Read		Status Register Data
Standby		Check SR.7 1=WSM Ready 0=WSM Busy
Repeat for subse	quent word/byte writes.	

SR full status check can be done after each word/byte write, or after a sequence of word/byte writes.

Write FFH after the last word/byte write operation to place device in

### Bus Operation Command Comments Standby Check SR.3 1=Vpp Error Detect Check SR.1 1=Device Protect Detect Standby Check SR.1 1=Device Protect Detect WP#=VIL,Block Lock-Bit is Set Only required for systems implementing lock-bit configuration Standby Check SR.4 1=Data Write Error SR.4,SR.3 and SR.1 are only cleared by the Clear Status Register command in cases where multiple locations are written before full status is checked. If error is detected, clear the Status Register before attempting retry or other error recovery.

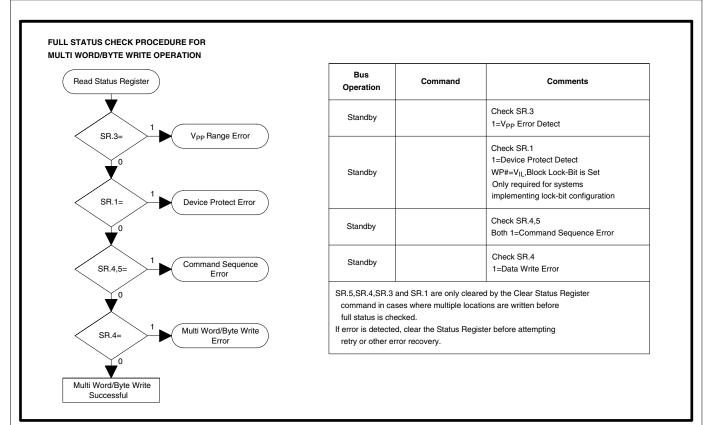
Figure 7. Automated Word/byte Write Flowchart



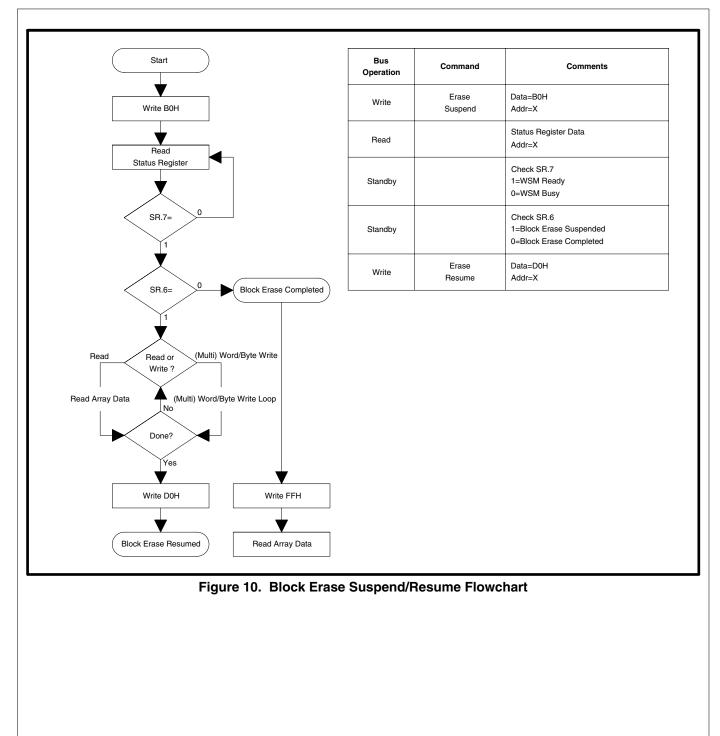
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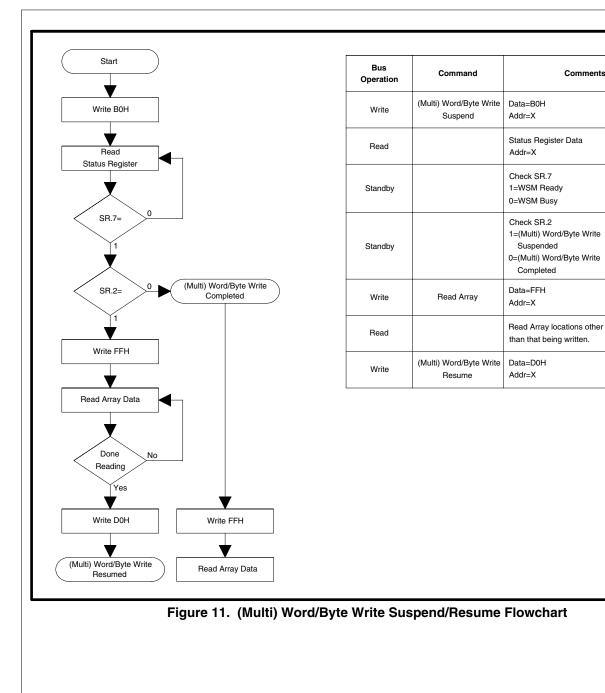








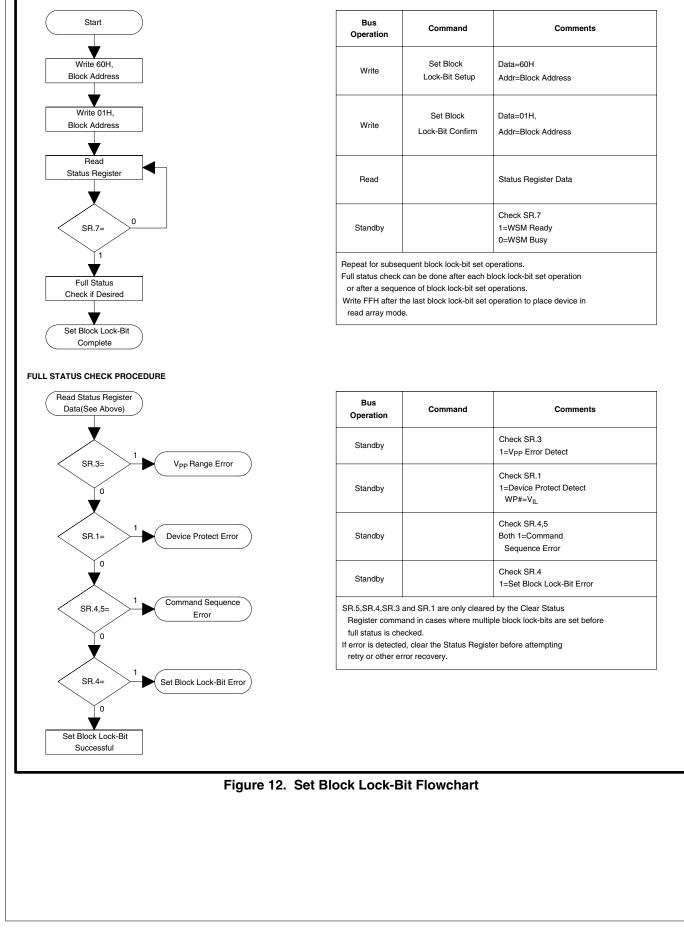
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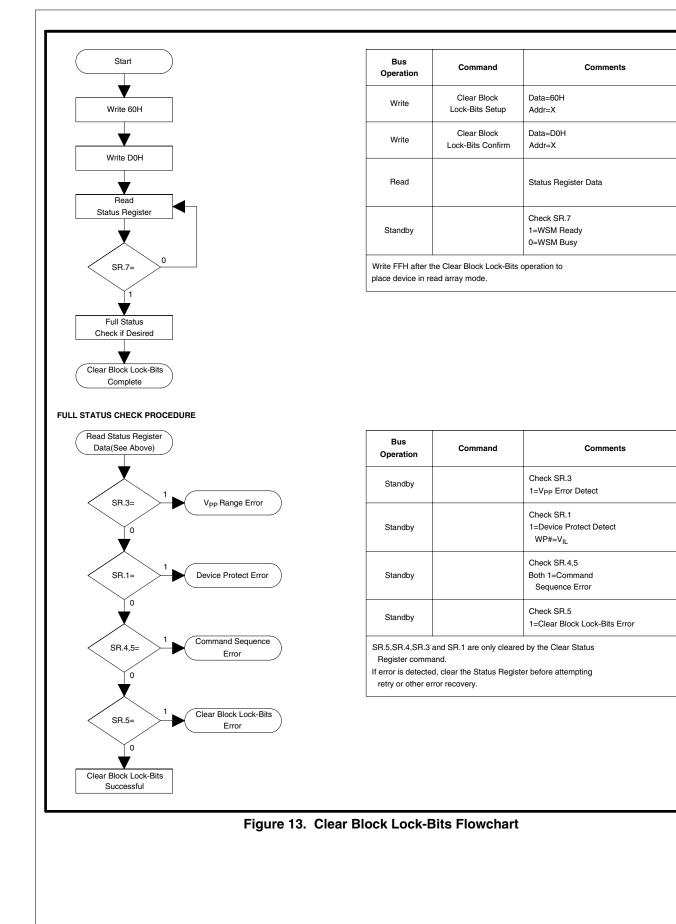




Comments

#### LHF32K10





#### **5 DESIGN CONSIDERATIONS**

#### 5.1 Three-Line Output Control

The device will often be used in large memory arrays. SHARP provides three control inputs to accommodate multiple memory connections. Three-Line control provides for:

- a. Lowest possible memory power dissipation.
- b. Complete assurance that data bus contention will not occur.

To use these control inputs efficiently, an address decoder should enable CE# while OE# should be connected to all memory devices and the system's READ# control line. This assures that only selected memory devices have active outputs while deselected memory devices are in standby mode. RP# should be connected to the system POWERGOOD signal to prevent unintended writes during system power transitions. POWERGOOD should also toggle during system reset.

#### 5.2 STS and Block Erase, Full Chip Erase, (Multi) Word/Byte Write and Block Lock-Bit Configuration Polling

STS is an open drain output that should be connected to V<sub>CC</sub> by a pullup resistor to provide a hardware method of detecting block erase, full chip erase, (multi) word/byte write and block lock-bit configuration completion. In default mode, it transitions low after block erase, full chip erase, (multi) word/byte write or block lock-bit configuration commands and returns to  $V_{OH}$  when the WSM has finished executing the internal algorithm. For alternate STS pin configurations, the see Configuration command.

STS can be connected to an interrupt input of the system CPU or controller. It is active at all times.

STS, in default mode, is also High-Z when the device is in block erase suspend (with (multi) word/byte write inactive), (multi) word/byte write suspend or deep power-down modes.

#### 5.3 Power Supply Decoupling

Flash memory power switching characteristics require careful device decoupling. System designers are interested in three supply current issues; standby current levels, active current levels and transient peaks produced by falling and rising edges of CE# and OE#. Transient current magnitudes depend on the device outputs' capacitive and inductive loading. Two-line control and proper decoupling capacitor selection will suppress transient voltage peaks. Each device should have a 0.1 µF ceramic capacitor connected between its  $V_{CC}$  and GND and between its  $V_{PP}$  and GND. These high-frequency, low inductance capacitors should be placed as close as possible to package leads. Additionally, for every eight devices, a 4.7 µF electrolytic capacitor should be placed at the array's power supply connection between V<sub>CC</sub> and GND. The bulk capacitor will overcome voltage slumps caused by PC board trace inductance.

#### 5.4 V<sub>PP</sub> Trace on Printed Circuit Boards

Updating flash memories that reside in the target system requires that the printed circuit board designer pay attention to the  $V_{PP}$  Power supply trace. The  $V_{PP}$  pin supplies the memory cell current for block erase, full chip erase, (multi) word/byte write and block lock-bit configuration. Use similar trace widths and layout considerations given to the  $V_{CC}$  power bus. Adequate  $V_{PP}$  supply traces and decoupling will decrease  $V_{PP}$  voltage spikes and overshoots.

#### 5.5 V<sub>CC</sub>, V<sub>PP</sub>, RP# Transitions

Block erase, full chip erase, (multi) word/byte write and block lock-bit configuration are not guaranteed if  $V_{\text{PP}}$  falls outside of a valid  $V_{\text{PPH1}}$  range,  $V_{\text{CC}}$  falls outside of a valid  $V_{CC1/2}$  range, or RP#= $V_{IL}$ . If  $V_{PP}$ error is detected, status register bit SR.3 is set to "1' along with SR.4 or SR.5, depending on the attempted operation. If RP# transitions to V<sub>II</sub> during block erase, full chip erase, (multi) word/byte write or block lock-bit configuration, STS(if set to RY/BY# mode) will remain low until the reset operation is complete. Then, the operation will abort and the device will enter deep power-down. The aborted operation may leave data partially altered. Therefore, the command sequence must be repeated after normal operation is restored. Device power-off or RP# transitions to V<sub>II</sub> clear the status register.

The CUI latches commands issued by system software and is not altered by  $V_{PP}$  or CE# transitions or WSM actions. Its state is read array mode upon power-up, after exit from deep power-down or after  $V_{CC}$  transitions below  $V_{LKO}$ .

After block erase, full chip erase, (multi) word/byte write or block lock-bit configuration, even after  $V_{PP}$  transitions down to  $V_{PPLK}$ , the CUI must be placed in read array mode via the Read Array command if subsequent access to the memory array is desired.

#### 5.6 Power-Up/Down Protection

The device is designed to offer protection against accidental block and full chip erasure, (multi) word/byte writing or block lock-bit configuration during power transitions. Upon power-up, the device is indifferent as to which power supply ( $V_{PP}$  or  $V_{CC}$ )

powers-up first. Internal circuitry resets the CUI to read array mode at power-up.

A system designer must guard against spurious writes for V<sub>CC</sub> voltages above V<sub>LKO</sub> when V<sub>PP</sub> is active. Since both WE# and CE# must be low for a command write, driving either to V<sub>IH</sub> will inhibit writes. The CUI's two-step command sequence architecture provides added level of protection against data alteration.

In-system block lock and unlock capability prevents inadvertent data alteration. The device is disabled while  $RP#=V_{II}$  regardless of its control inputs state.

#### 5.7 Power Dissipation

When designing portable systems, designers must consider battery power consumption not only during device operation, but also for data retention during system idle time. Flash memory's nonvolatility increases usable battery life because data is retained when system power is removed.

In addition, deep power-down mode ensures extremely low power consumption even when system power is applied. For example, portable computing products and other power sensitive applications that use an array of devices for solid-state storage can consume negligible power by lowering RP# to  $V_{IL}$  standby or sleep modes. If access is again needed, the devices can be read following the  $t_{PHQV}$  and  $t_{PHWL}$  wake-up cycles required after RP# is first raised to  $V_{IH}$ . See AC Characteristics— Read Only and Write Operations and Figures 17, 18, 19, 20 for more information.

#### 6 ELECTRICAL SPECIFICATIONS

#### 6.1 Absolute Maximum Ratings\*

Operating Temperature During Read, Erase, Write and Block Lock-Bit Configuration0°C to +70°C <sup>(1)</sup> Temperature under Bias10°C to +80°C
Storage Temperature65°C to +125°C
Voltage On Any Pin (except V <sub>CC</sub> , V <sub>PP</sub> )
$V_{CC}$ Suply Voltage0.2V to +7.0V <sup>(2)</sup>
V <sub>PP</sub> Update Voltage during Erase, Write and Block Lock-Bit Configuration0.2V to +7.0V <sup>(2)</sup>
Output Short Circuit Current 100mA <sup>(3)</sup>

\*WARNING: Stressing the device beyond the "Absolute Maximum Ratings" may cause permanent damage. These are stress ratings only. Operation beyond the "Operating Conditions" is not recommended and extended exposure beyond the "Operating Conditions" may affect device reliability.

#### NOTES:

- 1. Operating temperature is for commercial temperature product defined by this specification.
- 2. All specified voltages are with respect to GND. Minimum DC voltage is -0.5V on input/output pins and -0.2V on  $V_{CC}$  and  $V_{PP}$  pins. During transitions, this level may undershoot to -2.0V for periods <20ns. Maximum DC voltage on input/output pins and  $V_{CC}$  is  $V_{CC}$ +0.5V which, during transitions, may overshoot to  $V_{CC}$ +2.0V for periods <20ns.
- 3. Output shorted for no more than one second. No more than one output shorted at a time.

#### 6.2 Operating Conditions

#### Temperature and V<sub>CC</sub> Operating Conditions

Symbol	Parameter	Min.	Max.	Unit	Test Condition			
T <sub>A</sub>	Operating Temperature	0	+70	°C	Ambient Temperature			
V <sub>CC1</sub>	V <sub>CC</sub> Supply Voltage (5V±0.25V)	4.75	5.25	V				
V <sub>CC2</sub>	V <sub>CC</sub> Supply Voltage (5V±0.5V)	4.50	5.50	V				

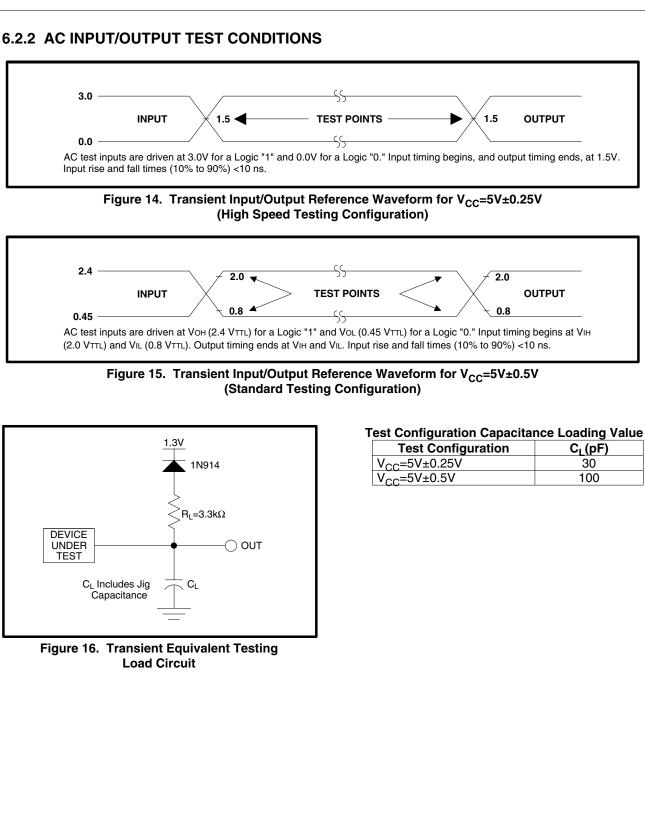
#### 6.2.1 CAPACITANCE<sup>(1)</sup>

T <sub>Δ</sub> =+25°C, f=1MHz								
Symbol	Parameter	Тур.	Max.	Unit	Condition			
C <sub>IN</sub>	Input Capacitance	7	10	pF	V <sub>IN</sub> =0.0V			
C <sub>OUT</sub>	Output Capacitance	9	12	pF	V <sub>OUT</sub> =0.0V			
NATE.					•••			

NOTE:

1. Sampled, not 100% tested.

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## 6.2.3 DC CHARACTERISTICS

		DC C	haracteris	tics		
			V <sub>C</sub>	<sub>C</sub> =5V		Test
Symbol	Parameter	Notes	Тур.	Max.	Unit	Conditions
ILI	Input Load Current	1		±1	μA	V <sub>CC</sub> =V <sub>CC</sub> Max. V <sub>IN</sub> =V <sub>CC</sub> or GND
I <sub>LO</sub>	Output Leakage Current	1		±10	μA	V <sub>CC</sub> =V <sub>CC</sub> Max. V <sub>OUT</sub> =V <sub>CC</sub> or GND
I <sub>CCS</sub>	V <sub>CC</sub> Standby Current	1,3,6	25	100	μΑ	CMOS Inputs V <sub>CC</sub> =V <sub>CC</sub> Max. CE#=RP#=V <sub>CC</sub> ±0.2V
			2	4	mA	TTL Inputs V <sub>CC</sub> =V <sub>CC</sub> Max. CE#=RP#=V <sub>IH</sub>
I <sub>CCD</sub>	V <sub>CC</sub> Deep Power-Down Current	1		20	μA	RP#=GND±0.2V I <sub>OUT</sub> (STS)=0mA
I <sub>CCR</sub>	V <sub>CC</sub> Read Current	1,5,6		60	mA	CMOS Inputs V <sub>CC</sub> =V <sub>CC</sub> Max., CE#=GND f=8MHz, I <sub>OUT</sub> =0mA
				75	mA	TTL Inputs $V_{CC}=V_{CC}Max., CE\#=V_{IL}$ f=8MHz, I <sub>OUT</sub> =0mA
I <sub>CCW</sub>	V <sub>CC</sub> Write Current ((Multi) W/B Write or Set Block Lock Bit)	1,7		35	mA	V <sub>PP</sub> =5.0V±0.5V
I <sub>CCE</sub>	V <sub>CC</sub> Erase Current (Block Erase, Full Chip Erase, Clear Block Lock Bits)	1,7		30	mA	V <sub>PP</sub> =5.0V±0.5V
I <sub>CCWS</sub> I <sub>CCES</sub>	V <sub>CC</sub> Write or Block Erase Suspend Current	1,2	1	10	mA	CE#=V <sub>IH</sub>
IPPS	V <sub>PP</sub> Standby Current	1	±2	±15	μA	V <sub>PP</sub> ≤V <sub>CC</sub>
I <sub>PPR</sub>	V <sub>PP</sub> Read Current	1	10	200	μA	V <sub>PP</sub> >V <sub>CC</sub>
I <sub>PPD</sub>	V <sub>PP</sub> Deep Power-Down Current	1	0.1	5	μA	RP#=GND±0.2V
I <sub>PPW</sub>	V <sub>PP</sub> Write Current ((Multi) W/B Write or Set Block Lock Bit)	1,7		80	mA	V <sub>PP</sub> =5.0V±0.5V
I <sub>PPE</sub>	V <sub>PP</sub> Erase Current (Block Erase, Full Chip Erase, Clear Block Lock Bits)	1,7		40	mA	V <sub>PP</sub> =5.0V±0.5V
I <sub>PPWS</sub> I <sub>PPES</sub>	V <sub>PP</sub> Write or Block Erase Suspend Current	1	10	200	μA	V <sub>PP</sub> =V <sub>PPH1</sub>

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			Vcc	≥=5V		Test
Symbol	Parameter	Notes	Min.	Max.	Unit	Conditions
V <sub>IL</sub>	Input Low Voltage	7	-0.5	0.8	V	
V <sub>IH</sub>	Input High Voltage	7	2.0	V <sub>CC</sub> +0.5	V	
V <sub>OL</sub>	Output Low Voltage	3,7		0.45	V	V <sub>CC</sub> =V <sub>CC</sub> Min. I <sub>OL</sub> =5.8mA
V <sub>OH1</sub>	Output High Voltage (TTL)	3,7	2.4		V	V <sub>CC</sub> =V <sub>CC</sub> Min. I <sub>OH</sub> =-2.5mA
V <sub>OH2</sub>	Output High Voltage (CMOS)	3,7	0.85 V <sub>CC</sub>		V	V <sub>CC</sub> =V <sub>CC</sub> Min. I <sub>OH</sub> =-2.5mA
			V <sub>CC</sub> -0.4		V	V <sub>CC</sub> =V <sub>CC</sub> Min. I <sub>OH</sub> =-100µA
V <sub>PPLK</sub>	V <sub>PP</sub> Lockout during Normal Operations	4,7		1.5	V	
V <sub>PPH1</sub>	V <sub>PP</sub> during Write or Erase Operations		4.5	5.5	V	
V <sub>LKO</sub>	V <sub>CC</sub> Lockout Voltage		2.0		V	

NOTES:

1. All currents are in RMS unless otherwise noted. Typical values at nominal V<sub>CC</sub> voltage and T<sub>A</sub>=+25°C.These currents are valid for all product versions (packages and speeds).

2. I<sub>CCWS</sub> and I<sub>CCES</sub> are specified with the device de-selected. If read or byte written while in erase suspend mode, the device's current draw is the sum of I<sub>CCWS</sub> or I<sub>CCES</sub> and I<sub>CCB</sub> or I<sub>CCW</sub>, respectively.

3. Includes STS.

4. Block erases, full chip erases, (multi) word/byte writes and block lock-bit configurations are inhibited when V<sub>PP</sub>≤V<sub>PPLK</sub>, and not guaranteed in the range between V<sub>PPLK</sub>(max.) and V<sub>PPH1</sub>(min.) and above V<sub>PPH1</sub>(max.).
5. Automatic Power Savings (APS) reduces typical I<sub>CCR</sub> to 1mA at 5V V<sub>CC</sub> in static operation.
6. CMOS inputs are either V<sub>CC</sub>±0.2V or GND±0.2V. TTL inputs are either V<sub>IL</sub> or V<sub>IH</sub>.
7. Sampled, not 100% tested.

## 6.2.4 AC CHARACTERISTICS - READ-ONLY OPERATIONS<sup>(1)</sup>

		<sub>CC</sub> =5V±0.25V <sub>CC</sub> =5V±0.5V	LH28F320	0S5-L90 <sup>(5)</sup>		320S5- 00 <sup>(6)</sup>	
Sym.	Parameter	Notes	Min.	Max.	Min.	Max.	Unit
AVAV	Read Cycle Time		90		100		ns
AVQV	Address to Output Delay			90		100	ns
ELQV	CE# to Output Delay	2		90		100	ns
PHQV	RP# High to Output Delay			400		400	ns
GLQV	OE# to Output Delay	2		30		35	ns
ELQX	CE# to Output in Low Z	3	0		0		ns
EHQZ	CE# High to Output in High Z	3		25		30	ns
GLQX	OE# to Output in Low Z	3	0		0		ns
GHQZ	OE# High to Output in High Z	3		10		10	ns
ОН	Output Hold from Address, CE# or OE# Change, Whichever Occurs Fir	st 3	0		0		ns
FLQV FHQV	BYTE# to Output Delay	3		90		100	ns
FLQZ	BYTE# to Output in High Z	3		25		30	ns
ELFL ELFH	CE# Low to BYTE# High or Low	3		5		5	ns

#### NOTES:

1. See AC Input/Output Reference Waveform for maximum allowable input slew rate.

2. OE# may be delayed up to t<sub>FLOV</sub>-t<sub>GLOV</sub> after the falling edge of CE# without impact on t<sub>FLOV</sub>.

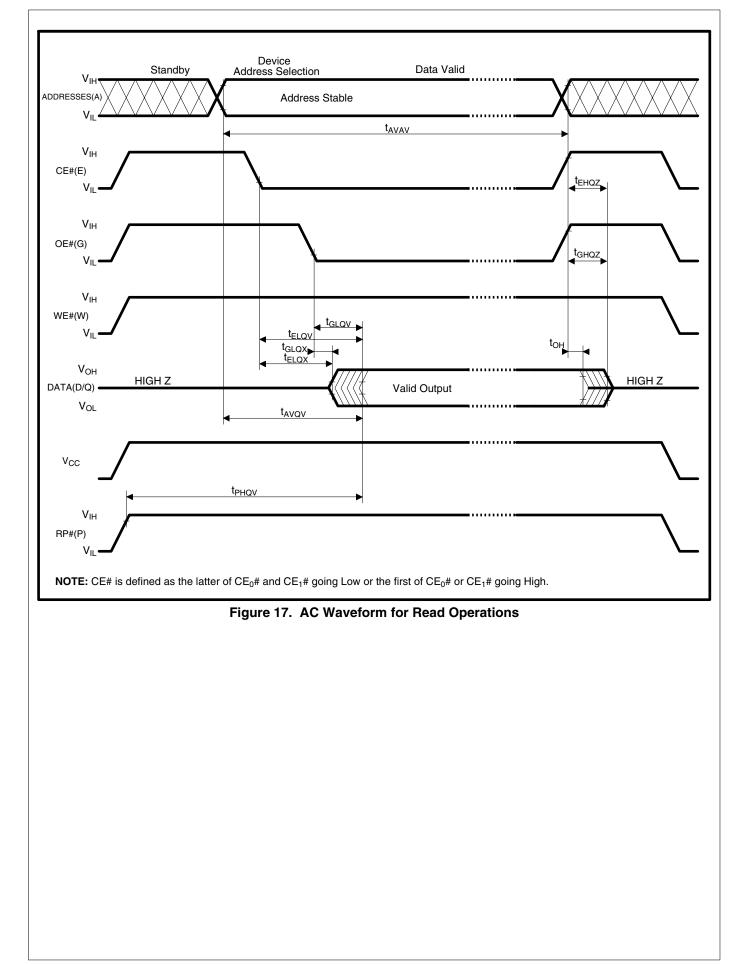
3. Sampled, not 100% tested.

4. See Ordering Information for device speeds (valid operational combinations).

5. See Transient Input/Output Reference Waveform and Transient Equivalent Testing Load Circuit (High Speed Configuration) for testing characteristics.

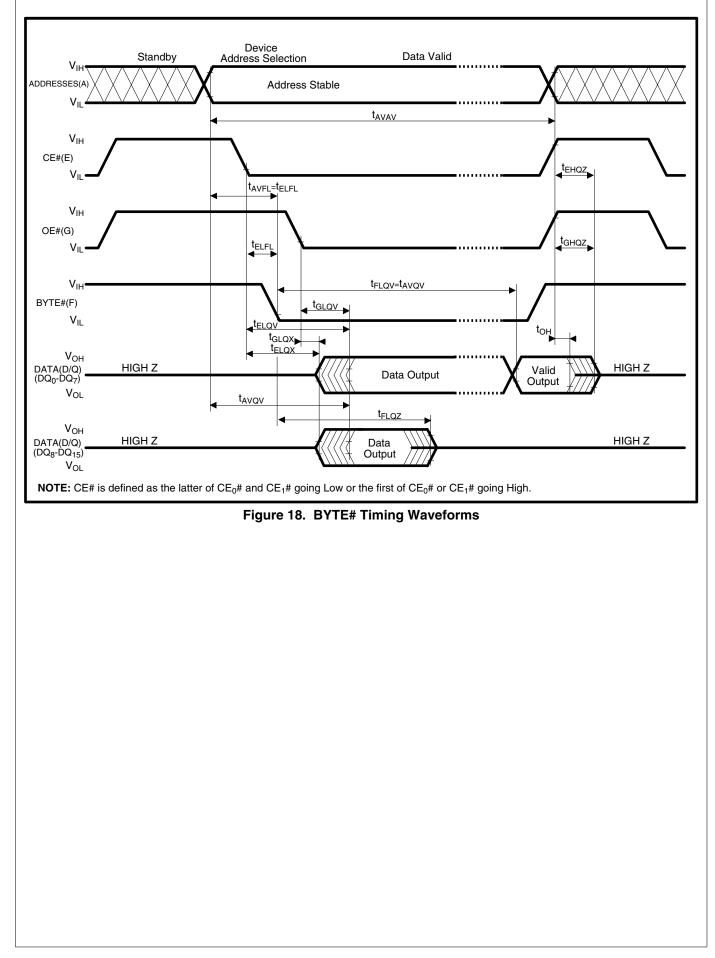
6. See Transient Input/Output Reference Waveform and Transient Equivalent Testing Load Circuit (Standard Configuration) for testing characteristics.

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	V	0.5V, 5V±0.25\ / <sub>CC</sub> =5V±0.25V	LH28F320	0S5-L90 <sup>(6)</sup>			_
Versions <sup>(5)</sup>		/ <sub>CC</sub> =5V±0.5V				320S5- )0 <sup>(7)</sup>	
Sym.	Parameter	Notes	Min.	Max.	Min.	Max.	Unit
t <sub>AVAV</sub>	Write Cycle Time		90		100		ns
t <sub>PHWL</sub>	RP# High Recovery to WE# Going Low	2	1		1		μs
t <sub>ELWL</sub>	CE# Setup to WE# Going Low		10		10		ns
t <sub>WLWH</sub>	WE# Pulse Width		40		40		ns
t <sub>SHWH</sub>	WP# VIH Setup to WE# Going High	1 2	100		100		ns
t <sub>VPWH</sub>	V <sub>PP</sub> Setup to WE# Going High	2	100		100		ns
t <sub>AVWH</sub>	Address Setup to WE# Going High	3	40		40		ns
t <sub>DVWH</sub>	Data Setup to WE# Going High	3	40		40		ns
t <sub>WHDX</sub>	Data Hold from WE# High		5		5		ns
t <sub>WHAX</sub>	Address Hold from WE# High		5		5		ns
t <sub>WHEH</sub>	CE# Hold from WE# High		10		10		ns
tWHWL			30		30		ns
t <sub>WHRL</sub>				90		90	ns
t <sub>WHGL</sub>	Write Recovery before Read		0		0		ns
t <sub>QVVL</sub>	V <sub>PP</sub> Hold from Valid SRD, STS Hig		0		0		ns
t <sub>QVSL</sub>	WP# V <sub>IH</sub> Hold from Valid SRD, STS High Z	5 2,4	0		0		ns

## 6.2.5 AC CHARACTERISTICS - WRITE OPERATIONS<sup>(1)</sup>

## NOTES:

1. Read timing characteristics during block erase, full chip erase, (multi) wrod/byte write and block lock-bit configuration operations are the same as during read-only operations. Refer to AC Characteristics for read-only operations.

2. Sampled, not 100% tested.

- 3. Refer to Table 4 for valid A<sub>IN</sub> and D<sub>IN</sub> for block erase, full chip erase, (multi) word/byte write or block lock-bit configuration.
- 4. V<sub>PP</sub> should be held at V<sub>PPH1</sub> until determination of block erase, full chip erase, (multi) word/byte write or block lock-bit configuration success (SR.1/3/4/5=0).
- 5. See Ordering Information for device speeds (valid operational combinations).

6. See Transient Input/Output Reference Waveform and Transient Equivalent Testing Load Circuit (High Seed Configuration) for testing characteristics.

7. See Transient Input/Output Reference Waveform and Transient Equivalent Testing Load Circuit (Standard Configuration) for testing characteristics.

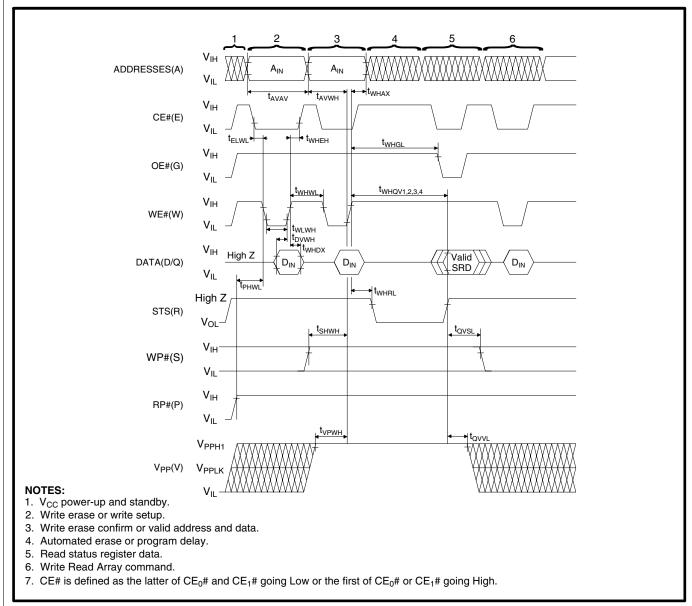


Figure 19. AC Waveform for WE#-Controlled Write Operations

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	V	<sub>CC</sub> =5V±0.25V	LH28F32	0S5-L90 <sup>(6)</sup>			
Versions <sup>(5)</sup>		V <sub>CC</sub> =5V±0.5V		LH28F320S5- L100 <sup>(7)</sup>			
Sym.	Parameter	Notes	Min.	Max.	Min.	Max.	Unit
t <sub>AVAV</sub>	Write Cycle Time		90		100		ns
t <sub>PHEL</sub>	RP# High Recovery to CE# Going L	.ow 2	1		1		μs
t <sub>WLEL</sub>	WE# Setup to CE# Going Low		0		0		ns
t <sub>ELEH</sub>	CE# Pulse Width		50		50		ns
t <sub>SHEH</sub>	WP# V <sub>IH</sub> Setup to CE# Going High	2	100		100		ns
t <sub>VPEH</sub>	V <sub>PP</sub> Setup to CE# Going High	2	100		100		ns
t <sub>AVEH</sub>	Address Setup to CE# Going High	3	40		40		ns
t <sub>DVEH</sub>	Data Setup to CE# Going High	3	40		40		ns
t <sub>EHDX</sub>	Data Hold from CE# High		5		5		ns
t <sub>EHAX</sub>	Address Hold from CE# High		5		5		ns
t <sub>EHWH</sub>	WE# Hold from CE# High		0		0		ns
	CE# Pulse Width High		25		25		ns
t <sub>EHRL</sub>	CE# High to STS Going Low			90		90	ns
EHGL	Write Recovery before Read		0		0		ns
t <sub>QVVL</sub>	V <sub>PP</sub> Hold from Valid SRD, STS High	ר Z 2,4	0		0		ns
t <sub>QVSL</sub>	WP# V <sub>IH</sub> Hold from Valid SRD, STS High Z	2,4	0		0		ns

## 6.2.6 ALTERNATIVE CE#-CONTROLLED WRITES<sup>(1)</sup>

### NOTES:

1. In systems where CE# defines the write pulse width (within a longer WE# timing waveform), all setup, hold and inactive WE# times should be measured relative to the CE# waveform.

2. Sampled, not 100% tested.

3. Refer to Table 4 for valid A<sub>IN</sub> and D<sub>IN</sub> for block erase, full chip erase, (multi) word/byte write or block lock-bit configuration.

4. V<sub>PP</sub> should be held at V<sub>PPH1</sub> until determination of block erase, full chip erase, (multi) word/byte write or block lock-bit configuration success (SR.1/3/4/5=0).

5. See Ordering Information for device speeds (valid operational combinations).

6. See Transient Input/Output Reference Waveform and Transient Equivalent Testing Load Circuit (High Seed Configuration) for testing characteristics.

7. See Transient Input/Output Reference Waveform and Transient Equivalent Testing Load Circuit (Standard Configuration) for testing characteristics.

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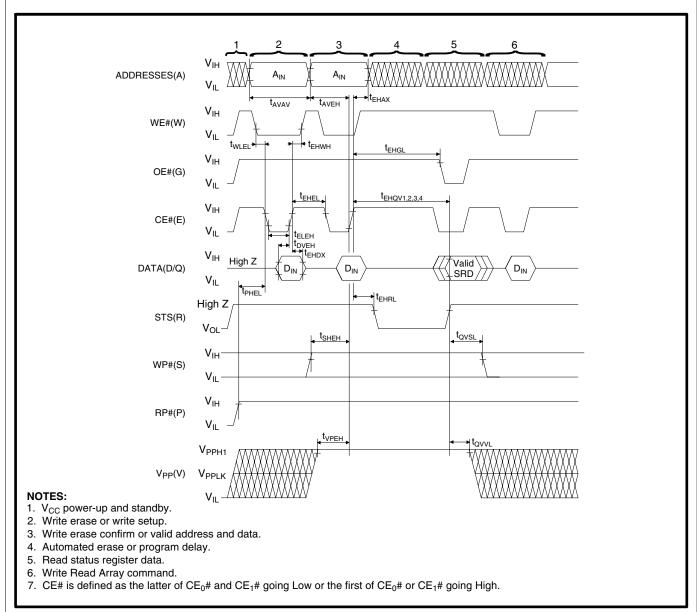
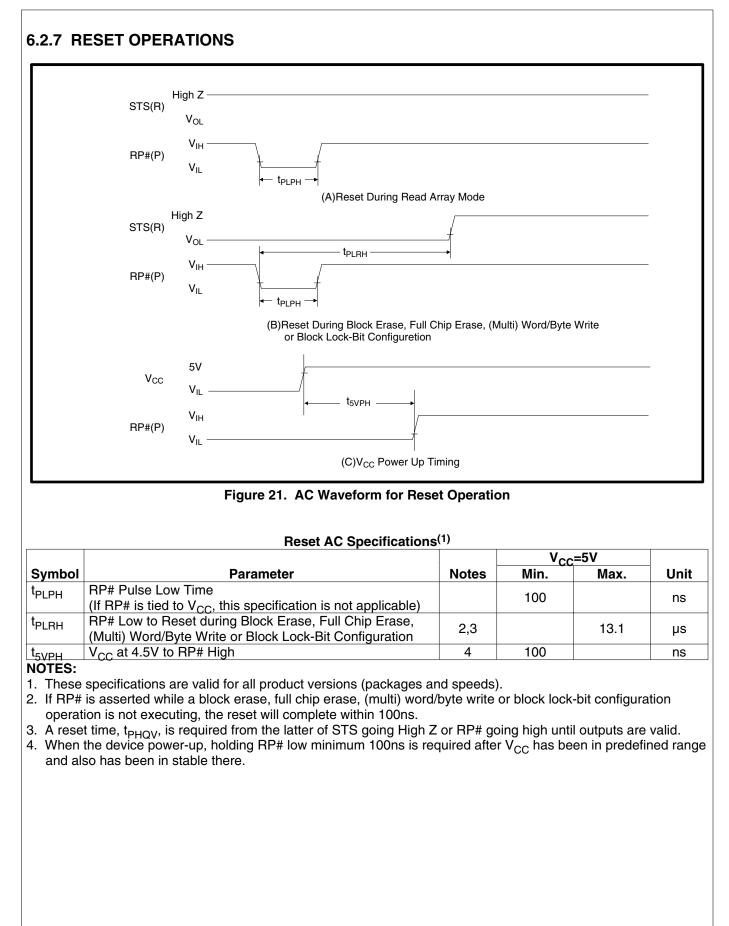


Figure 20. AC Waveform for CE#-Controlled Write Operations



## 6.2.8 BLOCK ERASE, FULL CHIP ERASE, (MULTI) WORD/BYTE WRITE AND BLOCK LOCK-BIT CONFIGURATION PERFORMANCE<sup>(3,4)</sup>

			V <sub>PP</sub> =4.5	5V-5.5V	
Sym.	Parameter	Notes	Typ.(1)	Max.	Unit
t <sub>WHQV1</sub> t <sub>EHQV1</sub>	Word/Byte Write Time (using W/B write, in word mode)	2	9.24	120	μs
twhqv1 t <sub>EHQV1</sub>	Word/Byte Write Time (using W/B write, in byte mode)	2	9.24	120	μs
	Word/Byte Write Time (using multi word/byte write)	2	2	120	μs
	Block Write Time (using W/B write, in word mode)	2	0.31	3.7	S
	Block Write Time (using W/B write, in byte mode)	2	0.61	7.5	s
	Block Write Time (using multi word/byte write)	2	0.13	1.5	S
t <sub>WHQV2</sub> t <sub>EHQV2</sub>	Block Erase Time	2	0.34	10	s
	Full Chip Erase Time		21.8	640	S
t <sub>WHQV3</sub> t <sub>EHQV3</sub>	Set Block Lock-Bit Time	2	9.24	120	μs
t <sub>WHQV4</sub> t <sub>EHQV4</sub>	Clear Block Lock-Bits Time	2	0.34	10	s
WHRH1 EHRH1	Write Suspend Latency Time to Read		5.6	7	μs
WHRH2 EHRH2	Erase Suspend Latency Time to Read		9.4	13.1	μs

#### .. 7000

### NOTES:

1. Typical values measured at T<sub>A</sub>=+25°C and nominal voltages. Assumes corresponding block lock-bits are not set. Subject to change based on device characterization.

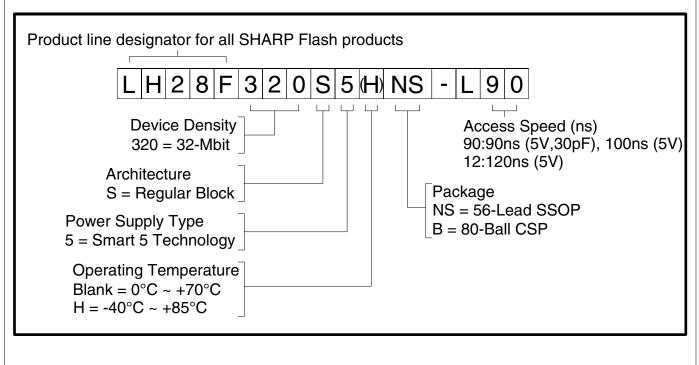
2. Excludes system-level overhead.

3. These performance numbers are valid for all speed versions.

4. Sampled but not 100% tested.

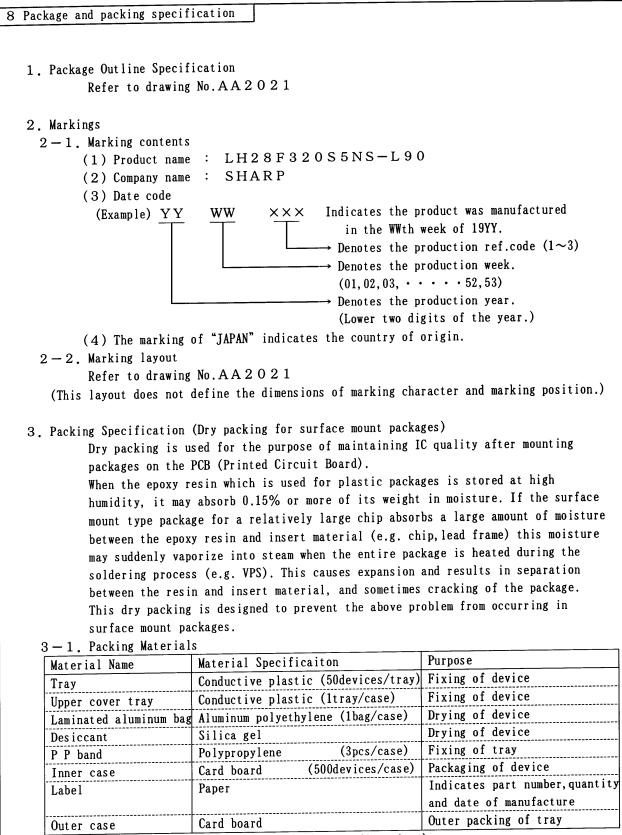
## 7 ADDITIONAL INFORMATION

## 7.1 Ordering Information



		Valid Operational Combinations			
		V <sub>CC</sub> =5V±0.5V	V <sub>CC</sub> =5V±0.25V		
		100pF load,	30pF load,		
Option	Order Code	TTL I/O Levels	1.5V I/O Levels		
1	LH28F320S5NS-L90	LH28F320S5-L100	LH28F320S5-L90		

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(Devices shall be placed into a tray in the same direction.)

	utline dimension of tray
R	efer to attached drawing
4. Storag	e and Opening of Dry Packing
4 - 1.5	tore under conditions shown below before opening the dry packing
(	1) Temperature range : $5\sim40^{\circ}$ C
(	(2) Humidity : 80% RH or less
4 - 2. N	lotes on opening the dry packing
(	(1) Before opening the dry packing, prepare a working table which is
	grounded against ESD and use a grounding strap.
(	(2) The tray has been treated to be conductive or anti-static. If the
	device is transferred to another tray, use a equivalent tray.
4 - 3.5	storage after opening the dry packing
F	Perform the following to prevent absorption of moisture after opening.
	(1) After opening the dry packing, store the ICs in an environment with a $1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 $
	temperature of $5\sim 25^{\circ}$ and a relative humidity of 60% or less and
	mount ICs within 3 days after opening dry packing. (2) To re-store the ICs for an extended period of time within 3 days after
	(2) To re-store the ICs for an extended period of time within 5 days after opening the dry packing, use a dry box or re-seal the ICs in the dry
	packing with desiccant (whoes indicater is blue), and store in an
	environment with a temperature of $5{\sim}40^\circ{\rm C}$ and a relative humidity of
	80% or less, and mount ICs within 2 weeks.
1	(3) Total period of storage after first opening and re-opening is within
	3 days, and store the ICs in the same environment as section 4-3.(1).
	First opening $\leftarrow X_1 \cdots \rightarrow re\text{-sealing} \leftarrow Y \cdots \rightarrow re\text{-opening} \leftarrow X_2 \cdots \rightarrow mounting$
	Its in dry 5,020 c and 500 DH ar loss
	packing 60%RH or less 80%RH or less 60%RH or less
	$X_1 + X_2$ : within 3 days
	Y :within 2 weeks
4-4. Ba	king (drying) before mounting
(1	) Baking is necessary
	(A) If the humidity indicator in the desiccant becomes pink
	(B) If the procedure in section $4-3$ could not be performed
(2	2) Recommended baking conditions
	If the above conditions (A) and (B) are applicable, bake it before mounting. The recommended conditions are $16 \sim 24$ hours at $120^{\circ}$ C.
	Heat resistance tray is used for shipping tray.
1.4	Heat resistance tray is used for shipping tray. 3) Storage after baking
( (	After baking ICs, store the ICs in the same environment as section
	4-3.(1).

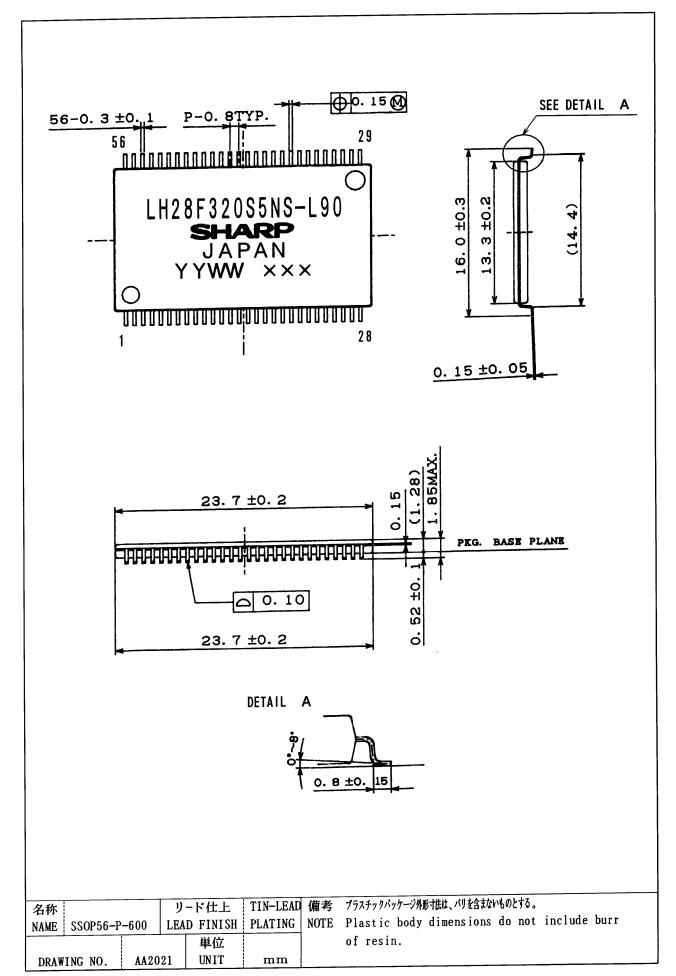
#### 5. Surface Mount Conditions

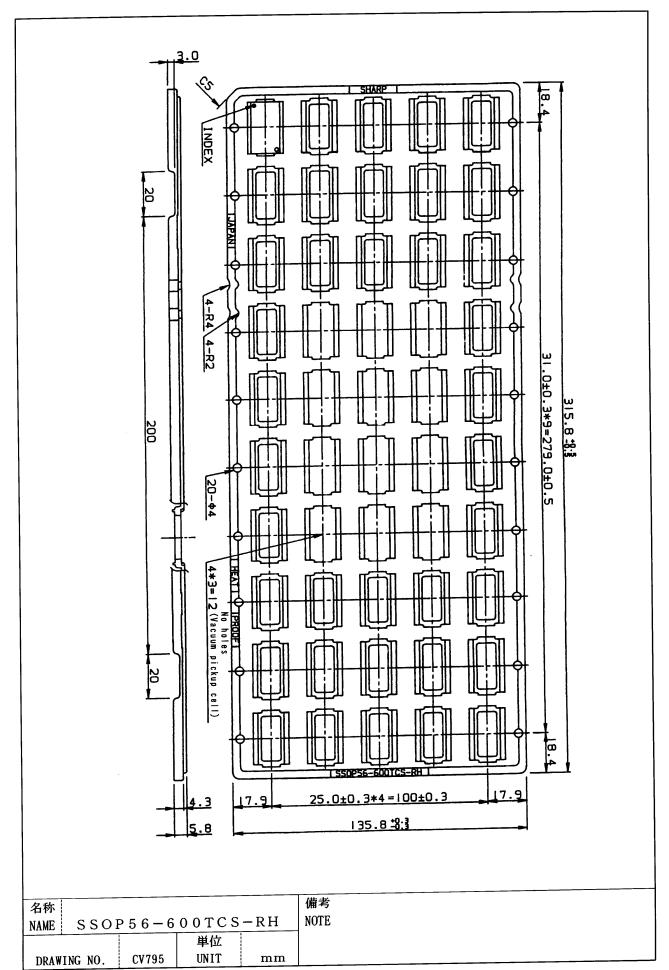
Please perform the following conditions when mounting ICs not to deteriorate IC quality.

## 5-1 . Soldering conditions (The following conditions are valid only for one time soldering.)

Mounting Method	Temperature and Duration	Measurement Point
Reflow soldering	Peak temperature of 240°C or less,	IC package surface
(air)	duration of less than 15 seconds	
	above 230℃, temperature	
	increase rate of $1 \sim 4^{\circ}$ C/second	
Manual soldering	260℃ or less, duration of less	IC outer lead surface
(soldering iron)	than 10 seconds	

- 5-2. Conditions for removal of residual flux
  - (1) Ultrasonic washing power : 25 Watts/liter or less
  - (2) Washing time
- : Total 1 minute maximum
- (3) Solvent temperature :  $15 \sim 40^{\circ}$ C

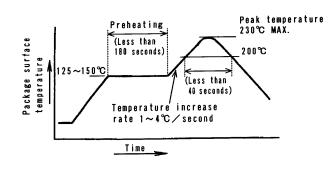




«Supplementary data»

	LHF32K10
Recommended mounting	g conditions for two time reflow soldering .
Product name(Package)	LH28F320S5NS-L90 (SSOP56-P-600)
Packing specification	Tray (Dry packing)
Mounting method	Reflow soldering (Air)
Reflow soldering conditions	Peak temperature of 230°C or less. 200°C or over, duration of less than 40 seconds. Preheat temperature of 125~150°C, duration of less than 180 seconds. Temperature increase rate of 1~4°C/second.
Measurement point	IC package surface
Storage conditions	After opening the dry packing, store the ICs in an environment with a temperature of $5\sim25^{\circ}$ and a relative humidity of 60% or less. If doing reflow soldering twice, do the first reflow soldering within 72 hours after opening dry packing and do the second reflow soldering within 72 hours after the first reflow soldering.
Note	If the above storage conditions are not applicable, bake it before reflow soldering. The recommended conditions are 16~24 hours at 120℃. (Heat resistance tray is used for shipping tray.)

## Recommended Reflow Soldering(Air) Temperature Profile



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(NO. 980727-X8)

#### LHF32K10

#### Flash memory LHFXXKXX family Data Protection

Noises having a level exceeding the limit specified in the specification may be generated under specific operating conditions on some systems.

Such noises, when induced onto WE# signal or power supply, may be interpreted as false commands, causing undesired memory updating.

To protect the data stored in the flash memory against unwanted overwriting, systems operating with the flash memory should have the following write protect designs, as appropriate:

### 1) Protecting data in specific block

Setting the lock bit of the desired block and pulling WP# low disables the writing operation on that block. By using this feature, the flash memory space can be divided into, for example, the program section(locked section) and data section(unlocked section).

By controlling WP#, desired blocks can be locked/unlocked through the software. For further information on setting/resetting block bit, refer to the specification. (See chapter 4.12 and 4.13.)

#### 2) Data protection through Vpp

When the level of Vpp is lower than VPPLK (lockout voltage), write operation on the flash memory is disabled. All blocks are locked and the data in the blocks are completely write protected.

For the lockout voltage, refer to the specification. (See chapter 6.2.3.)

#### 3) Data protection through RP#

When the RP# is kept low during power up and power down sequence such as voltage transition, write operation on the flash memory is disabled, write protecting all blocks.

For the details of RP# control, refer to the specification. (See chapter 5.6 and 6.2.7.)

## A-1 RECOMMENDED OPERATING CONDITIONS

### A-1.1 At Device Power-Up

AC timing illustrated in Figure A-1 is recommended for the supply voltages and the control signals at device power-up. If the timing in the figure is ignored, the device may not operate correctly.

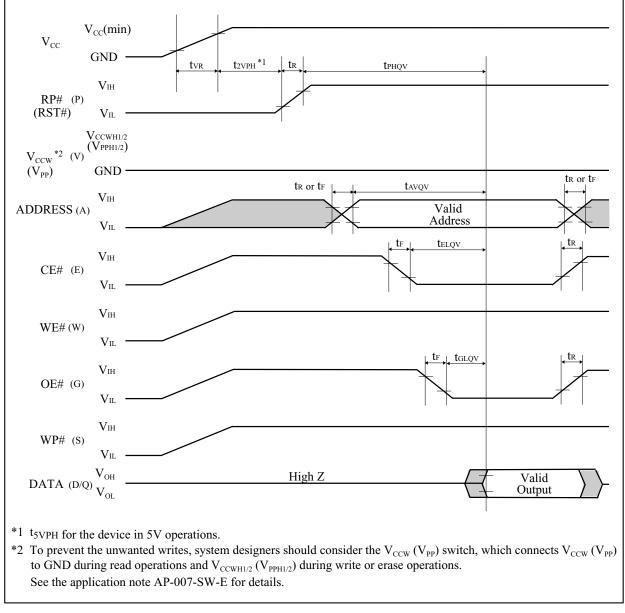


Figure A-1. AC Timing at Device Power-Up

For the AC specifications  $t_{VR}$ ,  $t_R$ ,  $t_F$  in the figure, refer to the next page. See the "ELECTRICAL SPECIFICATIONS" described in specifications for the supply voltage range, the operating temperature and the AC specifications not shown in the next page.

## A-1.1.1 Rise and Fall Time

Symbol	Parameter	Notes	Min.	Max.	Unit
t <sub>VR</sub>	V <sub>CC</sub> Rise Time	1	0.5	30000	μs/V
t <sub>R</sub>	Input Signal Rise Time			1	μs/V
t <sub>F</sub>	Input Signal Fall Time			1	μs/V

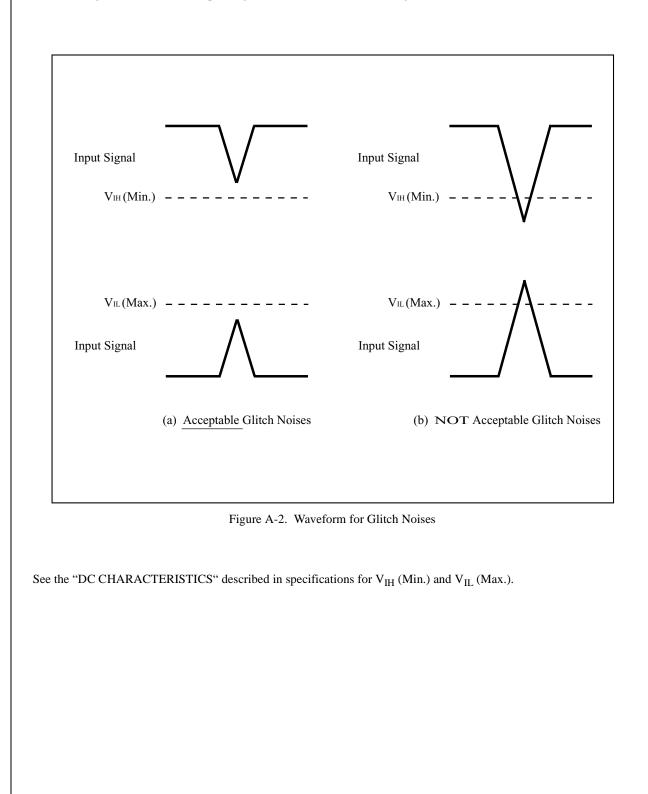
#### NOTES:

1. Sampled, not 100% tested.

2. This specification is applied for not only the device power-up but also the normal operations.  $t_R(Max.)$  and  $t_F(Max.)$  for RP# (RST#) are 100 $\mu$ s/V.

### A-1.2 Glitch Noises

Do not input the glitch noises which are below  $V_{IH}$  (Min.) or above  $V_{IL}$  (Max.) on address, data, reset, and control signals, as shown in Figure A-2 (b). The acceptable glitch noises are illustrated in Figure A-2 (a).



## A-2 RELATED DOCUMENT INFORMATION<sup>(1)</sup>

Document No.	Document Name
AP-001-SD-E	Flash Memory Family Software Drivers
AP-006-PT-E	Data Protection Method of SHARP Flash Memory
АР-007-SW-Е	RP#, V <sub>PP</sub> Electric Potential Switching Circuit

NOTE:

1. International customers should contact their local SHARP or distribution sales office.

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